

# *Apparent Singularities of Differential Systems with Rational Function Coefficients.*

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# Introduction - Preliminaries

# Apparent Singularities

$\mathbf{K} = \mathbb{C}(z)$ ,  $' = \frac{d}{dz}$ . For  $A \in \text{Mat}_n(\mathbf{K})$ , we denote by  $[A]$  the system:

$$[A] \quad \frac{dX}{dz} = A(z)X,$$

- The (finite) **singularities** of system  $[A]$  are the poles of the entries of  $A(z)$ .
- **Singularities of solutions of the system  $[A]$  are among the singularities of  $[A]$ , but the converse is not always true.**

**Def.** An **apparent singularity** of  $[A]$  is a singular point where the general solution of  $[A]$  is holomorphic.

## Question:

How to detect and remove the apparent singularities of a given system  $[A]$ ?

## Example

$$[A] \quad \frac{dX}{dz} = A(z)X, \quad A(z) = \begin{bmatrix} 0 & 1 \\ \frac{-2}{z} & 1 + \frac{2}{z} \end{bmatrix}.$$

- A fundamental matrix solution of  $[A]$  is

$$\begin{bmatrix} e^z & 1 + z + \frac{z^2}{2} \\ e^z & 1 + z \end{bmatrix}.$$

- Hence  $z = 0$  is an apparent singularity for  $[A]$ .
- The **polynomial gauge transformation**

$$X = T(z) Y, \quad T(z) = \begin{bmatrix} 1 & 0 \\ 1 & z^2 \end{bmatrix}$$

takes  $[A]$  into the equivalent system

$$[B] \quad \frac{d}{dz} Y = B Y$$

where

$$B = T[A] := T^{-1} \left( AT - \frac{dT}{dz} \right) = \begin{bmatrix} 1 & z^2 \\ 0 & 0 \end{bmatrix}.$$

## What this talk is about?

- A **general fact** : Any system  $[A]$  with rational coefficients can be reduced to a *gauge equivalent* system  $[B]$  with rational coefficients, such that the finite singularities of  $[B]$  coincide with the non-apparent singularities of  $[A]$ .
- Such a system  $[B]$  will be called a *complete desingularization* of  $[A]$
- We present an algorithm which outputs a desingularization for any input system  $[A]$ .
- More generally, given  $[A]$ , to find a **polynomial** matrix  $T \in \text{GL}_n(\mathcal{C}(z))$  such that  $B = T[A]$  satisfies  $\text{den}(B) \mid \text{den}(A)$ , and  $\text{den}(B)$  is as “small” as possible.

## *Previous and related works*

- Desingularization of linear difference/ differential (and more generally Ore) operators, e.g.
  - ▶ Abramov, van Hoeij 1999
  - ▶ Tsai 2000
  - ▶ Abramov-B.-van Hoeij'2006,
  - ▶ Chen-Jaroscsek-Kauers-Singer'2013, Chen-Kauers-Singer'2015
  - ▶ Yi Zhang, ISSAC'2016
- Desingularization of linear differential systems:
  - ▶ B.'2010,
  - ▶ B.-Maddah ISSAC'2015
- Desingularization of linear difference systems [Maximilian Jaroscsek' talk](#)

## Classification of Singularities

Consider a System of first order linear differential equations:

$$[A] \quad \frac{dX}{dz} = A(z)X, \quad A(z) \in \text{Mat}_n(\mathbb{C}(z))$$

- A pole  $z_0$  of  $A(z)$  is a **regular singular point** for  $[A]$  if there is a fundamental solution  $W$  of  $[A]$  has the form:

$$W(z) = \Phi(z)(z - z_0)^\Lambda$$

where  $\Phi(z)$  is holomorphic and  $\Lambda$  is a constant matrix.

- Otherwise  $z_0$  is called an **irregular singular point**.
- A system  $[A]$  has **regular singularity at  $z_0$**  iff it is gauge equivalent to a system  $[B]$  with a simple pole at  $z_0$ .
- We shall refer to simple poles of  $A(z)$  as **simple singularities** of  $[A]$ .
- if  $z_0$  is an apparent singularity then  $z_0$  is a regular singularity and thus can be reduced to a simple one.

## Apparent singularities are removable

**Prop.0** If  $z = z_0$  is a finite apparent singularity of  $[A]$  then there exists a polynomial matrix  $T(z)$  with

$$\det T(z) = c(z - z_0)^\alpha, \quad c \in \mathbb{C}^*, \alpha \in \mathbb{N}$$

such that  $[B] := T[A]$  has no pole at  $z = z_0$ .

