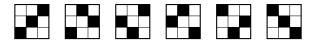
Sorting with C-machines

Jay Pantone Dartmouth College Hanover, NH



Erwin Schrödinger Institute Algorithmic and Enumerative Combinatorics October 20, 2017

Joint work with Michael Albert, Cheyne Homberger, Nathanial Shar, Vince Vatter

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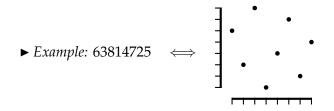
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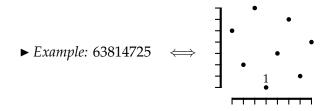
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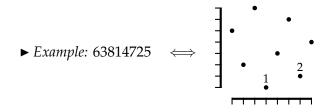
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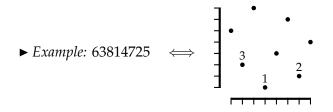
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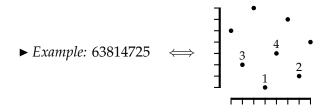
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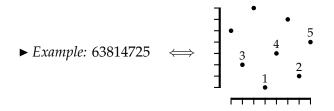
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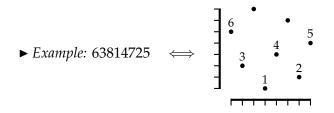
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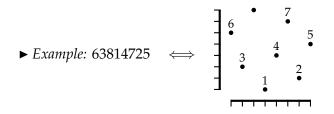
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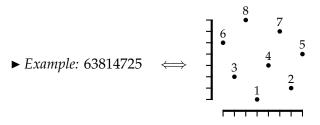
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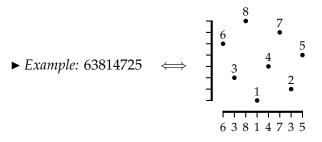
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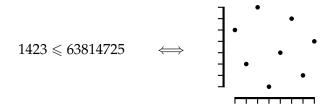
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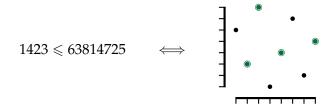


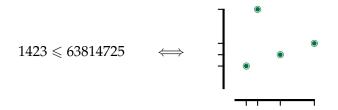
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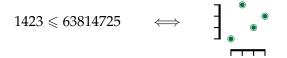
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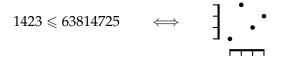












Exotic Machines

What else?

Permutation Poset

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Exotic Machines

What else?



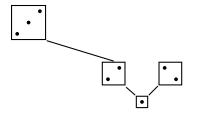
Exotic Machines

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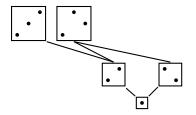
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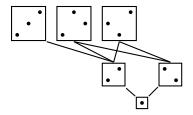
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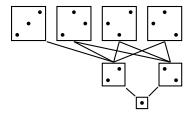
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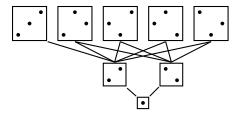
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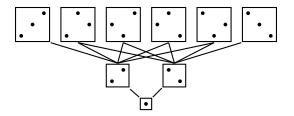
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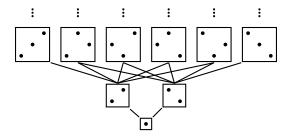
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Permutation Poset

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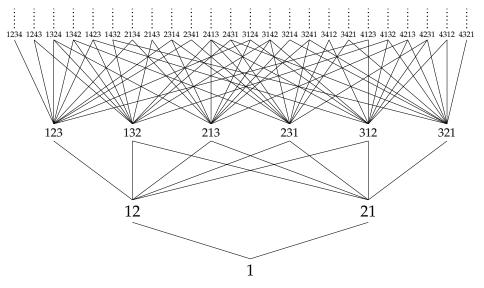
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The class with basis B is denoted Av(B).

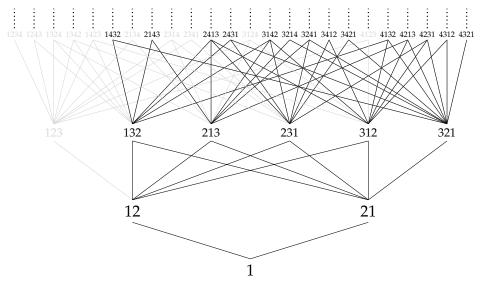
Exotic Machines

What else?



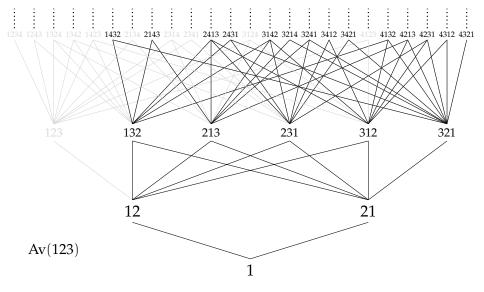
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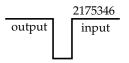
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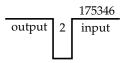


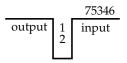
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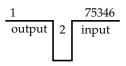
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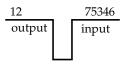


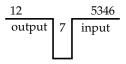


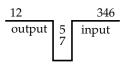


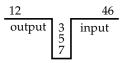


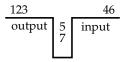


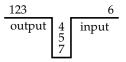


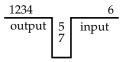


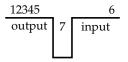


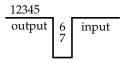


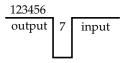


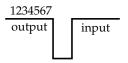




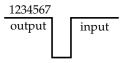




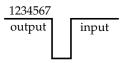


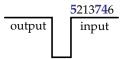


The permutation 2175346 can be sorted by a stack.

