

# Typing a semantic memory for mathematical content

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## Outline

- (1) The MoSMath project
- (2) The semantic memory
- (3) The type system
- (4) Applications

## MoSMath

“A modeling system for mathematics” (MoSMath).

Goal: a modeling system for the specification of models for the numerical work.

Input: a controlled natural language, formulas in a subset of  $\text{L}^{\text{A}}\text{T}_{\text{E}}\text{X}$ .

Output: Description in  $\text{L}^{\text{A}}\text{T}_{\text{E}}\text{X}$ , model-file in AMPL, etc.

Advantages:

No need to learn an algebraic modeling language

Specification is the least error prone, and the most natural

We expect that the framework of the MoSMath project will serve as a first step towards the FMathL project.

<http://www.mat.univie.ac.at/~neum/FMathL.html>

## **The semantic memory**

We want to be able to represent mathematical expressions and mathematical natural language.

We use a directed labeled graph, special case of a semantic network, implementable in the semantic web.

We assume an unrestricted set of **objects**.

Objects may, but need not have **names**.

Objects may, but need not have **external values**.

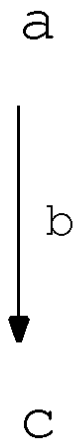
We refer to **unnamed objects** via a string beginning with a dollar-sign (\$).

The **semantic memory** (abbreviated SM) stores equations of the form

$$a.b=c$$

These are called **sems**.

Graphical: a sem  $a.b=c$  as an edge with label  $b$  from node  $a$  to node  $c$ .



A sem  $a.b=c$  usually means: “the  $b$  of  $a$  is  $c$ ”.

In terms of the semantic web: “the  $a$  has property  $b$  valued by  $c$ ”.

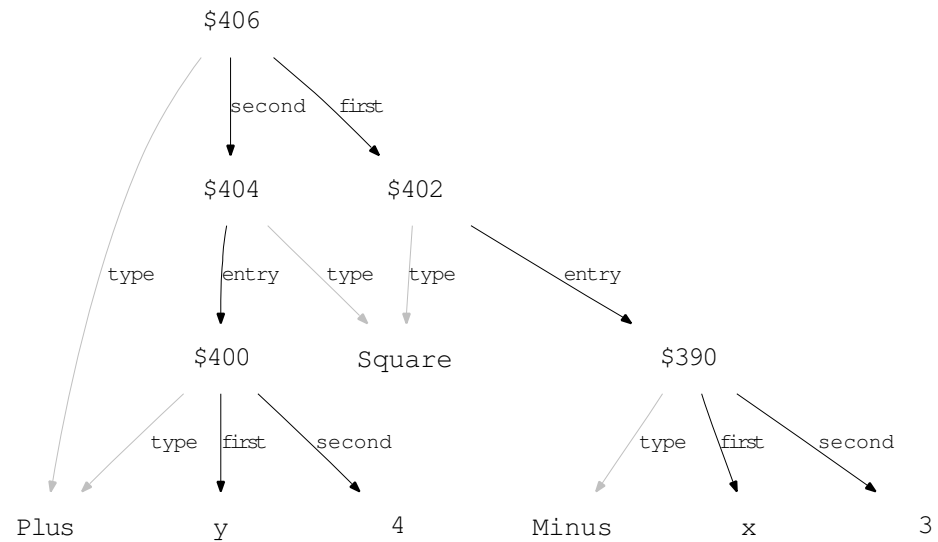
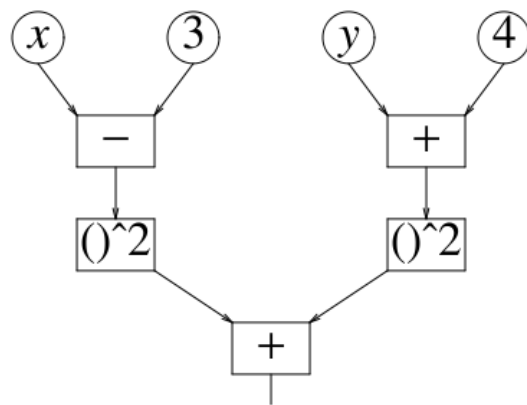
E.g., the sem  $formula27.label=CauchySchwarz$  would intuitively mean that the  $label$  of  $formula27$  is  $CauchySchwarz$ .

Only restriction: no two arcs beginning at the same node may have the same label.

In particular, cycles are allowed.

General idea for representing formulas: automatic differentiation.