### Proof.

- Every fundamental solution  $F$  of  $[A]$  is holomorphic (in a neighborhood of  $z_0$ );
- There exists matrices  $P(z) \in GL_n(\mathbb{C}[z])$ , and  $Q(z) \in GL_n(\mathbb{C}[[z - z_0]])$  such that

$$P(z)F(z)Q(z) = \text{Diag}((z - z_0)^{\alpha_1}, \dots, (z - z_0)^{\alpha_n})$$

where  $\alpha_1, \dots, \alpha_n \in \mathbb{N}$

- Take

$$T(z) = P^{-1}(z) \text{Diag}((z - z_0)^{\alpha_1}, \dots, (z - z_0)^{\alpha_n})$$



# Characterization of Regular Singularities

## How to recognize regular singularities?

**Problem 1:** Given a system  $[A]$  and a pole  $z_0$  of order  $p_{z_0}(A) > 1$  to decide whether  $z_0$  is regular singular or not.

In other words, to decide if the order of a given singularity can be reduced to 1 or not?

▶ There is no analogue of the Fuchs' Criterion.

**Problem 2:** Given a system  $[A]$  and a pole  $z_0$  of order  $p_{z_0}(A) > 1$ , to decide whether there exists  $T \in \text{GL}(n, \mathbb{C}((z - z_0)))$  such that  $p_{z_0}(T[A]) < p_{z_0}(A)$ .

▶ There is a method due to Moser (1960) which solve these two problems.

▶ **Rational Moser -Algorithm** Barkatou'1995: It transforms a given system over  $\mathbb{C}(z)$  into an equivalent one for which the orders of the finite poles are reduced to their minimal values.

▶ Other methods for reducing the rank of a singularity (to its minimal value) do exist: Levelt (1992), Wagenfurer (1989), ..., B.-Pfluegel (2007, 2008), B.-El Bacha (2012).

## An example

$$A = \begin{bmatrix} \frac{-1}{z} & \frac{-1+z^{17}-8z^{14}+24z^{11}-32z^8+16z^5}{z^3(z^3-2)^2} & 0 & \frac{1}{z^3(z^3-2)^2} \\ 0 & 0 & \frac{1}{z^3(z^3-2)^2} & 0 \\ \frac{1}{(z^3-2)^3 z^4} & 0 & \frac{4}{z} & 0 \\ 0 & -\frac{2}{z} & \frac{1}{z^3(z^3-2)^2} & \frac{2}{z} \end{bmatrix} \in M_4(\mathbb{Q}(z))$$

► When applied to  $[A]$  and the roots of the irreducible polynomial  $z^3 - 2$ , Algorithm [Bar'95] produces the equivalent matrix

$$B = T[A] = \begin{bmatrix} -\frac{10z^3+4}{z(z^3-2)} & 0 & 0 & 0 \\ \frac{1}{z^3(z^3-2)} & -\frac{8z^3+2}{z(z^3-2)} & 0 & \frac{z^2}{z^3-2} \\ 0 & \frac{1}{z^4(z^3-2)} & \frac{z^3-8}{z(z^3-2)} & 0 \\ 0 & 0 & \frac{1}{z^3(z^3-2)} & 0 \end{bmatrix}$$

and the gauge transformation

$$T = \begin{bmatrix} 0 & (z^3 - 2)^3 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & z^3 - 2 & 0 \\ (z^3 - 2)^4 & 0 & 0 & 1 \end{bmatrix}$$

- The denominator of the matrix  $B$  is  $z^4(z^3 - 2)$ .
- Hence the system  $[A]$  has regular singularities at the zeros of  $q = z^3 - 2$ .

Applied to  $[B]$  and  $z_0 = 0$ , Algorithm [Bar'95] produces the equivalent matrix

$$C = S[B] = \begin{bmatrix} 0 & 0 & \frac{1}{z(z^3-2)} & \frac{15z^3-6}{z(z^3-2)} \\ \frac{1}{z^2(z^3-2)} & -\frac{12z^3-6}{z(z^3-2)} & 0 & 0 \\ 0 & \frac{1}{z^2(z^3-2)} & -\frac{4+z^3}{z(z^3-2)} & 0 \\ 0 & 0 & 0 & -\frac{15z^3-6}{z(z^3-2)} \end{bmatrix}$$

and the transformation

$$S = \begin{bmatrix} 0 & 0 & 0 & -z^5 \\ 0 & z^4 & 0 & 0 \\ 0 & 0 & z^2 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix}$$

►  $z = 0$  is still a pole of order  $> 1$ , hence the point  $z = 0$  is an irregular singular point of the original system  $[A]$ .

