Global Optimization in the COCONUT project

Outline of Algorithm
API Design
Inference Engines
Examples

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Project funded by the Future and Emerging Technologies arm of the IST Programme
FET-Open scheme
The COCONUT project

- European Union research and development project
- Partners from six European universities: Nantes, Lausanne, Vienna, Louvain-la-Neuve, Coimbra, Darmstadt and an industrial partner: ILOG
- Aimed at the integration of the existing approaches to continuous global optimization and constraint satisfaction
- December 2000 – November 2003
Basic Modular Setup

Strategy Eng.

Main part of Algorithm, makes decisions

Inference Eng.

Problem structure analysis, local optimization, constraint satisfaction, interval analysis, linear relaxation, convex opt., bisection,...

Reports

Produces files (AMPL), human readable output, checkpointing

Management

Problem management, Resources, Parallelization

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The API is designed to make the development of the various module types independent of each other and independent of the internal model representation.

- A collection of C++ classes.
- Uses FILIB++ and MTL.
- Supports dynamic linking.
Modular API design (cont.d)

- All inference engines are subclasses of one class, so they have the same basic structure.
- Communication with the strategy engine by a database-like communication.
- The API implementation (w/o inference engines) consists of 44000 lines of C++ code and a few perl scripts, organized into 128 files, occupying 1.3 MB of disk space.
Search graph

- Starts at the **original model**
- Contains **relaxations**
- and **splits**.
- It is not a tree since it might also contain **glueings**.
- Some of the nodes will be **terminal**, since they can be solved completely.

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Model Relations

\[
\begin{align*}
\text{Relaxation} & : \min c^T x \\
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x]
\end{align*}
\]

\[
\begin{align*}
\text{Reduction} & : \min c^T x \\
\text{e.g.} & \quad \text{Add cut, prune box} \\
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x]
\end{align*}
\]

\[
\begin{align*}
\text{Split} & : \min c^T x \\
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x'] \subset [x]
\end{align*}
\]

\[
\begin{align*}
\text{s.t.} & \quad d^T x \leq d_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x]
\end{align*}
\]

\[
\begin{align*}
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x'] \subset [x]
\end{align*}
\]

\[
\begin{align*}
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x'] \subset [x]
\end{align*}
\]

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Model Reductions

\[
\begin{align*}
\min & \quad c^T x \\
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x]
\end{align*}
\]

Relaxation

\[
\begin{align*}
\min & \quad c^T x \\
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad x \in [x]
\end{align*}
\]

Split

\[
\begin{align*}
\min & \quad c^T x \\
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x'] \subset [x]
\end{align*}
\]

\[
\begin{align*}
\min & \quad c^T x \\
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x''] \subset [x]
\end{align*}
\]

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Model Glueing

\[
\begin{align*}
\min & \quad c^T x \\
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x]
\end{align*}
\]

Glue, Unsplit

\[
[x'] \cup [x''] = [x]
\]

\[
\begin{align*}
\min & \quad c^T x \\
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x'] \subset [x]
\end{align*}
\]

\[
\begin{align*}
\min & \quad c^T x \\
\text{s.t.} & \quad b^T x \leq b_0 \\
& \quad a(x) \leq 0 \\
& \quad x \in [x''] \subset [x]
\end{align*}
\]

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Models are organized in the **search graph**, represented by a Directed Acyclic Graph (DAG).

For every model in the search graph the following information is stored:

- Every equation/inequality is assigned a number of **annotations** describing its properties (e.g. linear, quadr., convex, separable, redundant, ...).
- Additional **local information** (e.g. local optima, active constraints, Lagrangian multipliers,...) is added.
- A description of the **relation** between the problem and its parent is provided.
Search graph implementation

- The DAG is implemented using the VGTL, a library following the generic programming spirit of the C++ STL.
- There are two types of nodes:
  - **Full nodes** contain complete descriptions of models,
  - **Delta nodes** contain only the changes to the parent model in order to save storage capacity
- The search graph has a **focus** pointing to the model which is worked upon. This model is copied into an enhanced structure – the **work node**. A reference to this work node is passed to the inference engines.
- The graph itself can be analyzed by **search inspectors**.
The internal mathematical representation of a problem is

\[
\begin{align*}
\min & \quad f_{\text{lin}}(x) + f_{\text{quad}}(x) + f_{\text{sep}}(x) + f_0(x) \\
\text{s.t.} & \quad G_{\text{lin}}(x) + G_{\text{quad}}(x) + G_{\text{sep}}(x) + G_0(x) \in S_c \\
& \quad x \in S_v
\end{align*}
\]

where (currently) the sets $S$ are boxes.

The algorithmic representation is in graph form using not a tree (or forest) as usual but a directed acyclic graph.

