

Forests and Parking Functions

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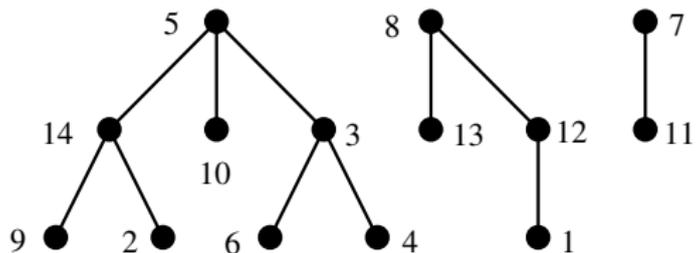
Outline

- 1 Forests
- 2 Parking Functions
- 3 Statistics
- 4 The Map $\varphi : F_n \rightarrow PF_n$
- 5 Further Result

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(Rooted Labeled) Forests



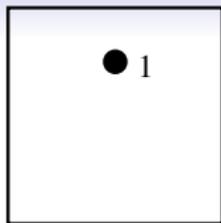


Figure: Forest on [1]

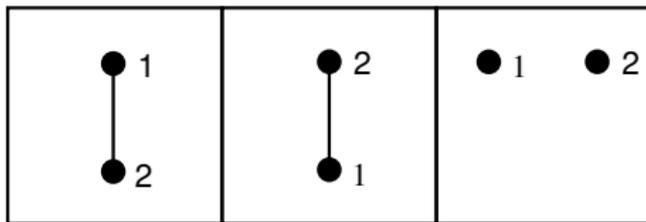


Figure: Forests on [2]

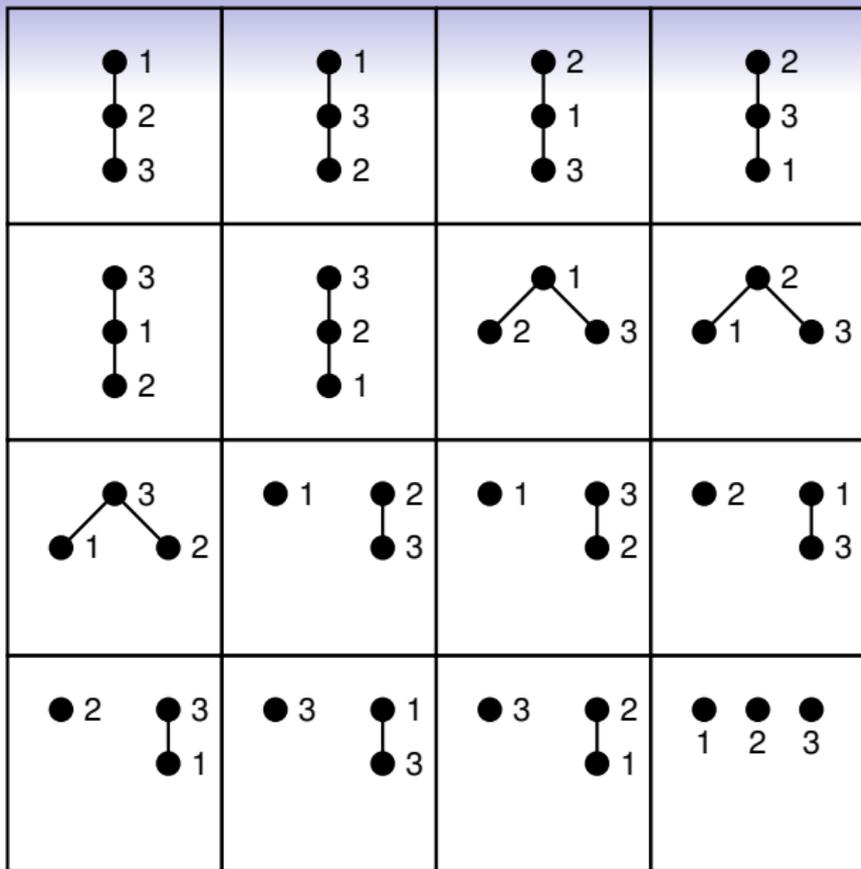


Figure: Forests on [3]

$$\begin{aligned}
 \# \text{ of forests on } [1] &= 1 = 2^0 \\
 \# \text{ of forests on } [2] &= 3 = 3^1 \\
 \# \text{ of forests on } [3] &= 16 = 4^2 \\
 \# \text{ of forests on } [4] &= 125 = 5^3 \\
 &\vdots \\
 \# \text{ of forests on } [n] &= (n + 1)^{n-1}
 \end{aligned}$$