## Moser's Method

► We assume here  $z_0 = 0$

$$A(z) = \frac{1}{z^p} \sum_{k=0}^{\infty} A_k z^k, \quad A_k \in M_n(\mathbb{C}), \quad A_0 \neq 0.$$

► A necessary condition that there exist a gauge transformation  $T \in GL(n, \mathbb{C}((z)))$  such that  $T[A] = \frac{1}{x^{p'+1}}(B_0 + B_1 z + \dots)$  with  $p' < p$  ( $B_0 \neq 0$ ), is that  $A_0$  is nilpotent.

**Moser rank:**  $m(A) = p - 1 + \frac{\text{rank}(A_0)}{n}$  if  $p > 1$ , otherwise  $m(A) = 1$ .

**Moser invariant:**  $\mu(A) = \min \{m(T[A]) \mid T \in GL(n, \mathbb{C}((z)))\}$

**Definition.**  $[A]$  is said to be Moser-reducible if  $m(A) > \mu(A)$ .

- $[A]$  is Moser-reducible  $\iff \exists T \in GL(n, \mathbb{C}((z)))$  such that  $m(T[A]) < m(A)$ .
- $z = 0$  is regular singular for  $[A] \iff \mu(A) = 1$ .

# A Criterion for Moser-reducibility

**Theorem.** [Moser 1960]

- 1 If  $p > 1$  then  $A$  is Moser-reducible iff the polynomial

$$\mathcal{B}_A(\lambda) := z^{\text{rank}(A_0)} \det(\lambda I - A_0/z - A_1)|_{z=0} \equiv 0.$$

- 2 If  $A$  is Moser-reducible then the reduction can be carried out with a transformation of the form  
 $T = (P_0 + zP_1)\text{diag}(1, \dots, 1, z, \dots, z)$ ,  $P_i \in \mathbb{C}^{n \times n}$ ,  $\det P_0 \neq 0$ .
- Applying Moser's Theorem several times, if necessary,  $\mu(A)$  can be determined.
- Further, a polynomial matrix  $T$  such that  $m(T[A]) = \mu(A)$  can be computed in this way

## *Review of Moser-reduction Algorithms*

- There are various algorithms to compute  $T$  such that  $T[A]$  is Moser-reduced.
- Moser's paper: no constructive algorithm given.
- Dietrich (1978), Hilali/Wazner (1987): first efficient algorithms,
- Bar'1995: version for rational function coefficients, implemented in ISOLDE
- B.-Pflügel (2007): New reduction algorithm + complexity analysis.



## Description of Moser Algorithm

- By a constant gauge transformation we can put  $A_0$  in the form:

$$A_0 = \begin{pmatrix} A_0^{11} & 0 \\ A_0^{21} & 0 \end{pmatrix}, \quad A_0^{11} \in \mathbb{C}^{r \times r} \quad r = \text{rank}(A_0).$$

- Let  $A_1$  be partitioned so that  $A_1^{11}$  is a square matrix of order  $r$ :

$$A_1 = \begin{pmatrix} A_1^{11} & A_1^{12} \\ A_1^{21} & A_1^{22} \end{pmatrix},$$

- Consider

$$G_\lambda(A) = \begin{pmatrix} A_0^{11} & A_1^{12} \\ A_0^{21} & A_1^{22} + \lambda I_{n-r} \end{pmatrix}.$$

- Then  $\det G_\lambda(A) = \mathcal{B}_A(\lambda)$ .
- $A$  is Moser-reducible  $\iff \det G_\lambda(A) \equiv 0$ .

Case 1:  $\text{rank}(A_0^{11} \ A_1^{12}) < r \quad (I)$

A is Moser-reducible  $\iff \begin{vmatrix} A_0^{11} & A_1^{12} \\ A_0^{21} & A_1^{22} + \lambda I_{n-r} \end{vmatrix} = 0.$

