## The Ingram Conjecture

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#### Tent maps and inverse limits:

Let  $T_s : I \to I$ ,  $x \mapsto \min\{sx, s(1-x)\}$  be the tent map with slope  $s \in [0, 1]$ .

#### The inverse limit space

 $\varprojlim([0,1],T_s) = \{(x_k)_{k \le 0} : T_s(x_k) = x_{k+1} \forall k < 0\}$ 

is the space of all backward orbits, equipped with product topology.

#### Theorem: (Ingram's Conjecture)

If  $1 \leq s < s' \leq 2$ , then  $\lim_{t \to 0} ([0,1],T_s)$  and  $\lim_{t \to 0} ([0,1],T_{s'})$  are non-homeomorphic.

#### Some History:

Tom Ingram posed the question in the early 1990s, specifically for the case that c is periodic.

Partial results for specific critical behaviour.

Barge & Diamond	three period 5 non-homeo $\exists$ complete algebraic inv.
Bruin	$\frac{\log s}{\log s'}$ irrational $\Rightarrow$ non-home
Kailhofer Block, Jakimovik,	all periodic cases non-homeo
Štimac	all (pre)periodic non-homeo
<b>X</b>	
Raines & Stimac	all non-recurrent non-homed
Barge & Diamond Brucks & Bruin	uncountable many non-homeo recurrent cases

#### Notation:

 $c = \frac{1}{2}$  is critical point.  $c_n = c_n(s) = T_s^n(c)$ I = [0, 1] and  $[c_2, c_1]$  is the **core** of  $T_s$ .

The shift homeomorphism is

$$\sigma(x) = (\dots, x_{-3}, x_{-2}, x_{-1}, x_0, T_s(x_0)).$$

The projection on coordinate -q is

 $\pi_q : \varprojlim([0,1],T_s) \to I, \qquad \pi_q(x) = x_{-q}$ commutes as  $T \circ \pi_q = \pi_q \circ \sigma$ .

#### The zero-composant

 $\mathfrak{C} = \{x \in \varprojlim([0, 1], T_s) : x_{-k} < c \forall k \text{ suff. large}\}$ is a ray disjoint from but converging to the **core** of the inverse limit space  $\varprojlim([c_2, c_1], T_s).$ 

#### Main Result:

If  $s \in (\sqrt{2}, 2]$ , then  $\lim_{t \to \infty} ([c_2, c_1], T_s)$  is indecomposable.

**Theorem:** If  $\sqrt{2} < s \leq 2$ , then every self-homeomorphism

 $h: \varprojlim([0,1],T_s) \rightarrow \varprojlim([0,1],T_s)$ 

is **isotopic** to a power of the shift.

**Remark:** We started proving that h is pseudoisotopic to  $\sigma^R$ , *i.e.*, h permutes the composants of  $\varprojlim([c_2, c_1], T_s)$  in the same way as  $\sigma^R$ .

This makes the Ingram Conjecture easier to prove as well.

### Chains:

Unimodal inverse limits are **chainable**, i.e., for every  $\varepsilon > 0$  there is a chain  $C = \{\ell_j\}_{j=1}^N$  such that

- $\varprojlim([0,1],T_s) \subset \cup_j \ell_j;$
- $\ell_j$  are open;
- $\ell_i \cap \ell_j \neq \emptyset$  if and only if  $|i j| \leq 1$ ;
- $\operatorname{mesh}(C) = \operatorname{max}_j \operatorname{diam}(\ell_j) < \varepsilon.$

We will use *q*-chains constructed as follows:

- $I \subset \cup \{I_j^q\}$ , where  $I_j^q$  are intervals of length  $< \varepsilon s^{-q}/2;$
- $\ell_j^q = \pi_q^{-1}(I_j^q);$
- The resulting chain  $C_q = \{\ell_j^q\}.$

#### More notation:

A point  $x \in \mathfrak{C}$  is a *p*-point if  $x_{-(p+j)} = c$  for some  $j \ge 0$ . The largest such  $j = L_p(x)$  is the *p*-level of *x*.

The set of *p*-points  $E_p$  can be ordered according to arc-length  $\overline{d}$  along  $\mathfrak{C}$ :

$$x \preceq y$$
 if  $\overline{d}(\alpha, x) \leq \overline{d}(\alpha, y)$ 

where  $\alpha$  is the end-point of  $\mathfrak{C}$ .

The folding pattern is the list of *p*-levels of *p*-points in an arc  $A \subset \mathfrak{C}$ :

$$FP(A) = L_p(x_1)L_p(x_2)\dots L_p(x_N)$$

if  $x_1 \preceq x_2 \preceq \cdots \preceq x_N$  are the *p*-points of *A*.

 $s_k$  is the k-th salient p-point if

• 
$$L_p(s_k) = k;$$

•  $L_p(x) < k$  for all  $\alpha \prec x \prec s_k$ .

## *p*-link-symmetry:

**Definition:** Given an arc  $A \subset \mathfrak{C}$  and a *p*-chain  $\mathcal{C}_p$ , let  $\ell^0, \ell^1, \ldots, \ell^N$  be the links successively visited by A. Then A is

• symmetric if the folding pattern FP(A) is a palindrome and  $\pi_p(\partial A)$  is a single point;

• *p*-link-symmetric if  $\ell^j = \ell^{N-j}$  for  $0 \le j \le N$ ;

• maximal *p*-link-symmetric if there is no *p*-link-symmetric arc  $A' \supset A$  passing though more links.

A symmetric arc has a well-defined midpoint.

If  $h(C_q)$  refines  $C_p$ , then *q*-link-symmetric arcs map to *p*-link-symmetric arcs.

#### Main Lemmas on *p*-link-symmetric Arcs:

Assume that c has an infinite orbit.

Let  $\kappa := \min\{i \ge 3 : c_i > c\}.$ 

**Lemma 1:** The maximal *p*-link-symmetric arc  $A_i$  centred at  $s_i$ ,  $i \ge \kappa$ , contains  $\kappa$  salient points, namely

 $s_{i-\kappa+2},\ldots,s_i,s_{i+1}$ 

and  $s_{i-\kappa+2}$  is interior to this arc.

**Lemma 2:** If a *p*-link-symmetric arc *J* is **not** centred at a salient point *y*, say  $s_{i-1} \prec y \prec s_i$ , then *J* contains at most one salient point, and  $J \subset A_i$ .

## How to use *p*-link Symmetry

#### Folding points.

A point is called **folding point** if it has no neighbourhood in the core  $\lim_{t \to \infty} ([c_2, c_1], T_s)$  homeomorphic with a Cantor set of arcs.

Characterisations of folding points:

- x is folding point iff  $x_k \in \omega(c)$  for all k.
- x is folding point iff there is a sequence of p-points  $x^k \to x$  such that  $L(x^k) \to \infty$ .
- x is folding point iff there is a sequence of salient points  $s_{i_k} \to x$ .

**Theorem:** If  $s \in (\sqrt{2}, 2]$  and  $h : \varprojlim([c_2, c_1], T_s) \rightarrow \varprojlim([c_2, c_1], T_s)$  is a homeomorphism, then there is R such that for every folding point x:  $h(x) = \sigma^R(x)$ 

A typical example of this is when T has a periodic critical point of period 3. The below picture illustrates the resulting inverse limit space; p is the fixed point of  $\hat{f}$ .



This representation has a single infinite Wada channel.