Actually,

$$\# \text{ of forests on } [n] = \# \text{ of trees on } [n + 1].$$

$$\begin{aligned} \# \text{ of forests on } [1] &= 1 = 2^0 \\ \# \text{ of forests on } [2] &= 3 = 3^1 \\ \# \text{ of forests on } [3] &= 16 = 4^2 \\ \# \text{ of forests on } [4] &= 125 = 5^3 \\ &\vdots \\ \# \text{ of forests on } [n] &= (n + 1)^{n-1} \end{aligned}$$

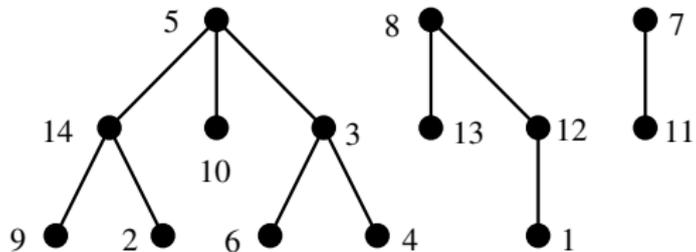
Actually,

$$\# \text{ of forests on } [n] = \# \text{ of trees on } [n + 1].$$

Inversion in Forests

$$\begin{aligned}\text{inv}(F; v) &= \#\{(v, u) \mid u \text{ is descendant of } v \text{ and } u < v\} \\ \text{inv}(F) &= \sum_v \text{inv}(F; v)\end{aligned}$$

For example,



$$\text{inv}(F; 5) = \#\{(5, 3), (5, 4), (5, 2)\} = 3$$

$$\text{inv}(F) = 7$$

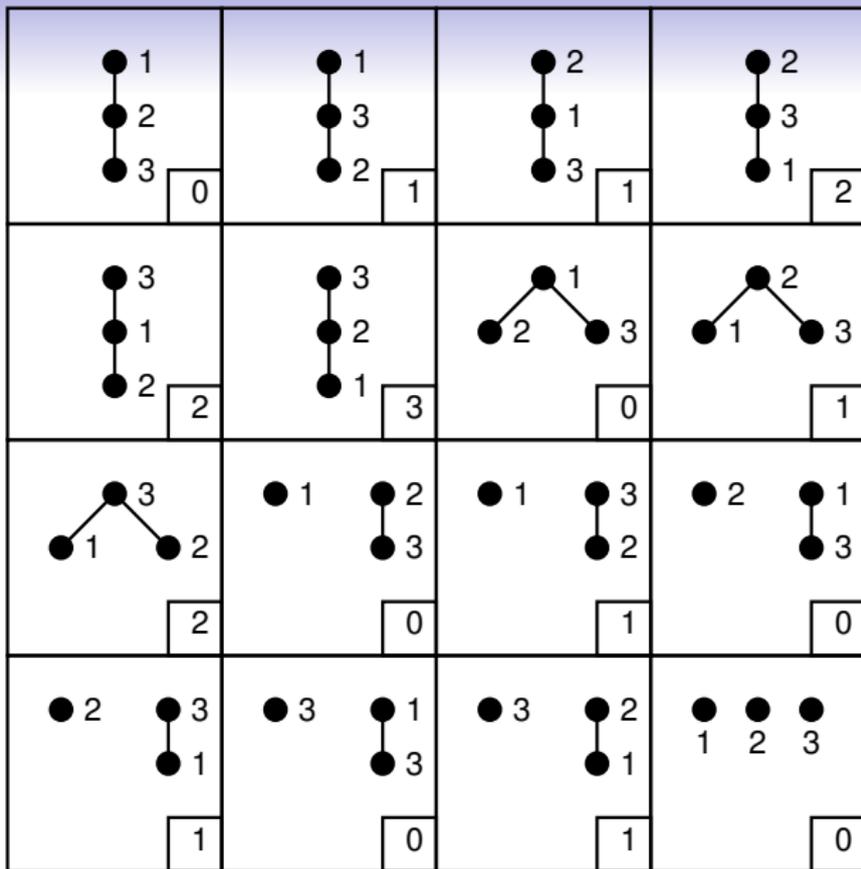
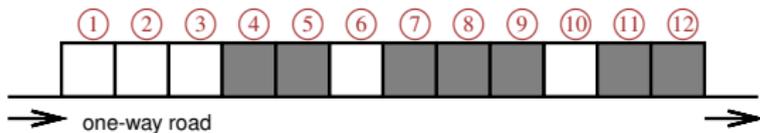


Figure: inv on forests on [3]

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Parking Rule



If you like parking at ⑦, then you could park at ⑩.
If you like parking at ⑪, then you could not park.