**Proposition 1** If  $m(A) > 1$  and  $\text{rank}(A_0^{11} \ A_1^{12}) < r$ , then A is M-reducible and the reduction can be carried out with the gauge transformation

$$T = \text{diag}(zI_r, I_{n-r}).$$

**Proof:** Let  $B = T[A] = T^{-1}AT - T^{-1}\frac{dT}{dz}.$

$$B = z^{-p}[B_0 + zB_1 + \dots] + z^{-1}\text{diag}(I_r, 0)$$

where

$$B_0 = \begin{pmatrix} A_0^{11} & A_1^{12} \\ 0 & 0 \end{pmatrix},$$

Since  $p > 1$ , then  $m(B) = p - 1 + \text{rank}(B_0)/n < m(A) = p - 1 + r/n.$

*Case 1:  $\text{rank}(A_0^{11} \ A_1^{12}) < r$  An example*

$$A = z^{-2} \begin{pmatrix} 0 & 0 \\ 2 & 0 \end{pmatrix} + z^{-1} \begin{pmatrix} 4 & 0 \\ 0 & -3 \end{pmatrix} + \begin{pmatrix} 0 & -4 \\ 0 & 0 \end{pmatrix}.$$

- ▶ Here  $p = 2, r = 1 \Rightarrow m(A) = 1 + 1/2 = 3/2 > 1$ .
- ▶  $\det G_\lambda(A) = \begin{vmatrix} 0 & 0 \\ 2 & -3 + \lambda \end{vmatrix} = 0 \Rightarrow A$  is Moser-reducible.
- ▶ Take

$$T = \begin{pmatrix} z & 0 \\ 0 & 1 \end{pmatrix}$$

$$B := T[A] = T^{-1}AT - T^{-1}T' = \frac{1}{z} \begin{pmatrix} 3 & -4 \\ 2 & -3 \end{pmatrix}.$$

- ▶ The system  $Z' = BZ$  has a singularity of first kind at  $z = 0$ .
- ▶ Hence  $Y' = AY$  has a regular singularity at  $z = 0$ .

Case 2:  $\text{rank}(A_0^{11} \ A_1^{12}) = r \quad (I)$

**Proposition 2** If  $A$  is M-reducible and  $\text{rank}(A_0^{11} \ A_1^{12}) = r$ , then there exists a constant matrix  $Q$  such that the matrix  $G_\lambda(Q[A])$  has the form has the following particular form:

$$G_\lambda(A) = \begin{pmatrix} A_0^{11} & U_1 & U_2 \\ V_1 & W_1 + \lambda I_{n-r-h} & W_2 \\ 0 & 0 & W_3 + \lambda I_h \end{pmatrix}, \quad (1)$$

where  $1 \leq h \leq n - r$ ,  $W_1$ ,  $W_3$  are square matrices of order  $(n - r - h)$  and  $h$  respectively,  $W_3$  is upper triangular with zero diagonal with the condition

$$\text{rank}(A_0^{11} \ U_1) < r \quad (2)$$

Case 2:  $\text{rank}(A_0^{11} \ A_1^{12}) = r \quad (II)$

**Proposition 3** If  $m(A) > 1$  and  $G_\lambda(A)$  has the form (1) with the condition (2), then  $A$  is reducible and the reduction can be carried out with the transformation

$$T = \text{diag}(zI_r, I_{n-r-h}, zI_h)$$

**Proof:** Put  $B = T[A] = T^{-1}AT - T^{-1}\frac{dT}{dz}$ . One has

$$B = z^{-p}[B_0 + zB_1 + \dots] + z^{-1}\text{diag}(I_r, 0, I_h)$$

where

$$B_0 = \begin{pmatrix} A_0^{11} & U_1 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{pmatrix},$$

and then  $\text{rank}(B_0) = \text{rank}(A_0^{11} \ U_1) < r = \text{rank}(A_0)$ . On the other hand since  $p > 1$ , then  $m(B) = p - 1 + \text{rank}(B_0)/n$ . Hence  $m(B) < m(A)$ .

## Moser-reduction (End)

If  $A$  is Moser-reducible and  $m(A) > 1$  then one can construct a matrix polynomial  $S$  of the form :

$$S = U \text{diag}(z, z, \dots, z, 1, 1, \dots, 1)$$

where  $U \in \text{GL}(n, \mathbb{C})$ , such that  $m(S[A]) < m(A)$ .