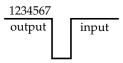


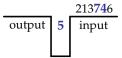
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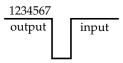


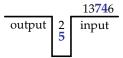
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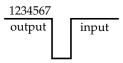


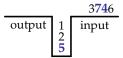
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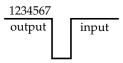


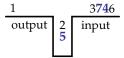
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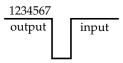


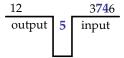
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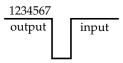


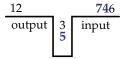
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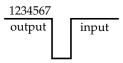


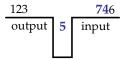
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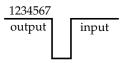


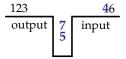
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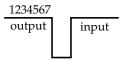


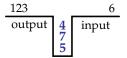
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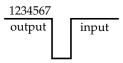


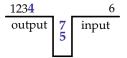
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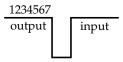


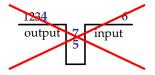
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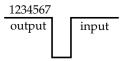


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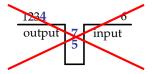




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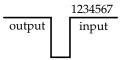


Not all permutations can be sorted by a stack.

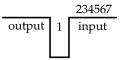


Permutations containing 231 cannot be sorted by a stack. In fact, the permutations that are stack-sortable are exactly those in Av(231).

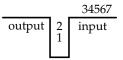
The permutation 2156473 can be generated by a stack.



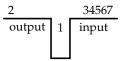
The permutation 2156473 can be generated by a stack.



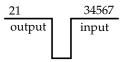
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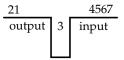
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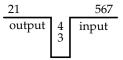
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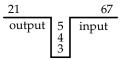
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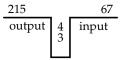
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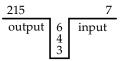
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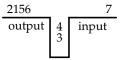
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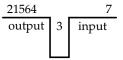
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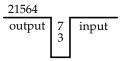
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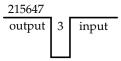
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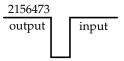
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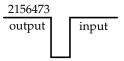
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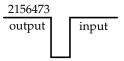
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(Notice the stack always holds a decreasing subpermutation when read top to bottom.)

The permutations that can be *generated* by a stack are exactly the inverses of those that can be *sorted* by a stack.

The permutation 2156473 can be generated by a stack.



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The permutations that can be *generated* by a stack are exactly the inverses of those that can be *sorted* by a stack.

A stack can generate the permutations in $Av(231^{-1}) = Av(312)$.



The stack can be thought of as a container that always holds a decreasing permutation, where at any time we can:

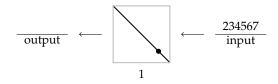
 push a new maximum entry into the container such that the container holds a decreasing permutation



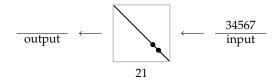
- push a new maximum entry into the container such that the container holds a decreasing permutation
- ► pop the leftmost entry out of the container.



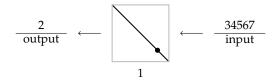
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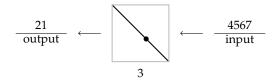
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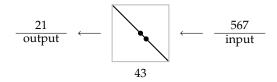
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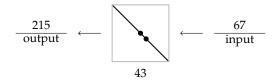
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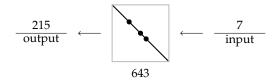
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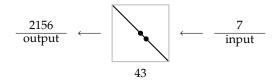
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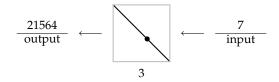
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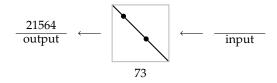
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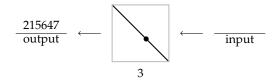
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The stack can be thought of as a container that always holds a decreasing permutation, where at any time we can:

- push a new maximum entry into the container such that the container holds a decreasing permutation
- pop the leftmost entry out of the container.

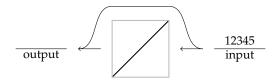


Since the container holds permutations that avoid the pattern 12, we call this the Av(12)-machine.

