# An analogue of Schensted's bumping algorithm in affine type $A$ 

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## Fock spaces representations of $\mathcal{U}_{q}\left(\widehat{\mathfrak{S L}_{e}}\right)$

Let $e \in \mathbb{Z}_{>1}, q$ an indeterminate.
$\mathcal{U}_{q}(\widehat{\mathfrak{s l}})=q$-deformation of the Lie algebra of affine type $A$.

- Generators $e_{i}, f_{i}, t_{i}^{ \pm 1}, \mathfrak{d}$ for $i=0 \ldots e-1$
- Some relations.


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Let $I \in \mathbb{Z}_{>0}$ and $\mathbf{s} \in \mathbb{Z}^{\prime}$.

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Theorem (Jimbo, Misra, Miwa, Okado 1991)
$\mathcal{F}_{\mathbf{s}}$ is an (integrable) $\mathcal{U}_{q}(\widehat{\mathfrak{s l}})$-module.

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\hline & & \\
\hline
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\hline & \\
& \\
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- Symbol of $\boldsymbol{\lambda}$ of shape $\mathbf{s}$ :

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\left(\begin{array}{lllllllll}
\ldots & -4 & -3 & -2 & & & & & \\
\ldots & -4 & -3 & -2 & -1 & 0 & 1 & 2 & 3 \\
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\left(\begin{array}{lllllll}
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\mathcal{S}(\boldsymbol{\lambda}, \mathbf{s})=\left(\begin{array}{llllllll}
0 & 4 & & & & & \\
0 & 1 & 2 & 3 & 4 & 7 & 8 \\
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## Definition

$\mathcal{S}(\boldsymbol{\lambda}, \mathbf{s})$ is semistandard if $s_{c} \leq s_{c+1}$ for all $c=1, \ldots, /$ and if the columns (resp. rows) of $\mathcal{S}(\boldsymbol{\lambda}, \mathbf{s})$ are non-decreasing (resp. increasing).

## Crystal graphs

According to Kashiwara, $\mathcal{F}_{\mathbf{s}}$ has a crystal graph $B\left(\mathcal{F}_{\mathbf{s}}\right)$.

- Vertices $=$ all $l$-partitions.
- Edges of $B=$ directed arrows colored by $i \in\{0, \ldots e-1\}$.


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## Properties

- $B\left(\mathcal{F}_{\mathbf{s}}\right)=\bigsqcup B$ (connected components).
- Each $B$ has a unique vertex $\dot{\boldsymbol{\lambda}}$ with zero-indegree : the source vertex. We denote $B=B(\boldsymbol{\lambda}, \mathbf{s})$.
- Each $\boldsymbol{\lambda} \in B(\dot{\boldsymbol{\lambda}}, \mathbf{s})$ writes $\dot{\boldsymbol{\lambda}} \xrightarrow{i_{1}} \ldots \xrightarrow{i_{p}} \boldsymbol{\lambda}$
- $\emptyset$ is always a source vertex.
- There is a natural graph isomorphism between $B(\dot{\lambda}, \mathbf{s})$ and $B(\emptyset, \mathbf{r})$ for some $\mathbf{r} \in \mathscr{S}_{e}$ where $\mathscr{S}_{e}=\left\{\mathbf{t} \in \mathbb{Z}^{\prime} \mid 0 \leq t_{d}-t_{c}<e\right.$ for $\left.c<d\right\}$.


## An example

Take $e=3$,
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This crystal graph is isomorphic to $B(\emptyset, \mathbf{r})$ with $\mathbf{r}=(0,1)$ :


## Notation

We write $|\boldsymbol{\lambda}, \mathbf{s}\rangle \sim|\boldsymbol{\nu}, \mathbf{t}\rangle$ if $B(\dot{\boldsymbol{\lambda}}, \mathbf{s}) \simeq B(\dot{\boldsymbol{\nu}}, \mathbf{t})$ and if $\boldsymbol{\lambda}$ and $\boldsymbol{\nu}$ appear at the same place in their respective crystal.

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Problem: Give a combinatorial description of the relation $\sim$. In other terms, starting from arbitrary $|\boldsymbol{\lambda}, \mathbf{s}\rangle$, determine $\mathbf{r} \in \mathscr{S}_{e}$ and $\boldsymbol{\mu} \in B(\emptyset, \mathbf{r})$ such that $|\boldsymbol{\lambda}, \mathbf{s}\rangle \sim|\boldsymbol{\mu}, \mathbf{r}\rangle$.

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$\rightsquigarrow$ Interesting because only $B(\emptyset, \mathbf{r})$ has an explicit combinatorial description (Foda-Leclerc-Okado-Thibon-Welsh).
$\rightsquigarrow$ Natural interpretations in terms of representations of the complex reflection group $G(I, 1, n)$ (Lascoux-Leclerc-Thibon, Ariki, Shan, Losev).