- Moser's Theorem allows us to check whether  $A$  is Moser-reducible.
- If  $A$  is Moser-reducible then by the above theorem we can find a matrix  $S$  such that  $m(S[A]) < m(A)$ .
- After this reduction has been carried out we can apply Moser's Theorem to check whether further reduction is possible and so on.
- After **at most  $n(p - 1)$**  steps we obtain an equivalent matrix  $B$  such that  $m(B) = \mu(A)$ .
- **The nature of the singularity depends on the first  $n(p - 1)$  coefficients in the series expansion of  $A$**

# Removal of Apparent Singularities

## How to detect and remove an apparent singularity? (I)

**Prop.1:** If  $z = z_0$  is a finite apparent singularity of  $[A]$  then one can construct a polynomial matrix  $T(z)$  with  $\det T(z) = c(z - z_0)^\alpha$ ,  $c \in \mathbb{C}^*$  and  $\alpha \in \mathbb{N}$  such that  $T[A]$  has at worst a simple pole at  $z = z_0$ .

This follows from the fact that:

- if  $z_0$  is an apparent singularity then  $z_0$  is a regular singularity,
  - and that a system with a regular singularity at  $z_0$  is equivalent to a system with a simple pole at  $z_0$ .
- The transformation  $T$  can be constructed using the *rational Moser algorithm* (developed in Bar'1995).



## How to detect and remove an apparent singularity? (II)

**Prop.2** Suppose that  $A(z)$  has simple pole at  $z = z_0$  and let

$$A(z) = \frac{A_0}{(z - z_0)} + \sum_{k \geq 1} A_k (z - z_0)^{k-1}, \quad A_k \in \mathbb{C}^{n \times n}.$$

If  $z_0$  is an apparent singularity then the eigenvalues of  $A_0$  are nonnegative integers and  $A_0$  is diagonalizable.

This follows from the fact that:

- A system having a simple singularity at  $z = z_0$  with residue matrix  $A_0$  possesses a local fundamental solution of the form:

$$\Phi(z)(z - z_0)^\Lambda$$

where  $\Phi(z)$  is holomorphic at  $z = z_0$  and  $\Lambda$  is a constant matrix with

$$\text{spec}(\Lambda) \subset \text{spec}(A_0) - \mathbb{N}$$

- When  $A_0$  is not diagonalizable, the local solutions at  $z_0$  involve logarithmic terms.

## How to detect and remove an apparent singularity? (III)

**Prop.3:** Suppose that  $z = z_0$  is a simple pole of  $A(z)$  and that its residue matrix  $A_0$  has only nonnegative integer eigenvalues. Then one can construct a polynomial matrix  $T(z)$  with

$$\det T(z) = c(z - z_0)^\alpha$$

for some  $c \in \mathbb{C}^*$  and  $\alpha \in \mathbb{N}$  such that

$$B := T[A] = B_0(z - z_0)^{-1} + \dots$$

has at worst a simple pole at  $z = z_0$  with

$$B_0 = mI_n + N$$

where  $m \in \mathbb{N}$  and  $N$  nilpotent.

- Moreover  $z_0$  is an apparent singularity iff  $N = 0$ .
- In this case the gauge transformation  $Y = (z - z_0)^m \tilde{Y}$  leads to a system for which  $z = z_0$  is an ordinary point.

## How to detect and remove an apparent singularity? (IV)

### Skech of the proof of Prop. 3:

- The eigenvalues of  $A_0$  of are nonnegative integers:

$$m_1 > m_2 > \dots > m_s, \quad m_i - m_{i+1} = \ell_i \in \mathbb{N}^*, \quad i = 1, \dots, s-1.$$

- By applying a constant gauge transformation we can assume that:

$$A_0 = \begin{pmatrix} A_0^{11} & 0 \\ 0 & A_0^{22} \end{pmatrix},$$

where  $A_0^{11}$  is an  $\nu_1$  by  $\nu_1$  matrix having one single eigenvalue  $m_1$

$$A_0^{11} = m_1 I_{\nu_1} + N_1$$

$N_1$  being a nilpotent matrix.

- Apply the gauge transformation  $U = \text{diag}((z - z_0)I_{\nu_1}, I_{n-\nu_1})$  yields the new system:

$$Z' = (z - z_0)^{-1} \tilde{A}(z)Z, \quad \tilde{A}(z) = (z - z_0)U^{-1}A(z)U - (z - z_0)U^{-1}U'$$

with the leading matrix:

$$\tilde{A}(z_0) = (A_0 + (z - z_0)U^{-1}A_1U - (z - z_0)U^{-1}U')|_{z=z_0}.$$

- Let  $A_1$  be partitioned as  $A_0$  :

$$A_1 = \begin{pmatrix} A_1^{11} & A_1^{12} \\ A_1^{21} & A_1^{22} \end{pmatrix}, \quad A_1^{11} \in \mathbb{C}^{\nu_1 \times \nu_1}$$

Then

$$\tilde{A}(z_0) = \begin{pmatrix} A_0^{11} - I_{\nu_1} & A_1^{12} \\ 0 & A_0^{22} \end{pmatrix}.$$

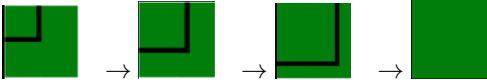
Hence the eigenvalues of  $\tilde{A}(0)$  are:  $m_1 - 1, m_2, \dots, m_s$ , each with the same initial multiplicity  $\nu_j$ .