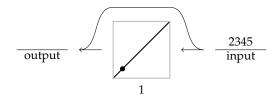
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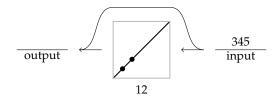
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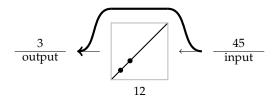
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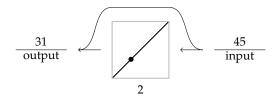
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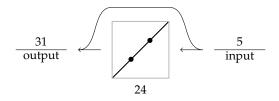
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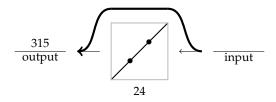


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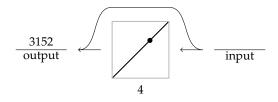
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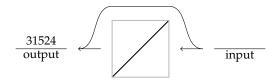
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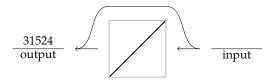
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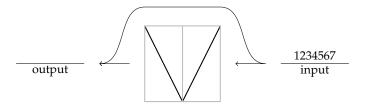


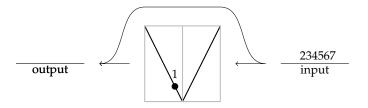
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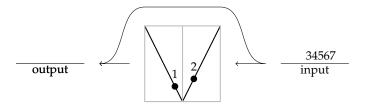
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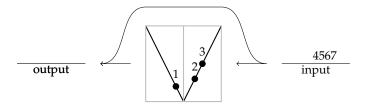


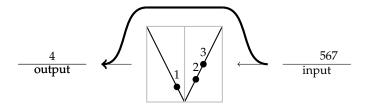
The Av(21)-machine generates the class Av(321).

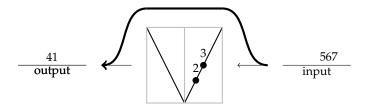


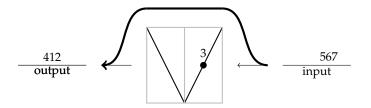








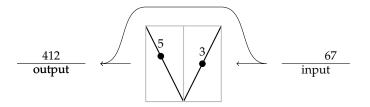


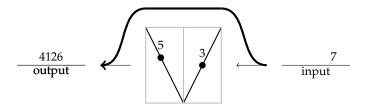


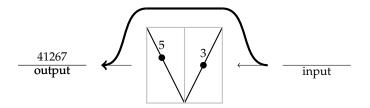
Exotic Machines

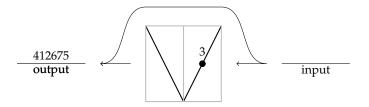
What else?

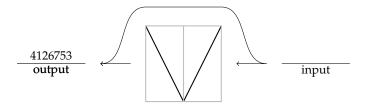
Example: The Av(231, 132)-Machine











Whenever the machine is nonempty, there are always two places to push a new entry.

The Av(231, 132)-machine generates the class Av(4231, 4132).

► The Av(12)-machine generates the class Av(312).

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Theorem. The Av(B)-machine generates the class

 $Av(\{{}^+\beta:\beta\in B\}),$

where ${}^+\beta$ is formed by adding a new maximum in front of β .

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 $Av(\{{}^+\beta:\beta\in B\}),$

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Example: The Av(123, 4132)-machine generates the class Av(4123, 54132).

Structural Descriptions:

Every permutation in C can be formed by ...

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Av(312): $\approx 4^n$

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Instead of remembering all $\approx 4^n$ permutations of length n in Av(312) to build those of length n + 1, we only need to remember the n + 1 possible states of the machine:

[machine has no entries], [machine has 1 entry],

[machine has n entries]

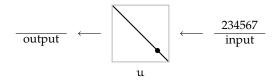
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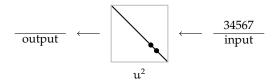
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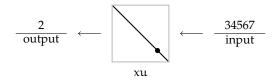
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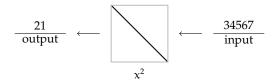
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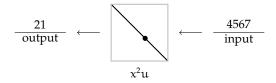
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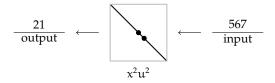
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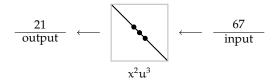
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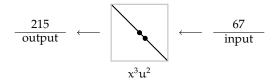
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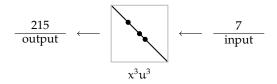
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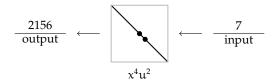
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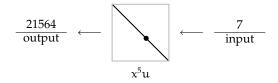
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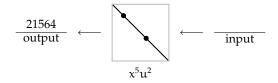
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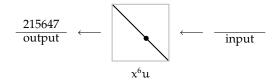
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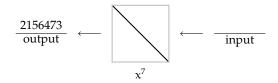


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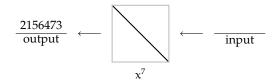
Let f(x, u) be the generating function for valid states of the Av(12)-machine where u tracks the number of entries in the machine and x tracks the number that have been output so far.



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Hence the generating function for these states is f(x, 0).

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This translates to a functional equation:

$$f(x, u) = 1 + uf(x, u) + \frac{x}{u} (f(x, u) - f(x, 0))$$

The kernel method solves: $f(x, 0) = \frac{1 - \sqrt{1 - 4x}}{2x}$.