Parking Function

Suppose (p_1, \dots, p_n) is a sequence of favorite parking spaces for each cars. If no car failed in parking, (p_1, \dots, p_n) is called **parking function**.

For example,

$$(4, 3, 3, 1, 4) \rightarrow \boxed{4} \boxed{\emptyset} \boxed{2} \boxed{1} \boxed{3}$$

is not a parking function.

$$(4, 3, 3, 1, 1) \rightarrow \boxed{4} \boxed{5} \boxed{2} \boxed{1} \boxed{3}$$

is a parking function.

Criteria of Parking Function

$$q_i \leq i \text{ for all } i$$

where (q_1, \dots, q_n) is rearrangement of (p_1, \dots, p_n) by order.

For example,

$$(4, 3, 3, 1, 4) \rightarrow (1, 3, 3, 4, 4)$$

is not a parking function.

$$(4, 3, 3, 1, 1) \rightarrow (1, 1, 3, 3, 4)$$

is a parking function.

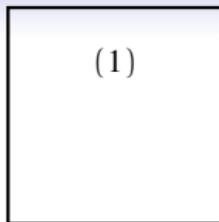


Figure: Parking Function with length 1

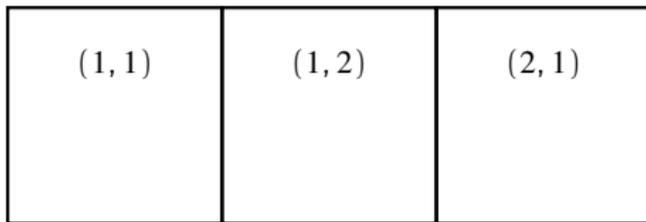


Figure: Parking Functions with length 2

| | | | |
|-----------|-----------|-----------|-----------|
| (1, 1, 1) | (1, 1, 2) | (1, 1, 3) | (1, 2, 1) |
| (1, 2, 2) | (1, 2, 3) | (1, 3, 1) | (1, 3, 2) |
| (2, 1, 1) | (2, 1, 2) | (2, 1, 3) | (2, 2, 1) |
| (2, 3, 1) | (3, 1, 1) | (3, 1, 2) | (3, 2, 1) |

Figure: Parking Functions with length 3

Jump in Parking Functions

- 1 $PA(p_1, \dots, p_n) = (q_1, \dots, q_n)$
where q_i is the space parked actually by i -th car. [▶ here](#)
- 2 $\text{jump}(P; i) = q_i - p_i$
- 3 $\text{jump}(P) = \sum_i \text{jump}(P; i)$

Note that,

$$\begin{aligned}\text{jump}(P) &= \sum_i \text{jump}(P; i) \\ &= (q_1 + \dots + q_n) - (p_1 + \dots + p_n) \\ &= \binom{n+1}{2} - (p_1 + \dots + p_n)\end{aligned}$$

For example,

① $PA(4, 3, 3, 1, 1) = (4, 3, 5, 1, 2)$

② $\text{jump}(4, 3, 3, 1, 1; 3) = 5 - 3 = 2$

③ $\text{jump}(4, 3, 3, 1, 1) = 0 + 0 + 2 + 0 + 1 = 3$

| | | | |
|-----------|-----------|-----------|-----------|
| (1, 1, 1) | (1, 1, 2) | (1, 1, 3) | (1, 2, 1) |
| 3 | 2 | 1 | 2 |
| (1, 2, 2) | (1, 2, 3) | (1, 3, 1) | (1, 3, 2) |
| 1 | 0 | 1 | 0 |
| (2, 1, 1) | (2, 1, 2) | (2, 1, 3) | (2, 2, 1) |
| 2 | 1 | 0 | 1 |
| (2, 3, 1) | (3, 1, 1) | (3, 1, 2) | (3, 2, 1) |
| 0 | 1 | 0 | 0 |