Variables appearing left of an assignment are substituted out.
Directed Acyclic Graph (DAG)

**Objective**

\[ \min \quad 3x^2z + 4xy^3z \]

**Constraints**

\[ 7x + 3x^2z + 7y^3 \leq 6 \]

\[ y \in [1, 2], z \in [-1, 4] \]

- DAG representation of the model
- similar to computational tree
- every node is an expression
- a node may have more than one parent
- Constants and variables are sources, objective and constraints are sinks

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Expressions

- Every vertex represents a function $F$ mapping a vector $x \in \mathbb{R}^n$ to a value $F(x) \in \mathbb{R}$.
- Predefined functions include \texttt{sum}, \texttt{product}, \texttt{max}, \texttt{min}, elementary real functions (\texttt{exp, log, pow, sqrt, ...})
- Variable indicator contains the indices of the variables this vertex depends on.
- Additional information is added (ranges, semantics, variable name, vertex number, ...)

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Evaluation of a DAG

- Evaluation works similar to computation trees by performing a **graph walk**.
- **Caching** keeps evaluation work minimal.
- The whole model is stored in **one** graph.
- Defining **short-cuts** makes it possible to replace graph walks by evaluation functions. Short-cuts may be defined at every node.
- Additional elementary functions can easily be incorporated.
Graphs and Evaluators

- Generic Graph Library (Vienna Graph Template Library) in C++ to construct and manipulate DAGs, and forests (trees).
- Generic Programming approach with containers, walkers, function objects, and generic algorithms.
- For expression graphs (DAG or tree) special visitors are provided — (cached) forward and backward evaluators.
- Currently implemented Evaluators:
  - Real Function Values and Function Ranges
  - Gradients (Real, Interval)
  - Slopes
- In the near future Evaluators for:
  - Hessians (Real, Interval)
  - Second order Slopes

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Example

- Interval evaluation and constraint propagation produce bounds on each node
- No reduction on the domain of the variables
- The bounds on intermediate nodes are improved compared to interval evaluation
Linear enclosures produced using slopes give redundant constraints, e.g.

\[24(x_1 - 2) - 48(x_2 - 4) - 32(x_3 - 4) \leq 0\]
Example (ctd.)

- Now constraint propagation leads to a reduction of the domain of the variables

\[ x_2 \in [3.4, 4] \]
\[ x_3 \in [3.4, 4] \]

- With previously known techniques but without (expensive) higher order consistency, such a reduction would have required a split of the box.
Corresponding to every type of problem change, a class of inference engines is designed:
- Model analysis (e.g. find convex part)
- Model reduction (e.g. pruning, fathoming)
- Model relaxation (e.g. linear relaxation)
- Model splitting (e.g. bisection)
- Model glueing (e.g. undo excessive splitting)
- Update local information (e.g. probing, local optimization)
- Check certificate (check correctness of calculation)

Inference engines never change the model but calculate which changes may be made and are considered useful.
Inference Engines: General features

- All inference modules only advertise changes.
- There is a fixed documentation structure defined.
  - Services Provided
  - Limits
  - Structure, Prerequisites of Input
  - Structure, Features of Output
  - Control Parameters
  - Termination Reason

- They produce a **result** where every possible change is listed together with a **weight** (the higher the weight the more important the change) and a **certificate** for correctness.

- They collect **statistical data** to support the strategy engine in making decisions.

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Several state of the art techniques were implemented as inference engines:

- STOP (starting point generator)
- DONLP2-INTV (local optimizer)
- Karush-John-Condition Generator
- Point Verifier
- Exclusion Box
- Interval constraint propagation
- Linear Relaxation
- CPLEX (linear programming solver)
- Basic Splitter
- BCS (box covering solver)
Inference Engine: STOP

- **Heuristic** Global Optimization Algorithm
- Combines Multi-Level-Coordinate-Search and Constraint Propagation
- Produces Starting Points for Local Optimization

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FET-Open scheme
General purpose **non-linear local optimizer** for continuous variables
- SQP method
- Dense Linear Algebra
- Envelope uses standard evaluators, gradients are computed by automatic differentiation
Inference Engine:
Karush-John Conditions

- Generates the DAG representation of the Karush-John first order optimality conditions
- Every constraint (even two-sided) gets associated one Lagrange multiplier
- Constructed by symbolic differentiation of the DAG representation
Inference Engine: Point Verifier

- Computes a **uniqueness region** around an approximate solution, in particular a **verified point**
- Uses Karush-John conditions
- Tries to maximize the uniqueness region by inflation

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Derives a large **exclusion box** and a tiny **inclusion box** such that the area between these two boxes does not contain a local optimizer.

They are computed around an approximate local optimizer to **get rid of the cluster effect**.

Does not focus on uniqueness.