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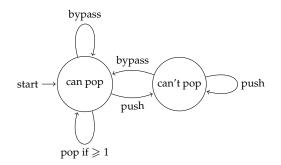
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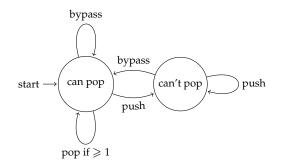
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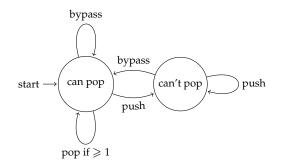
- ► a bypass is always preferred when possible,
- an entry should be popped as soon as possible.

Equivalently: once you push an entry, you cannot pop any entries until you have bypassed an entry.





[can pop]: $f(x, u) = 1 + x(f(x, u) + g(x, u)) + \frac{x}{u}(f(x, u) - f(x, 0)),$ [can't pop]: g(x, u) = u(f(x, u) + g(x, u)).



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Solve g(x, u) in terms of $f(x, u) \rightarrow$ substitute into $f(x, u) \rightarrow$ kernel method \rightarrow same answer!

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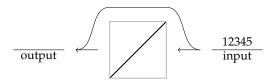
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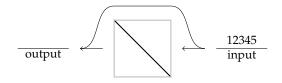
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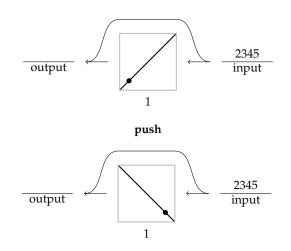
(To make this easier to describe, we tack the bypass back onto the Av(12)-Machine, even though it makes no difference.)

TATION CLASSES	C-Machines	Exotic Machines	What else?
	000000000000000000000000000000000000000		

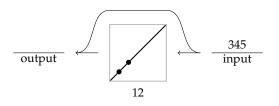




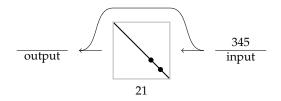
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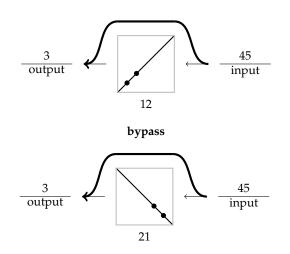


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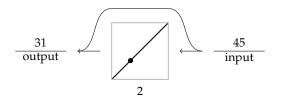




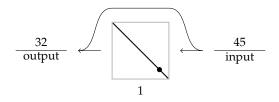




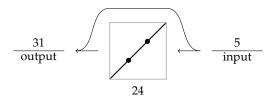
Permutation Classes	C-Machines	Exotic Machines	What else?
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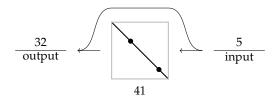




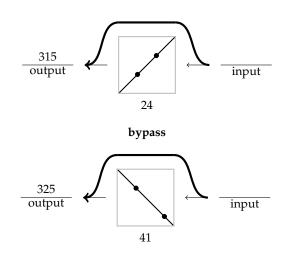
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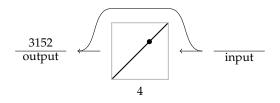




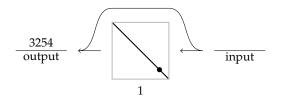
ERMUTATION CLASSES	C-Machines	Exotic Machines	WHAT ELSE?
	000000000000000000000000000000000000000		



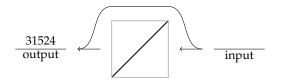
ERMUTATION CLASSES	C-Machines	Exotic Machines	WHAT ELSE?
	000000000000000000000000000000000000000		00000000



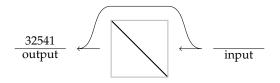




Permutation Classes	C-Machines	Exotic Machines	What else?
	000000000000000000000000000000000000000		00000000







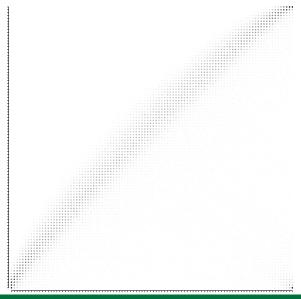
RANDOM SAMPLING

Any combinatorial family that can be counted with *dynamic programming* can be sampled from uniformly at random.

C-Machines

Exotic Machines

RANDOM SAMPLING

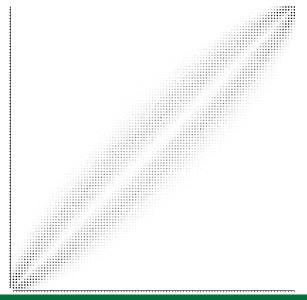


22 of 47

C-Machines

Exotic Machines

RANDOM SAMPLING



23 of 47

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Of these 38 classes, the exact enumerations are known for 35 of them. (More on this later!)

There are ten 2×4 classes that are enumerated by the Schröder numbers: $1, 2, 6, 22, 90, \ldots$

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Example: Av(231, 132).