Figure: jump on Parking Functions with length 3

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Generating Function

- GF for inv on forests

$$\sum_{F \in F_1} q^{\text{inv}(F)} = 1$$

$$\sum_{F \in F_2} q^{\text{inv}(F)} = 2 + q$$

$$\sum_{F \in F_3} q^{\text{inv}(F)} = 6 + 6q + 3q^2 + q^3$$

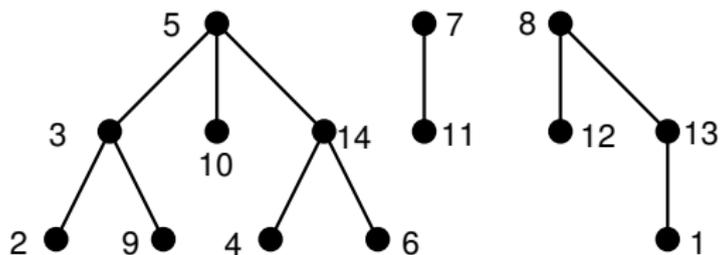
- GF for jump on parking function

$$\sum_{P \in PF_1} q^{\text{jump}(P)} = 1$$

$$\sum_{P \in PF_2} q^{\text{jump}(P)} = 2 + q$$

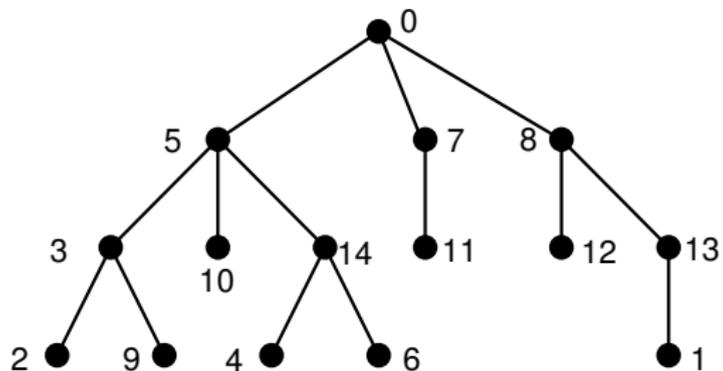
$$\sum_{P \in PF_3} q^{\text{jump}(P)} = 6 + 6q + 3q^2 + q^3$$

One Map



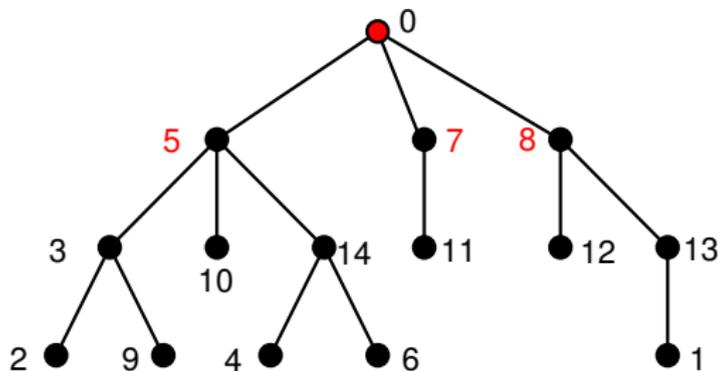
(_ _ _ _ _ _ _ _ _ _ _ _ _ _)

One Map



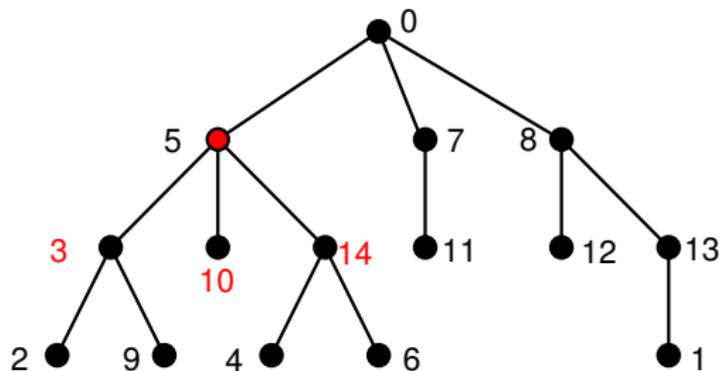
(_ _ _ _ _)

One Map



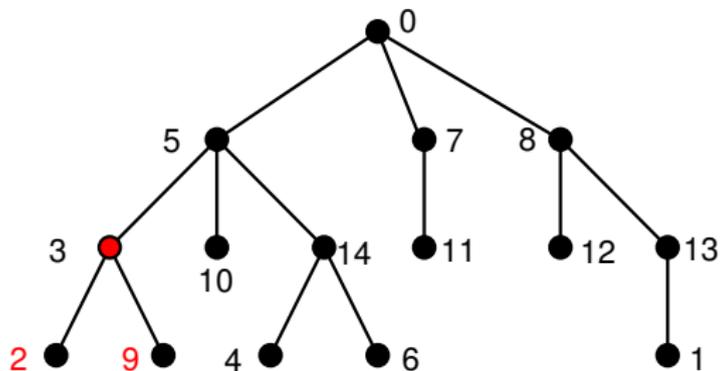
(, , , , 1 , , 1 , 1 , , , , , ,)