Uses slopes and H-matrix techniques.
Inference Engine: Constraint Propagation

- Performs the hull-consistency algorithm for **constraint propagation**.
- Reduces the possible range of the variables
- Might detect infeasibility of the problem
Inference Engine: Linear Relaxation

- Computes a **linear relaxation** of the problem.
- Uses centered forms and slopes to compute the linear inequalities.
- Makes use of the DAG enhancements to improve the slopes.
- Either adds the linear relaxation as cuts or generates a full linear model.
Inference Engine: CPLEX

- Solves linear problems.
- Interfaces the state-of-the-art commercial linear solver CPLEX.
- Extremely good performance.
Inference Engine: Basic Split

- Provides splitting coordinates and split points.
- Computes a difficulty estimate for the variables involved.
- Suggests splits for the \( n \) most difficult variables.
- Uses exclusion box and solution information to improve the choice of split points.
- Cuts exclusion boxes out of the search area by careful choice of splitting coordinates.

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Inference Engine: BCS

- Covers the feasible area by boxes.
- Uses DMBC (dichotomous maintaining bound-consistency) and UCA6 (union-conservative approximation) in both basic and enhanced variants.
- Distinguishes between boxes in the interior and at the border of the feasible region.
- Uses the commercial ILOG Solver, or the constraint propagator provided by IRIN, but can work with any constraint propagator.
Contributions from the outside of the COCONUT project

We are happy that researchers and companies from outside the COCONUT project agreed to complement our efforts in integrating the known techniques:

- Bernstein modules by J. Garloff & A. Smith (U. Konstanz)
- Verified lower bounds for convex relaxations by Ch. Jansson (TU Hamburg-Harburg)
- GAMS reader by the GAMS consortium
- Taylor arithmetic by G. Corliss (Marquette U.)
- Asymptotic arithmetic by K. Petras (U. Braunschweig)
- XPRESS commercial LP-solver (Dash Optimization)
- Hopefully additional contributions by you!

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Worst case finite element analysis

- Linear FEM equations become non-convex if material data is uncertain.
- Typical size of uncertainty is 10-20% in elasticity and cross-section area.
- Law requires the computation of the worst case.
- Industry relevant problems have some thousand variables.

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Worst case FEM structural analysis
Promising result

- Worst case analysis on the displacements $u$ for a 20x20 wall in the non-linear system
  $$A(x)u = b$$
- 1620 material parameters $x$ with 16.4% uncertainty, 840 displacements $u$
- Traditional methods fail for 0.01% uncertainty
- Exploiting the special structure, within 30s on a 1.6 Ghz Pentium 4, without bisection we get

<table>
<thead>
<tr>
<th>Uncertainty (%)</th>
<th>0.01</th>
<th>0.05</th>
<th>0.5</th>
<th>1</th>
<th>2.5</th>
<th>5</th>
<th>10</th>
<th>16.4</th>
</tr>
</thead>
<tbody>
<tr>
<td>Overestimation</td>
<td>1.03</td>
<td>1.15</td>
<td>2.55</td>
<td>4.12</td>
<td>8.92</td>
<td>17.26</td>
<td>35.33</td>
<td>61.59</td>
</tr>
</tbody>
</table>

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This setup allows for highest flexibility and extensibility
- the modules are split into inference engines (calculation) and management parts
- additional modules for model handling are added
- The strategy engine decides which components are called in every algorithmic step of this type
This class of modules produces **output**. Various types of files and human readable output will have to be created.

**Examples:**
- Solution Report (humans, AMPL, GAMS)
- Progress Information
- Checkpointing
- Debugging Information
- Error Messages
Management Modules

- Corresponding to every internal part of the program, a class of **management modules** is designed:
  - Model management
  - Data collection
  - Resource management
  - Initialization management

- Management modules **never calculate** anything. They just **perform some** of the changes which have been advertised by inference modules.
Strategy Engine

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Strategy Engine (ctd.)

- It is the core of the algorithm and consists of
  - The logic core ("search") which is essentially the main solution loop,
  - Special decision makers (very specialized inference engines) for
determining the next action at every point in the algorithm.
- It calls the management modules, the report modules, and
the inference engines in succession.
- It can be programmed using a simple strategy language
(interpreted, Python based).
  - (Semi-)interactive and automatic solution process
  - Debugging and single-stepping of strategies
  - Object oriented, dynamically typed objects, garbage collected
  - Easily extendable

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Strategy Engine (ctd.)

- Manages the search graph via the **search graph manager**, 
- Manages the search database via the **database manager**, 
- Uses a **component framework** to communicate with the inference engines, 
- Launches inference engines dynamically (on need) to avoid memory overload, 
- Provides a management interface, 
- Strategy engine is itself a component, so multilevel strategies are possible, 
- Prepared for **distributed and parallel computing**, and **distributed memory**

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Extensibility

- The strategy language makes it easy to change the strategy.
- A registration phase during initialization removes the need to recompile the program when new inference engines are added.
- Registration also allows us to balance scientific and commercial interests:
  - Free but reduced core version with open API specification
  - Free strategy engine with basic strategy
  - Advanced commercial components
- Extending the system by external contributors is made easy by this modular design.
We hope that the community will contribute to this algorithmic framework.
Thank you for your attention!

COCONUT Website:
http://www.mat.univie.ac.at/coconut

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