## $1^{\text {st }}$ tool: the cyclage

For $\mathbf{s}=\left(s_{1}, \ldots, s_{l}\right)$ and $\boldsymbol{\lambda}=\left(\lambda^{1}, \ldots, \lambda^{\prime}\right)$, we define

$$
\xi(\mathbf{s})=\left(s_{l}-e, s_{1}, \ldots, s_{l-1}\right) \quad \text { and } \quad \xi(\lambda)=\left(\lambda^{\prime}, \lambda^{1}, \ldots, \lambda^{\prime-1}\right) .
$$

Example: Take $e=3, \mathbf{s}=(2,0,1)$ and $\boldsymbol{\lambda}=\left(3.2,1,4^{2}\right)$. Then

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\xi(\mathbf{s})=(-2,2,0) \quad \text { and } \quad \xi(\boldsymbol{\lambda})=\left(4^{2}, 3.2,1\right) .
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## Proposition

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Aim: construct a semistandard symbol, starting from an arbitrary symbol.
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Theorem (Kashiwara 90's)

$$
|\boldsymbol{\lambda}, \mathbf{s}\rangle \sim|\nu, \mathbf{t}\rangle .
$$

Denote RS : $|\boldsymbol{\lambda}, \mathbf{s}\rangle \longmapsto|\boldsymbol{\nu}, \mathbf{t}\rangle$.

There exists $M \in \mathbb{N}$ such that $|\boldsymbol{\nu}, \mathbf{t}\rangle=(\mathbf{R S} \circ \xi)^{M}(|\boldsymbol{\lambda}, \mathbf{s}\rangle)$ verifies:
(1) $\mathcal{S}(\boldsymbol{\nu}, \mathbf{t})$ is semistandard, and
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## Theorem (Foda-Leclerc-Okado-Thibon-Welsh 1999)

$|\boldsymbol{\nu}, \mathbf{t}\rangle \in B(\emptyset, \mathbf{t})$ except if $\boldsymbol{\nu}$ contains e parts of the same size $\alpha>0$ such that

$$
\left\{\alpha-a_{i}+s_{c_{i}} \bmod e ; i=1, \ldots, e\right\}=\{0, \ldots, e-1\}
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where $a_{i}=$ row fo the $i$-th part, and $c_{i}=$ component of the $i$-th part.

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$\rightsquigarrow$ The numbers $\alpha-a_{i}+s_{c_{i}}$ are the contents of the rightmost boxes in the parts of size $\alpha$.
$\rightsquigarrow$ How to get rid of the "bad" parts?
$3^{\text {rd }}$ tool: The reduction isomorphism
Let $e=4, I=3, \mathbf{t}=(5,6,8)$, and $\boldsymbol{\nu}=\left(4^{2} .2 .1^{2}, 3.2^{2} .1^{2}, 4.2^{3} .1\right)$.
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| 5 | 6 | 7 | 8 |
| :---: | :---: | :---: | :---: |
| 4 | 5 | 6 | 7 |
| 3 | 4 |  |  |
| 2 |  |  |  |
| 1 |  |  |  |


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| 5 | 6 | 7 | 8 |
| :--- | :--- | :--- | :--- |
| 4 | 5 | 6 | 7 |


| 6 | 7 | 8 |
| :--- | :--- | :--- |$\quad$| 8 | 9 | 10 | 11 |
| :--- | :--- | :--- | :--- | :--- |


| 3 | 4 |
| :--- | :--- |
| 2 |  |
| 1 |  |
|  |  |


| 3 |
| :--- |
| 2 |


| 5 | 6 |
| :--- | :--- |
| 4 |  |
|  |  |

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Theorem (G. 2013)

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$\rightsquigarrow$ There exists $N \in \mathbb{N}$ such that $\left|\rho^{N}(\boldsymbol{\nu}), \xi^{N}(\mathbf{t})\right\rangle \in B\left(\emptyset, \xi^{N}(\mathbf{t})\right)$.

## Analogy with the non-affine case

|  | Finite type $A$ | Affine type $A$ |
| :---: | :---: | :---: |
| Quantum group | $\mathcal{U}_{q}\left(\mathfrak{s l}_{e}\right)$ | $\mathcal{U}_{q}\left(\widehat{\widehat{s l}_{e}}\right)$ |
| Fock space | $\begin{aligned} & \text { Basis }=\text { Young } \\ & \text { tableaux } \end{aligned}$ | Basis = charged l-partitions/symbols |
| Interesting connected component of the crystal | $\begin{gathered} \text { Vertices }= \\ \text { semistandard tableaux } \end{gathered}$ | Vertices $=$ FLOTW symbols |
| Crystal isomorphisms | RS for tableaux | RS for symbols |
|  |  | cyclage $\xi$ |
|  |  | reduction $\rho$ |
| Equivalence relation expected | $\begin{gathered} \mathbf{T}_{1} \sim \mathbf{T}_{2} \text { iff } \\ \mathbf{R S}\left(\mathbf{T}_{1}\right)=\mathbf{R S}\left(\mathbf{T}_{2}\right) \end{gathered}$ | $\begin{gathered} \boldsymbol{\lambda}_{1} \sim \boldsymbol{\lambda}_{2} \text { iff } \\ \Phi\left(\boldsymbol{\lambda}_{1}\right)=\Phi\left(\boldsymbol{\lambda}_{2}\right) \text { where } \\ \Phi=\rho^{N} \circ(\mathbf{R S} \circ \xi)^{M} \end{gathered}$ |