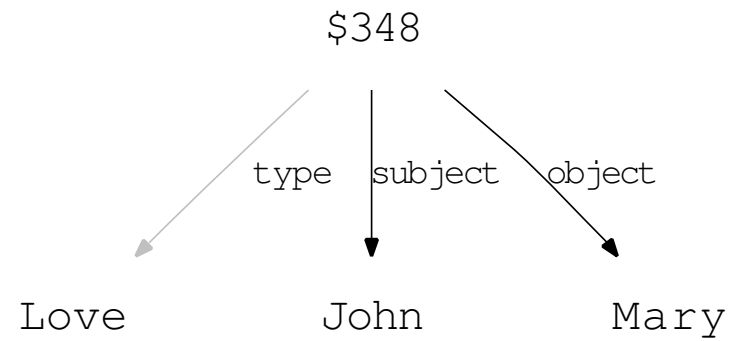
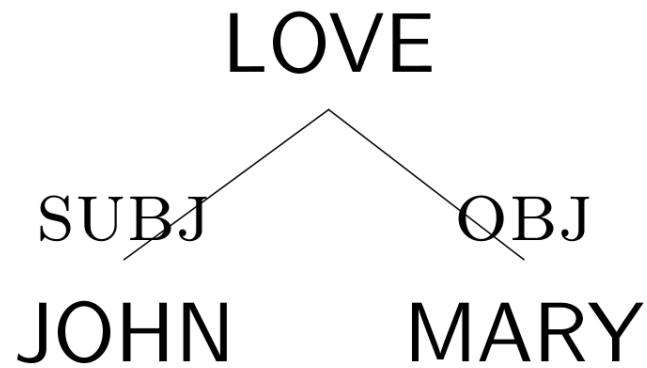
$$(x - 3)^2 + (y + 4)^2$$



Source: D. Gay, Using Expression Graphs In Optimization Algorithms

General idea for representing natural language: dependency grammar.

“John loves Mary.”



### **The semantic virtual machine:**

A virtual machine that operates on the SM.

The algorithms for the semantic virtual machine are also graphs in the semantic memory.

## Typing

An essential step that brings formal structure into the semantic memory is the introduction of **types**.

The set of sems reachable from object  $o$  are called the **record** with handle  $o$ .

We want to give a criterion when a record is “well-formed”.

Categories are objects, ordered by a transitive partial relation  $<$ . ( $C1 < C2$  means  $C1$  is **contained** in  $C2$ ).

A category is called a **type** or an **atomic** if it is minimal in the ordering  $<$ , and a **union** otherwise.

Atomics do not have any children, they have semantic meaning in itself, types pose requirements on records.

If  $x.type=y$  then the requirements of  $y$  apply to  $x$ .

Categories are defined in text documents called **type sheets**.

Example: we define types `BinaryRel` and `Integer`,  
atomics `LessEq`, `Equal` and `Less`  
and the union `RelationSym`.

`RelationSym`:

```
atomic> LessEq, Equal, Less
```

`BinaryRel`:

```
allOf> lhs = Integer  
      rhs = Integer  
      relation = RelationSym
```

`Integer`:

```
nothingElse
```

All information about the categories is represented in the semantic memory  
→ type checking is an algorithm on the semantic memory.

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operator	usage
<code>allOf</code>	requirements to all of a set of fields
<code>oneOf</code>	requirements to exactly of a set of fields
<code>someOf</code>	requirements to at least one of a set of fields
<code>optional</code>	requirements to a set of fields, if nonempty
<code>fixed</code>	requires not types but certain objects
<code>someOfType</code>	requirements to all fields of a type
<code>itself</code>	requires a set of fields which are also the entries
<code>nothingElse</code>	forbids non-required fields
<code>nothing</code>	defines an atomic type
<code>union</code>	defines a union
<code>complete</code>	closes a union

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The type definitions are themselves typed, i.e., the creation of a type sheet that allows to type the representation of types in the semantic matrix, which gives the type of types.

The type sheet for types has only 40 lines.

Meta-schema for RelaxNG schema has over 300 lines, for Meta-DTD over 700 lines.

## Application 1: The OR Library

A significant fraction of the OR Library was represented manually.

Algorithms produce  $\text{\LaTeX}$  and AMPL.

Example: multi-dimensional knapsack problem.

Let the integer  $N$  be the number of contracts, let the integer  $M$  be the number of budgets. Let  $c_j$  be the contract volume of project  $j$  for  $j = 1, \dots, N$ , let  $A_{i,j}$  be the estimated cost of budget  $i$  for project  $j$  for  $i = 1, \dots, M$  and  $j = 1, \dots, N$ , and let  $B_i$  be the available amount of budget  $i$  for  $i = 1, \dots, M$ . For  $j = 1, \dots, N$ , let  $x_j = 1$  if project  $j$  is selected, and let  $x_j = 0$  otherwise.

**Problem :** Given integers  $N$  and  $M$ , vector  $c$ , matrix  $A$  and vector  $B$ , find the binary vector  $x$  such that

$$\sum_{j=1}^N c_j x_j$$

is maximal under the constraint  $\sum_{j=1}^N A_{i,j} x_j \leq B_i$  for  $i = 1, \dots, M$ .

Represented in the semantic memory using 449 sems.

Automatically generated AMPL-output:

```
param N ;
param M ;
param c{j in 1..N} ;
param A{i in 1..M , j in 1..N} ;
param B{i in 1..M} ;
var x{j in 1..N} binary ;

maximize target : sum{j in 1..N}(c[j] * x[j]);
subject to constraint_1{i in 1..M} : sum{j in 1..N}(A[i ,
j] * x[j]) <= B[i];
```

## Application 2: A collection of formulas

- mathematical formulas were extracted from lecture notes
- manually fed into the semantic memory
- automatic  $\text{\LaTeX}$ -output

## Application 3: The TPTP Library

- Automatically import formulas from the TPTP library (“Thousands of Problems for Theorem Provers”).
- formulas from different branches of mathematics
- more than 10.000 problem files.

Thank you for your attention!

<http://www.mat.univie.ac.at/~neum/FMathL.html>