A small adjustment to the earlier functional equations gives:

$$\begin{split} f(x,u) &= 1 + x(f(x,u) + g(x,u)) + \frac{x}{u}(f(x,u) - f(x,0)), \\ g(x,u) &= 2u((f(x,u) - f(x,0)) + g(x,u)) + uf(x,0). \end{split}$$

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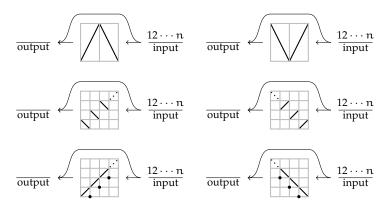
An application of the kernel method gives the generating function for the Schröder numbers.

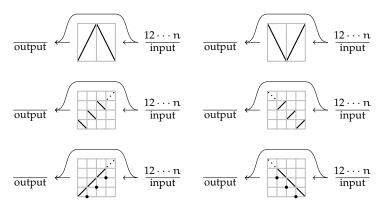
$$f(x,0) = \frac{3 - x - \sqrt{1 - 6x + x^2}}{2}$$

C-Machines

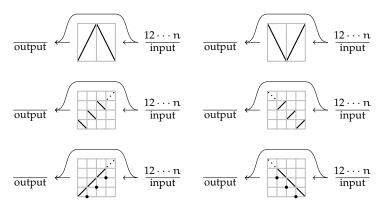
Exotic Machines

Enumerating the Schröder Classes





All are enumerated by the Schröder numbers.



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Free bijections between every pair of classes.

$2{\times}4\ C \text{Lasses}$

Enumerations are known for all but three:

- ► Av(4123,4231)
- ► Av(4123,4312)
- ► Av(4231,4321)

$2{\times}4\ C \text{Lasses}$

Enumerations are known for all but three:

- ► Av(4123,4231)
- ► Av(4123,4312)
- ► Av(4231,4321)

These three can all be modeled with C-machines.

Av(4123, 4231)

Av(4123, 4231) is generated by the Av(123, 231)-machine.

The class Av(123, 231) is a geometric grid class.

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Av(4123, 4231) is generated by the Av(123, 231)-machine.

The class Av(123, 231) is a geometric grid class.



The states of the machine can be represented by 4-tuples (a, b, c, P) where a, b, and c are the number of a, b, c entries, and P is a boolean representing whether we can or can't pop.

C-Machines

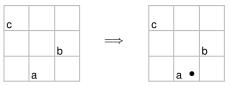
Exotic Machines

What else?

Av(4123, 4231)



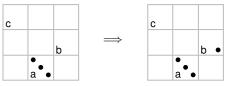
Av(4123, 4231)



We can describe the state transitions:

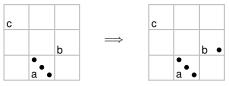
• $(0,0,0,T) \longrightarrow \{(1,0,0,F), (0,0,0,T)\}$

Av(4123, 4231)



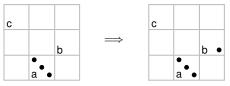
- ► $(0,0,0,T) \longrightarrow \{(1,0,0,F), (0,0,0,T)\}$
- $(a,0,0,F) \longrightarrow \{(a+1,0,0,F), (a,1,0,F), (a,0,0,T)\}$

Av(4123, 4231)



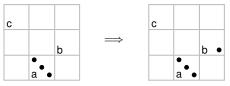
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$\operatorname{Av}(4123,4231)$



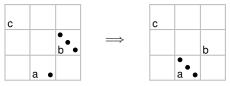
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- $(a, b, 0, F) \longrightarrow \{(a, b+1, 0, F), (a, b, 1, F), (a, b, 0, T)\}$

Av(4123, 4231)



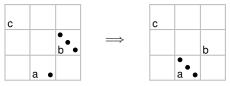
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$\operatorname{Av}(4123,4231)$



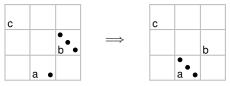
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Av(4123, 4231)



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Av(4123, 4231)



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With a little work, we can translate the state transitions into a set of functional equations:

$$\begin{aligned} A(a,x) &= 1 + \frac{x}{a} (A(a,x) - A(0,x)) + aA(a,x) + xB(0,a,x), \\ B(a,b,x) &= \frac{1}{a} (A(a,x) - A(0,x)) \frac{bC}{1-b} + B(a,b,x) \frac{bC}{1-b} \\ &+ \frac{x}{a} (1+C) (B(a,b,x) - B(0,b,x)). \end{aligned}$$

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The generating function for the class is A(0, x), but we have no idea how to solve these equations.

Using the state transitions, dynamic programming, and Amazon cloud computing, we have computed the first 1000 terms of Av(4123, 4231).

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We've written software that takes initial terms of a sequence and tries to produce a conjecture for the actual generating function. It searches the space of rational, algebraic, D-finite, and Dalgebraic functions.

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We've written software that takes initial terms of a sequence and tries to produce a conjecture for the actual generating function. It searches the space of rational, algebraic, D-finite, and Dalgebraic functions.

Despite having 1000 initial terms of Av(4123, 4231), we can't guess any generating function!

Exotic Machines

Three More Exotic Machines



600 terms Av(4231,4321)

Exotic Machines

What else?