One Map



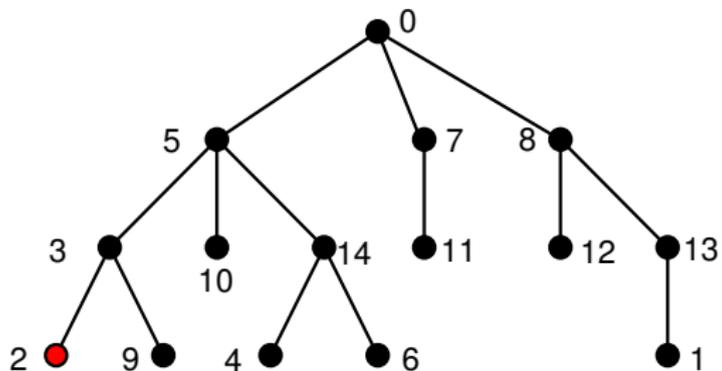
(_, _, 2, _, 1, _, 1, 1, _, 2, _, _, 2)

One Map



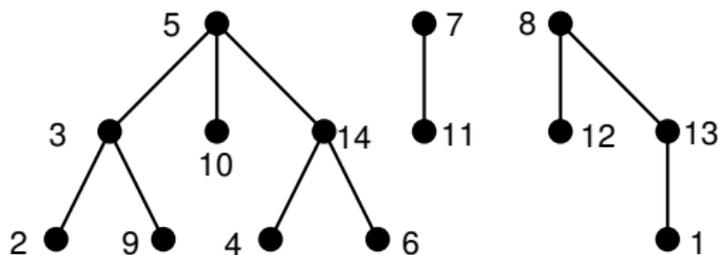
(_, 3, 2, _, 1, _, 1, 1, 3, 2, _, _, _, 2)

One Map



(_, 3, 2, _, 1, _, 1, 1, 3, 2, _, _, _, 2)

One Map



(14, 3, 2, 7, 1, 7, 1, 1, 3, 2, 10, 12, 12, 2)

Theorem (G.Kreweras 1980)

$$\sum_{F \in \mathcal{F}_n} q^{\text{inv}(F)} = \sum_{P \in \mathcal{PF}_n} q^{\binom{n+1}{2} - |P|} \quad (1)$$

NOTE. In 2004, R. Stanley notices that a **nonrecursive** bijection between forests and parking functions would be greatly preferred, which yields (1)

Leader in Forests

Definition (Leader in Forests)

- 1 $v = \text{leader}$ in $F \Leftrightarrow \text{inv}(F; v) = 0$.
- 2 $\text{lead}(F) =$ the # of leaders in F

