

## SOME COMBINATORIAL PROPERTIES OF COMPLETE SEMI-THUE SYSTEMS

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A *reduction system*  $(R, \longrightarrow)$  consists of a set  $R$  and a binary relation  $\longrightarrow$  on  $R$ . Let  $\xrightarrow{*}$  be the reflexive-transitive closure of  $\longrightarrow$  and  $[x]$  the class of an  $x \in R$  with respect to the equivalence generated by  $\longrightarrow$ .  $x$  is called *irreducible* (or *in normal form*) if there is no  $y \in R$  such that  $x \longrightarrow y$ . A reduction system can have the following properties:

- *Chain Condition*: There is no infinite chain  $x_1 \longrightarrow x_2 \longrightarrow x_3 \longrightarrow \dots$  in  $R$ . (Then  $\longrightarrow$  is called *terminating* or *Noetherian*.)
- *Confluence*:  $\forall w, x, y \in R: (w \xrightarrow{*} x \wedge w \xrightarrow{*} y \Rightarrow \exists z \in R: x \xrightarrow{*} z \wedge y \xrightarrow{*} z)$ .
- *Completeness*: Chain condition and confluence.

If a reduction system is complete, normal forms always exist and are unique. See [3] for further details.

Let  $\Sigma$  be a finite alphabet.  $\Sigma^*$  denotes the free monoid over  $\Sigma$  and  $\square$  the empty word. A *semi-Thue system* (STS) on  $\Sigma$  is a subset  $S \subseteq \Sigma^* \times \Sigma^*$ . Each element  $(u, v)$  of  $S$  is called a *rule* and written in the form  $u \longrightarrow v$ . An STS  $S$  defines a reduction relation  $\longrightarrow$  on  $\Sigma^*$  by  $xuy \longrightarrow xvy \Leftrightarrow (u, v) \in S$ .

Let  $OV(u) = \{x \in \Sigma^* \mid \exists y, z \in \Sigma^* : u = yx = xz\} \setminus \{\square, u\}$  be the set of *non-trivial self-overlaps* of  $u \in \Sigma^*$ . Generalizing results of Book [2], Otto and Wrathall [6], one obtains the following.

**Theorem.** *Let the single-rule STS  $u \longrightarrow v$  satisfy the chain condition, and let  $u = u_0u_1u_2 \dots u_k$  ( $k \geq 0$ ), such that  $OV(u) = \{u_1u_2 \dots u_k, u_2 \dots u_k, \dots, u_k\}$ . The STS is confluent if and only if one of the following two conditions is satisfied:*

- (a)  $v$  has  $u_1u_2 \dots u_k$  as a self-overlap, or
- (b) there is a  $j \in \{1, 2, \dots, k+1\}$ , such that  $v = u_ju_{j+1} \dots u_k$  (for  $j = k+1$ :  $v = \square$ ) and  $u = u_{j-1}^j u_j u_{j+1} \dots u_k$ .

For the case  $v = \square$ , this means that  $u$  must be a power of a word  $y$  without proper self-overlap [2]. The classes  $|w|$  of such complete systems  $y^r \longrightarrow \square$  are deterministic context-free languages [1]. One can show that the unambiguous grammar  $(\Sigma \cup \{S\}, \Sigma, P, S)$  with  $P = \{S \longrightarrow \square, S \longrightarrow (a_1Sa_2S \dots a_{k-1}Sa_k)^r S\}$ , where  $y = a_1a_2 \dots a_k$  ( $a_i \in \Sigma$ ), generates  $|w|$ . (There is a similar grammar for the general case  $|w|$ .) From this presentation it follows that the structure generating function  $S(z)$  (cf. [4]) of  $|w|$  is the unique solution of the equation  $S(z) = 1 + z^{rk}(S(z))^{r(k-1)+1}$  in  $\mathbb{Z}[[z]]$ ,

which is a variant of the well-known “trinomial equation” (T):  $A(x) = 1 + x(A(x))^t$  ( $t \in \mathbb{N}$ ). (T) has the unique solution  $A(x) = \sum_{n=0}^{\infty} a_n x^n$  with  $a_n = \frac{1}{n(t-1)+1} \binom{tn}{n}$  (the  $a_n$  having a lot of combinatorial interpretations, see, e.g., [5]). This is usually proved by the Lagrange inversion formula, but it can also be deduced from the set equation  $[a^p] = [a^{p-1}]a + [a^{p+k-1}]b$  for the special STS  $a^{k-1}b \rightarrow \square$ . The enumeration of words in the classes  $[w]$ ,  $w$  irreducible, can be carried out in the same manner.

## REFERENCES

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