Three More Exotic Machines



600 terms Av(4231, 4321)



1000 terms Av(4123,4312)

Exotic Machines

Three More Exotic Machines



600 terms Av(4231, 4321)



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5000 terms Av(4123,4231,4312)

Exotic Machines

What else?

Still...

Even though we don't have a generating function, we can still:

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Sample uniformly at random

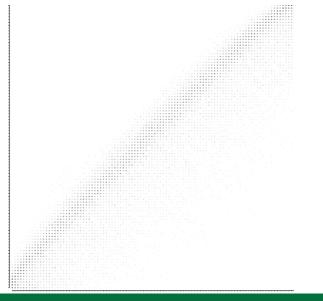
Still...

Even though we don't have a generating function, we can still:

- Sample uniformly at random
- Use many initial terms to approximate the singularity structure and asymptotic behavior

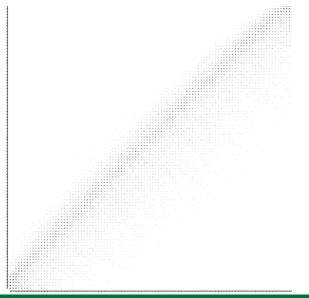
Exotic Machines

Av(4123, 4231)



Exotic Machines

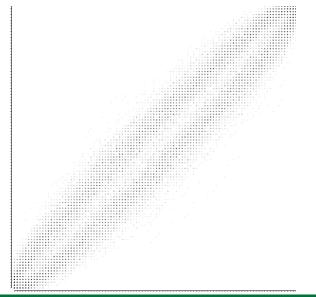
Av(4123, 4312)



Exotic Machines

What else?

Av(4231,4321)



SINGULARITY STRUCTURE – AV(4123,4231,4312) Dominant singularity at

with critical exponent

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Other singularities:

- 0.2380637598080281627623423301122200693506416354296582730015127601332...
- ► 0.2425743178578198539792244185837648673615631138554928145039424175060...
- ► 0.244744854382140527597173051613305433663865887080858073339928...
- ► 0.24600369400712937477975290028099807844398...
- 0.2468170919259163688566314680...
- ► 0.2473814067574507777...
- ► 0.247793229299...
- ▶ 0.2482...

SINGULARITY STRUCTURE – AV(4123,4231,4312) Dominant singularity at

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Other singularities:

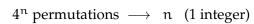
- 0.2380637598080281627623423301122200693506416354296582730015127601332...
- ► 0.2425743178578198539792244185837648673615631138554928145039424175060...
- ► 0.244744854382140527597173051613305433663865887080858073339928...
- 0.24600369400712937477975290028099807844398...
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- 0.2473814067574507777...
- ▶ 0.247793229299...
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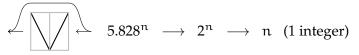
 $\implies \sim C(4.468408...)^n$



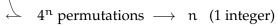
 4^n permutations $\longrightarrow n$ (1 integer)

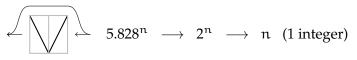


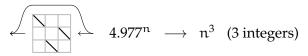




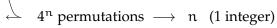


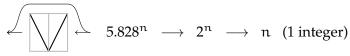


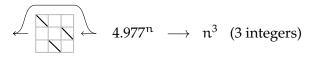




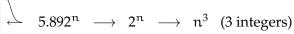








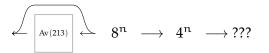


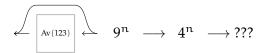


$$\overbrace{\mathsf{Av}(231)}^{\mathsf{Av}(231)} \xleftarrow{} 11.6^{\mathsf{n}}(?) \longrightarrow 4^{\mathsf{n}} \longrightarrow ???$$

$$\overbrace{\text{Av}(231)}^{\text{Av}(231)} \xleftarrow{} 11.6^{n}(?) \longrightarrow 4^{n} \longrightarrow ???$$

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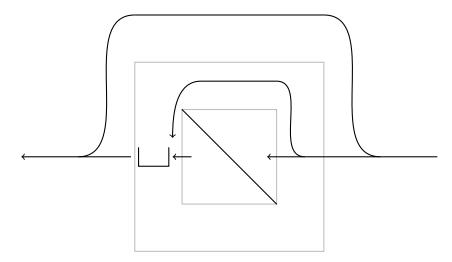
$$\overbrace{Av(213)}^{Av(213)} \xleftarrow{} 8^n \longrightarrow 4^n \longrightarrow ???$$

$$\overbrace{A^{v(123)}}^{A^{v(123)}} \xleftarrow{} 9^{n} \longrightarrow 4^{n} \longrightarrow ???$$

$$\overbrace{A^{V(312)}}^{A^{V(312)}} \xleftarrow{} 9^{n} \longrightarrow 4^{n} \longrightarrow n^{2} \quad (2 \text{ integers}) !!!$$