## How to detect and remove an apparent singularity? (V)

- ▶ By repeating this process  $\ell_1$  times, the eigenvalues become:

$$m_1 - \ell_1 = m_2, m_2, \dots, m_s.$$

- ▶ Thus, after  $\ell_1 + \dots + \ell_{s-1}$  steps, the eigenvalues  $m_1, \dots, m_s$  are reduced to one single eigenvalue  $m_s$  of multiplicity  $\nu_1 + \dots + \nu_s = n$ .


$$A_0 = \begin{matrix} \blacksquare & & & \\ & \blacksquare & & \\ & & \blacksquare & \\ & & & \blacksquare \end{matrix} \rightarrow \begin{matrix} \blacksquare & & & \\ & \blacksquare & & \\ & & \blacksquare & \\ & & & \blacksquare \end{matrix} \rightarrow \begin{matrix} \blacksquare & & & \\ & \blacksquare & & \\ & & \blacksquare & \\ & & & \blacksquare \end{matrix} \rightarrow \blacksquare = m_s I_n + N$$

- ▶  $z_0$  is an apparent singularity iff  $N = 0$ .
- ▶ In this case the gauge transformation  $Y = (z - z_0)^{m_s} \tilde{Y}$  leads to a system for which  $z = z_0$  is an ordinary point.

# Complete Desingularization

- The matrix  $T$  in Prop3 is obtained as a product of invertible constant matrices or diagonal matrices of the form  $U = \text{diag}((z - z_0)I_\nu, I_{n-\nu})$ .

Hence  $T$  is a polynomial matrix with  $\det T(z) = c(z - z_0)^\alpha$  for some  $c \in \mathbb{C}$  and  $\alpha \in \mathbb{N}$ .

- Due to the form of its determinant, the gauge transformation  $T(z)$  in the above proposition **does not affect the other finite singularities** of  $[A]$ .

**Theorem** One can construct a polynomial matrix  $T(z)$  which is invertible in  $\mathbb{C}(z)$  such that the finite poles of  $B := T[A]$  are exactly the true singularities for  $[A]$ .

# Algorithm of Desingularization

**Step 1** Reduce the rank of each singularity to its minimal value :

Compute a polynomial matrix  $T(z)$  such that

- ▶ the zeros of  $\det T(z)$  are in  $\mathcal{P}(A)$
- ▶  $T[A]$  has the same poles as  $A$  with minimal orders.

**Step 2** Select simple singularities with nonnegative exponents:

For each simple pole  $z_0$  compute  $A_{0,z_0} := \text{res}_{z=z_0} A(z)$  and its spectrum.

$\mathcal{App}(A) := \{z_0 \text{ simple singularities such that } \text{spec}(A_{0,z_0}) \subset \mathbb{N}\}$

**Step 3** Make all exponents equal:

For  $z_0 \in \mathcal{App}(A)$  compute a polynomial matrix  $T_{z_0}(z)$  with  $\det T_{z_0}(z) = c(z - z_0)^\alpha$  such that  $T_{z_0}[A]$  has at worst a simple pole at  $z = z_0$  with residue matrix of the form  $R_{z_0} = m_{z_0} I_n + N_{z_0}$  where  $m_{z_0} \in \mathbb{N}$  and  $N_{z_0}$  nilpotent.

**Step 4** Determine the apparent singularities of  $[A]$ :

Keep in  $\mathcal{App}(A)$  only the points  $z_0$  for which  $N_{z_0} = 0$ .

**Step 5** Shift the exponent to 0 :

Apply the scalar transformation  $T = \prod_{z_0 \in \mathcal{App}(A)} (z - z_0)^{m_{z_0}} I_n$ .