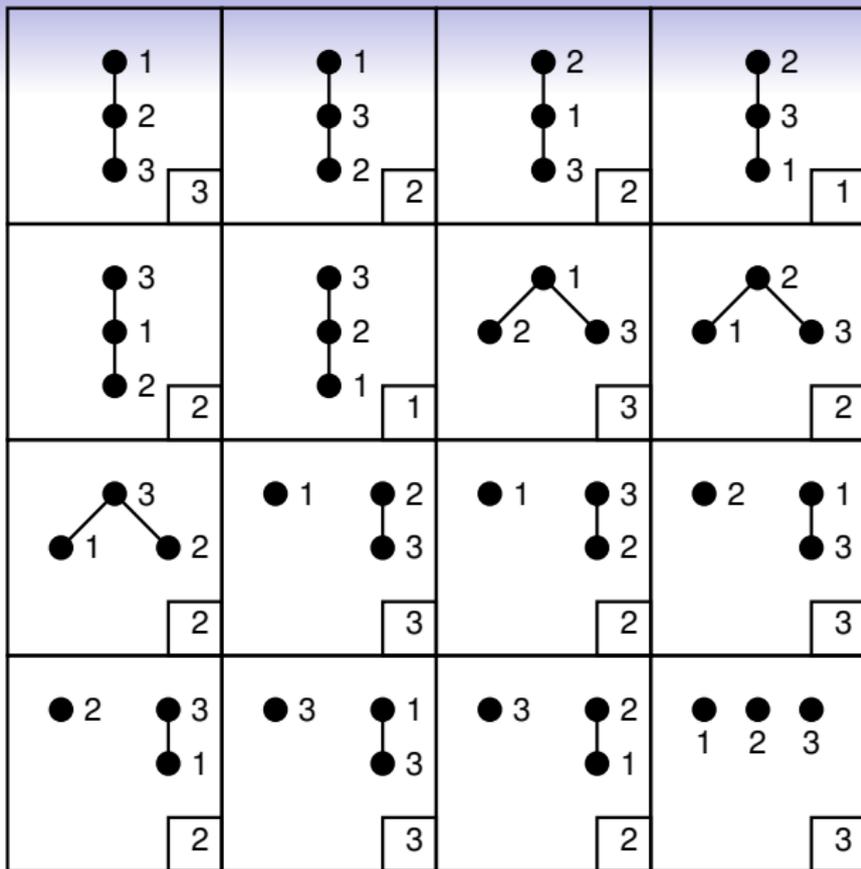


Figure: lead on forests on [3]

Definition (Lucky in Parking Functions)

- 1 $c = \text{lucky}$ in $P \Leftrightarrow \text{jump}(P; c) = 0$.
- 2 $\text{lucky}(P) =$ the # of luckys in P

| | | | |
|-----------|-----------|-----------|-----------|
| (1, 1, 1) | (1, 1, 2) | (1, 1, 3) | (1, 2, 1) |
| 1 | 1 | 2 | 2 |
| (1, 2, 2) | (1, 2, 3) | (1, 3, 1) | (1, 3, 2) |
| 2 | 3 | 2 | 3 |
| (2, 1, 1) | (2, 1, 2) | (2, 1, 3) | (2, 2, 1) |
| 2 | 2 | 3 | 2 |
| (2, 3, 1) | (3, 1, 1) | (3, 1, 2) | (3, 2, 1) |
| 3 | 2 | 3 | 3 |

Figure: lucky on Parking Functions with length 3

Theorem (Gessel-Seo 2004)

$$\sum_{F \in F_n} u^{\text{lead } F} = \sum_{P \in PF_n} u^{\text{lucky } P} \quad (2)$$

NOTE. They do not have a direct proof of (2). They found each of two GFs for lead and lucky, that are neither bijective.

$$\begin{array}{ccc}
 & u \prod_{i=1}^{n-1} (i + (n - i + 1)u) & \\
 \text{GS 2004} \swarrow & & \swarrow \text{GS 2004} \\
 \sum_{F \in F_n} u^{\text{lead } F} & \stackrel{\text{SS 2007}}{=} & \sum_{P \in PF_n} u^{\text{lucky } P} \\
 & \stackrel{?}{=} &
 \end{array}$$

Objective

We'll construct the **nonrecursive** bijection between forests and parking functions such that

$$\begin{aligned}\varphi & : F_n \rightarrow PF_n \\ F & \mapsto P = \varphi(F) \\ \text{inv}(F) & = \text{jump}(P) \\ \text{lead}(F) & = \text{lucky}(P)\end{aligned}$$

Theorem (S. 2008)