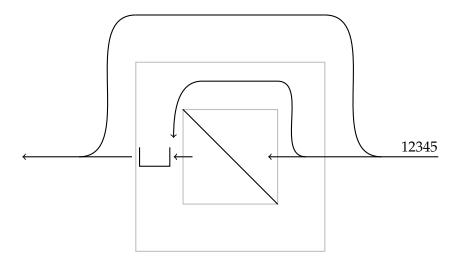
Exotic Machines

Nested Av(12)-Machine



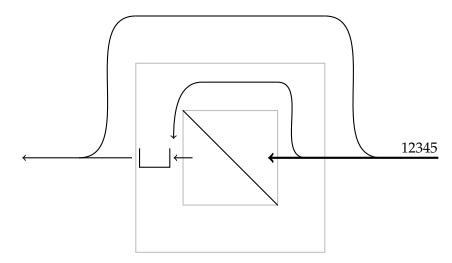
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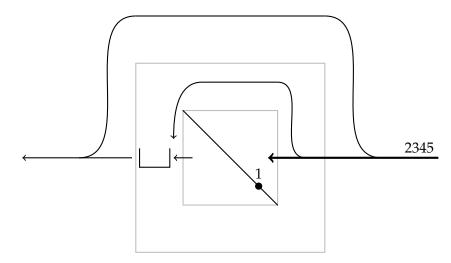
Exotic Machines

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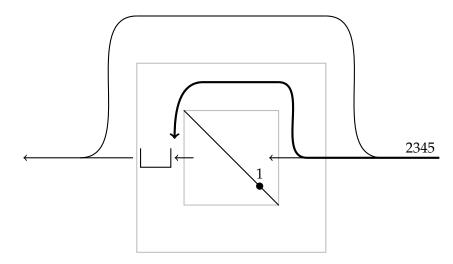
Exotic Machines

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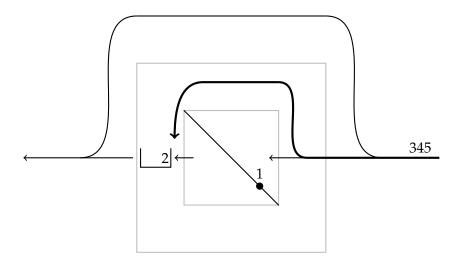
Exotic Machines

Nested Av(12)-Machine



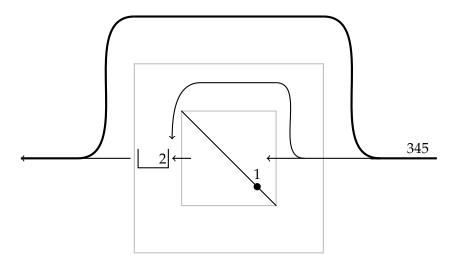
Exotic Machines

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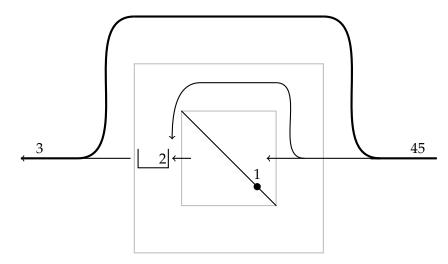
Exotic Machines

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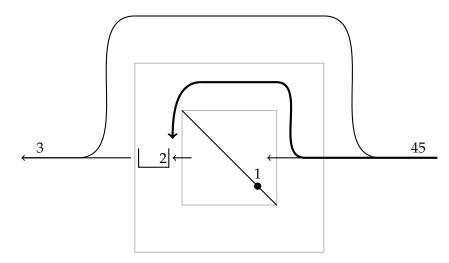
Exotic Machines

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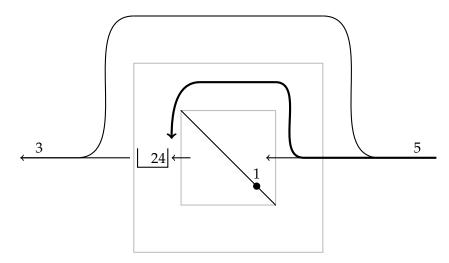
Exotic Machines

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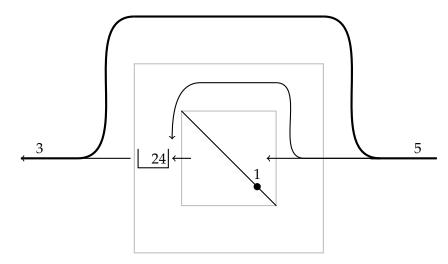
Exotic Machines

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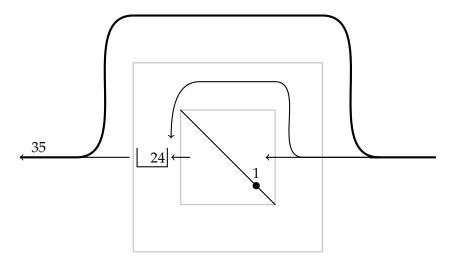
Exotic Machines

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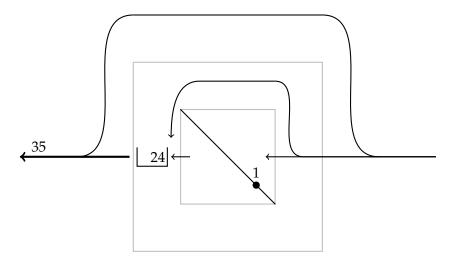
Exotic Machines