# Application to Desingularization of Scalar Differential Equations



# Desingularization of Scalar differential Equations

- **Removing** apparent singularities of  $L \in \mathbb{C}(z)[\partial]$ :

→ to **construct another operator**  $\tilde{L} \in \mathbb{C}(z)[\partial]$  such that:

- any solution of  $L(y) = 0$  is a solution of  $\tilde{L}(y) = 0$ ,**  
i.e.  $\tilde{L} = R \circ L$  for some  $R \in \mathbb{C}(z)[\partial]$
  - and the singularities of  $\tilde{L}$  are exactly the singularities of  $L$  that are not apparent.**
- Such an operator  $\tilde{L}$  is called a **desingularization** of  $L$ .

**Example:**  $L = \partial - \frac{\mu}{z}$ ,  $\mu \in \mathbb{N}$ .

The operator  $\tilde{L} = \partial^{\mu+1}$  is a desingularization of  $L$ .

- Several algorithms have been developed for linear differential (and more generally Ore) operators, e.g.
  - Abramov-B.-van Hoeij'2006,
  - Chen-Jaroschek-Kauers-Singer'2013, Chen-Kauers-Singer'2015
  - We developed, in [ABH 2006]<sup>1</sup> an algorithm that, given an operator  $L$  of order  $n$ , produces a **desingularization  $\tilde{L}$  with minimal order  $m \geq n + 1$** .
  - This algorithm has been implemented in Maple.
  - I will refer to this algorithm as **ABH method**.

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<sup>1</sup> S. ABRAMOV, M. BARKATOU and M. van HOEIJ AAEECC 2006

## Example 1

Consider the second order operator

$$L := \partial^2 - \frac{(z+2)}{z}\partial + \frac{2}{z}.$$

- $z = 0$  is a singularity of  $L$ .
- The general solution of  $L(y) = 0$  is given by

$$c_1 e^z + c_2 \left(1 + z + \frac{z^2}{2}\right) \quad c_1, c_2 \in \mathbb{C}.$$

- $L$  has an apparent singularity at  $z = 0$ .
- The desingularization computed by ABH method is **of order 4**

$$\tilde{L} = \partial^4 + \left(-1 + \frac{z}{4}\right)\partial^3 + \left(-\frac{1}{4} - \frac{3z}{8}\right)\partial^2 + \left(\frac{1}{2} + \frac{z}{8}\right)\partial - \frac{1}{4}$$

- The apparent singularity of  $L$  at  $z = 0$  can be removed by computing a *gauge equivalent* first-order differential system with coefficient in  $\mathbb{C}(z)$  of size  $\text{ord}(L) = 2$ .
- Consider the first-order differential system associated with  $L$

$$[A] \quad \frac{d}{dz} X = A(z)X, \quad A(z) = \begin{bmatrix} 0 & 1 \\ \frac{-2}{z} & 1 + \frac{2}{z} \end{bmatrix}.$$

- Set

$$X = T(z) Y, \quad \text{where} \quad T(z) = \begin{bmatrix} 1 & 0 \\ 1 & z^2 \end{bmatrix}.$$

- The new variable  $Y$  satisfies the *gauge equivalent* first-order differential system of the same dimension given by

$$[B] \quad \frac{d}{dz} Y = B Y$$

where

$$B := T^{-1}AT - T^{-1} \frac{d}{dz} T = \begin{bmatrix} 1 & z^2 \\ 0 & 0 \end{bmatrix}.$$

## Example 2

- Consider

$$L = \partial^2 + \frac{(3z^2 - 4)}{z(z^2 + 2)} \partial - 2 \frac{-1 + 2z^2}{z^2 + 2}$$

- $L$  has an apparent singularity at  $z = 0$  with local exponents 0 and 3.
- The desingularization computed by ABH method is of order 4

$$\begin{aligned} \tilde{L} = & \partial^4 + 1/2 \frac{z(24 + 7z^2)}{z^2 + 2} \partial^3 + 1/2 \frac{(58z^2 + 88 + 27z^4)}{(z^2 + 2)^2} \partial^2 \\ & - 1/2 \frac{z(-4z^2 + 4 + 93z^4 + 28z^6)}{(z^2 + 2)^3} \partial - 4 \frac{44z^2 + 16 + 42z^4 + 7z^6}{(z^2 + 2)^3}. \end{aligned}$$

- The companion matrix of  $L$  is

$$A = \begin{bmatrix} 0 & 1 \\ 2 \frac{-1+2z^2}{z^2+2} & -\frac{3z^2-4}{z(z^2+2)} \end{bmatrix}$$

- It has a simple pole at  $z = 0$  with a residue matrix  $A_0 = \begin{bmatrix} 0 & 0 \\ 0 & 2 \end{bmatrix}$ .
- Our algorithm computes the following gauge transformation  $T$

$$T = \begin{bmatrix} 1 & 0 \\ z & -z^2 \end{bmatrix}$$

- The matrix of the new equivalent system is

$$B = T^{-1}(AT - T') = \begin{bmatrix} z & -z^2 \\ 1 & -\frac{z(z^2+7)}{z^2+2} \end{bmatrix}$$

- It has  $z = 0$  as ordinary point.