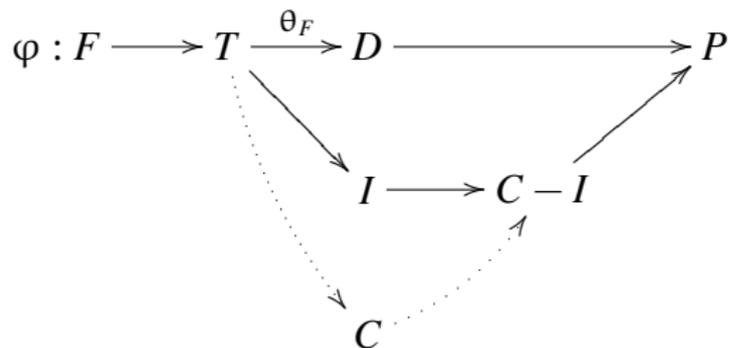
We have

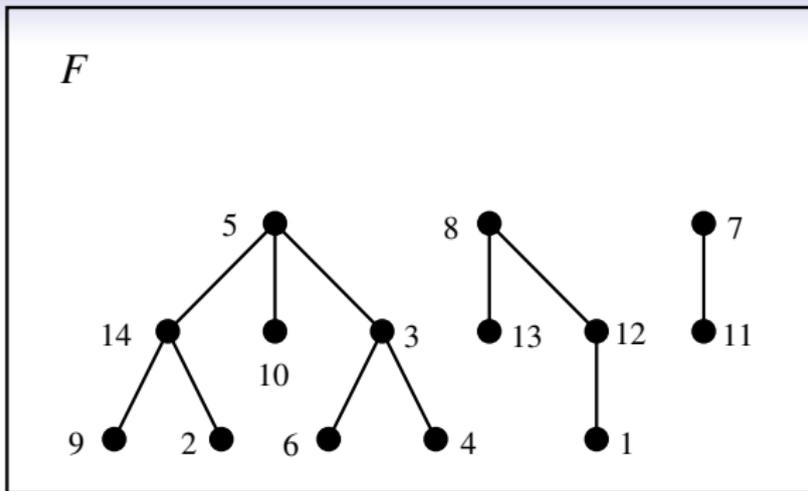
$$\sum_{F \in F_n} q^{\text{inv}(F)} u^{\text{lead}(F)} = \sum_{P \in PF_n} q^{\text{jump}(P)} u^{\text{lucky}(P)}.$$

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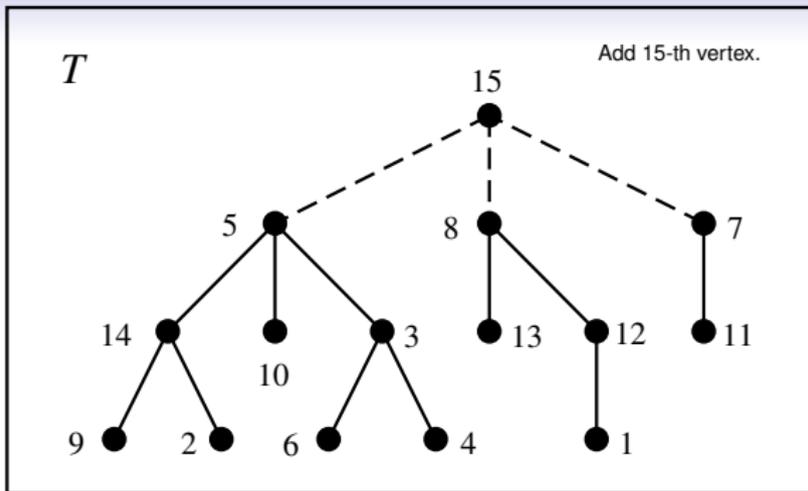
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Diagram of φ

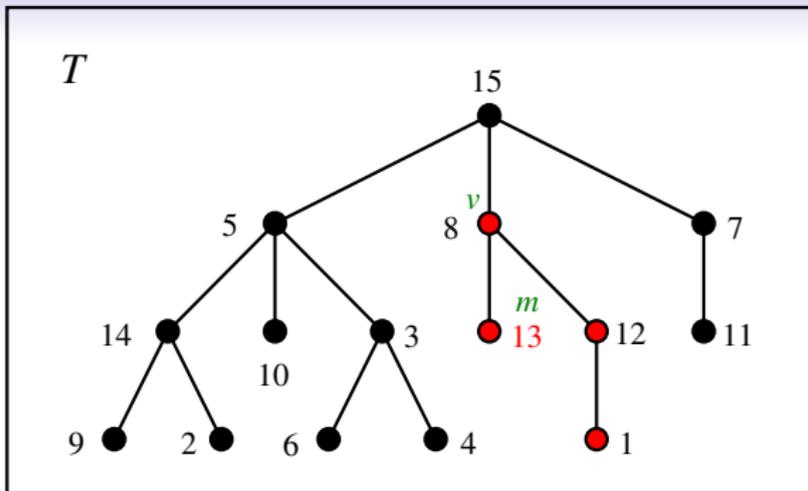




We consider $F \in F_{14}$ for example. Of course, F is drawn in the method we decide.



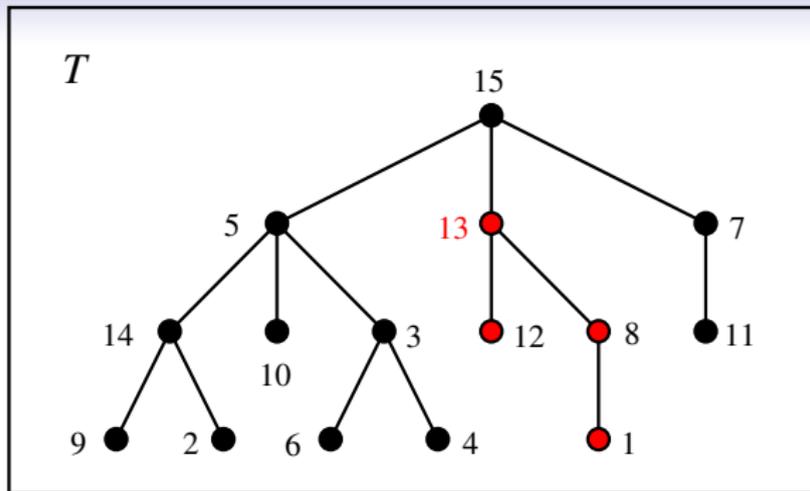
Add the vertex 15 at the top and change the forest F to the tree T .



for all $v \in V$ do

- 1 find the maximum label m on descendants of v .
- 2 label m on v .
- 3 rearrange the other labels in descendants of v by order-preserving.