Nested Av(12)-Machine



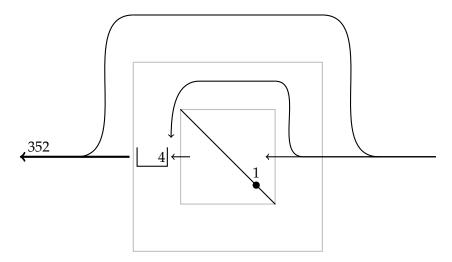
Exotic Machines

Nested Av(12)-Machine



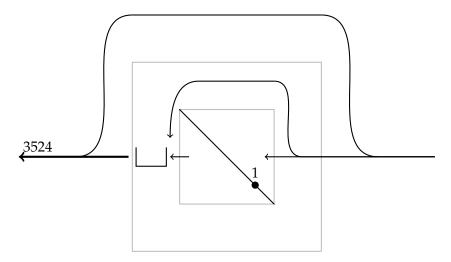
Exotic Machines

Nested Av(12)-Machine



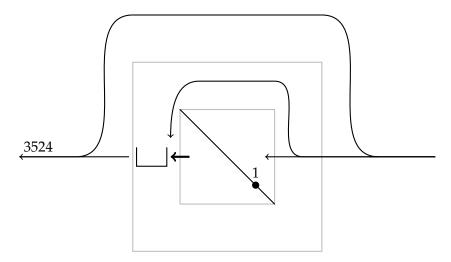
Exotic Machines

Nested Av(12)-Machine



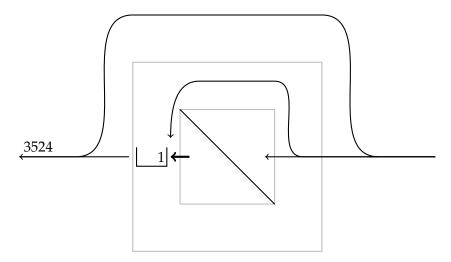
Exotic Machines

Nested Av(12)-Machine



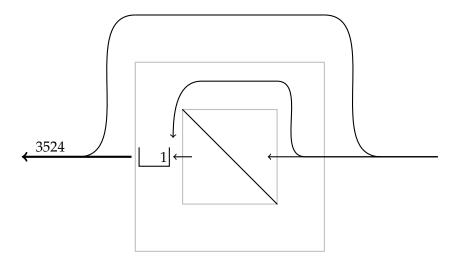
Exotic Machines

Nested Av(12)-Machine



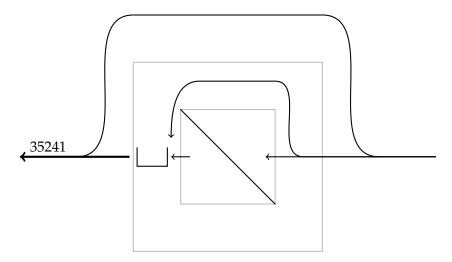
Exotic Machines

Nested Av(12)-Machine



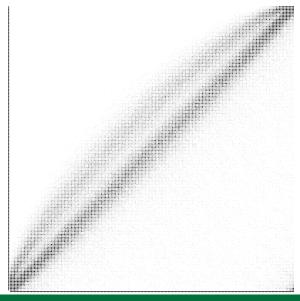
Exotic Machines

Nested Av(12)-Machine



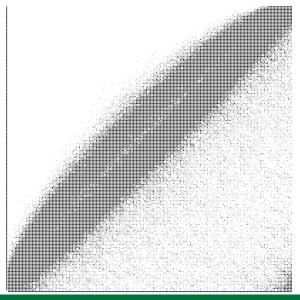
Exotic Machines

RANDOM SAMPLING



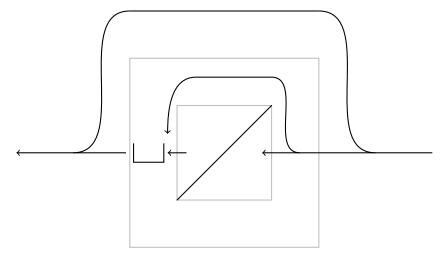
Exotic Machines

RANDOM SAMPLING



Exotic Machines

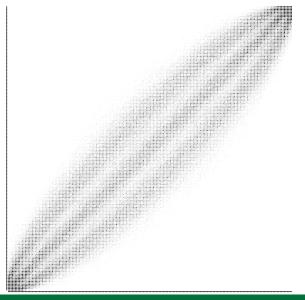
Nested Av(21)-Machine



Generates Av(4321) Free bijection to Av(4312)!

Exotic Machines

RANDOM SAMPLING



Exotic Machines

What about Av(4231)?

Nesting doesn't work.

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Some heuristics suggest state compression isn't possible with a standard \mathcal{C} -Machine.

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Some heuristics suggest state compression isn't possible with a standard \mathcal{C} -Machine.

Perhaps a more complicated machine: second bypass, add a buffer or a stack, multiple moves at once, ...

Thanks!