## Example 3

- Let  $\partial = \frac{d}{dz}$  and consider

$$L = \partial^2 - \frac{(z^2 - 3)(z^2 - 2z + 2)}{(z - 1)(z^2 - 3z + 3)z} \partial + \frac{(z - 2)(2z^2 - 3z + 3)}{(z - 1)(z^2 - 3z + 3)z}.$$

- $L$  has apparent singularities at  $z = 0$  and the roots of  $z^2 - 3z + 3 = 0$ .
- A desingularization computed by the classical algorithm<sup>2</sup> is given by:

$$\begin{aligned} \tilde{L}_{\text{Classical}} &= (z - 1)(z^4 - z^3 + 3z^2 - 6z + 6)\partial^4 \\ &\quad - (z^5 - 2z^4 + z^3 - 12z^2 + 24z - 24)\partial^3 \\ &\quad - (3z^3 + 9z^2)\partial^2 + (6z^2 + 18z)\partial - (6z + 18). \end{aligned}$$

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<sup>2</sup>Exm 1, Chen-Kauers-Singer'14

- A desingularization computed by the probabilistic method of CKS14<sup>3</sup> is given by:

$$\begin{aligned}
 \tilde{L}_{CKS} = & (z - 1)(z^6 - 3z^5 + 3z^4 - z^3 + 6)\partial^4 \\
 & - (2z^6 - 9z^5 + 15z^4 - 11z^3 + 3z^2 - 24)\partial^3 \\
 & - (z^7 - 4z^6 + 6z^5 - 4z^4 + z^3 + 6z - 6)\partial \\
 & + (2z^6 - 9z^5 + 15z^4 - 11z^3 + 3z^2 - 24).
 \end{aligned}$$

- The removal of one apparent singularity introduces new singularities. The latter can then be removed by using a trick introduced in ABH algorithm.

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<sup>3</sup>Exm 7(1), Chen-Kauers-Singer'14



- The desingularization computed by ABH method is:

$$\begin{aligned} \tilde{L}_{ABH} = & \partial^4 + \frac{(16z^4 - 55z^3 + 63z^2 - 42z + 36)}{9(z-1)} \partial^3 \\ & - \frac{(64z^5 - 316z^4 + 591z^3 - 468z^2 + 123z + 42)}{9(z-1)^2} \partial^2 \\ & - \frac{96z^5 - 570z^4 + 1333z^3 - 1597z^2 + 993z - 219}{9(z-1)^3} \\ & + \frac{\beta}{9(z-1)^3} \partial, \end{aligned}$$

where

$$\beta = (48z^6 - 197z^5 + 148z^4 + 488z^3 - 1162z^2 + 999z - 288).$$

- The companion matrix of  $L$  is

$$A = \begin{bmatrix} 0 & 1 \\ \frac{(z-2)(2z^2-3z+3)}{(z-1)(z^2-3z+3)z} & \frac{(z^2-3)(z^2-2z+2)}{(z-1)(z^2-3z+3)z} \end{bmatrix}$$

- Our new algorithm computes the following gauge transformation  $T$

$$T = \begin{bmatrix} 1 & 0 \\ 1 & (-z^2 + 3z - 3)z^2 \end{bmatrix}$$

- The matrix of the new equivalent system is

$$B = T^{-1}(AT - T') = \begin{bmatrix} 1 & -z^2(z^2 - 3z + 3) \\ 0 & \frac{2}{1-z} \end{bmatrix}$$

- It has  $z = 0$  and roots of  $z^2 - 3z + 3 = 0$  as ordinary points.
- No new apparent singularities are introduced.

- The desingularization algorithms developed specifically for scalar equations are based on computing a least common left multiple of the operator in question and an appropriately chosen operator.
- This outputs an equation whose solution space contains strictly the solution space of the input equation.
- The new algorithm is based on an adequate choice of a gauge transformation.
- The desingularized output system is always equivalent to the input system and the dimension of the solution space is preserved.
- The transformations and the equivalent systems computed by our algorithm, have rational function coefficients.