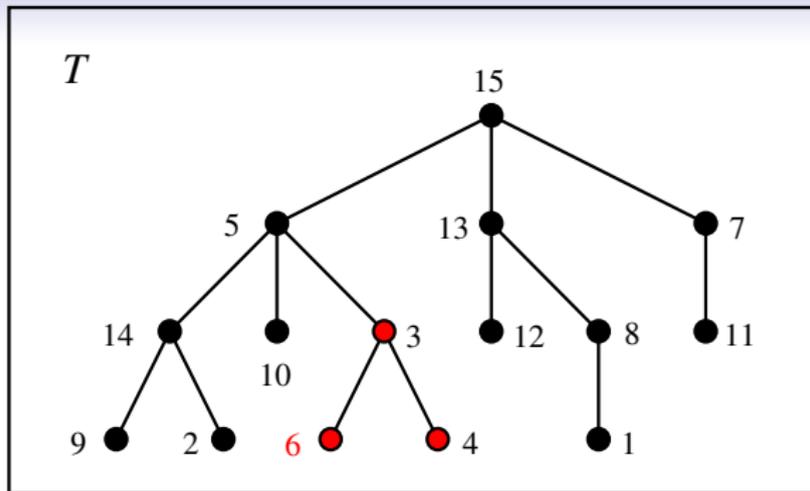
end do



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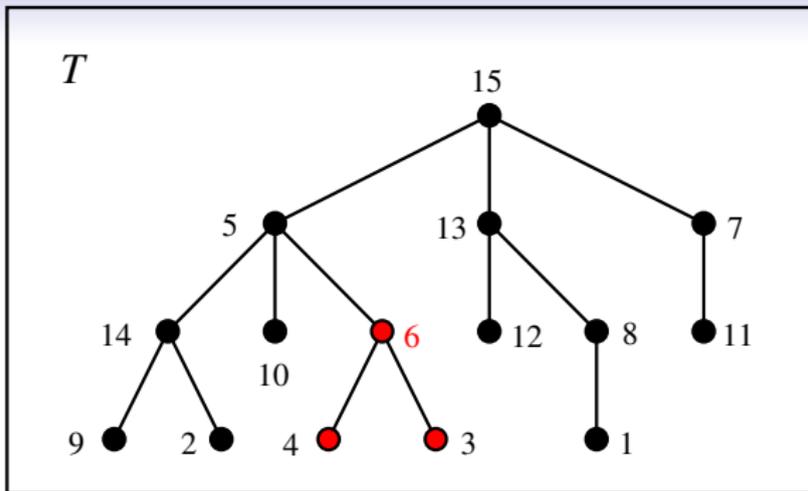
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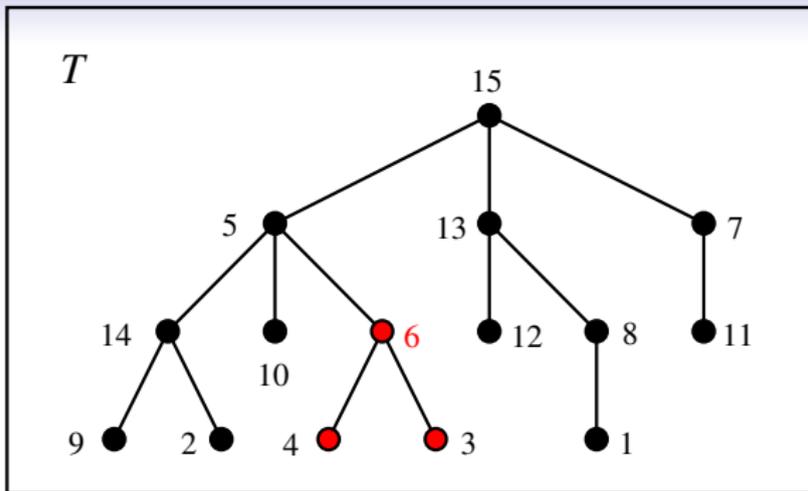
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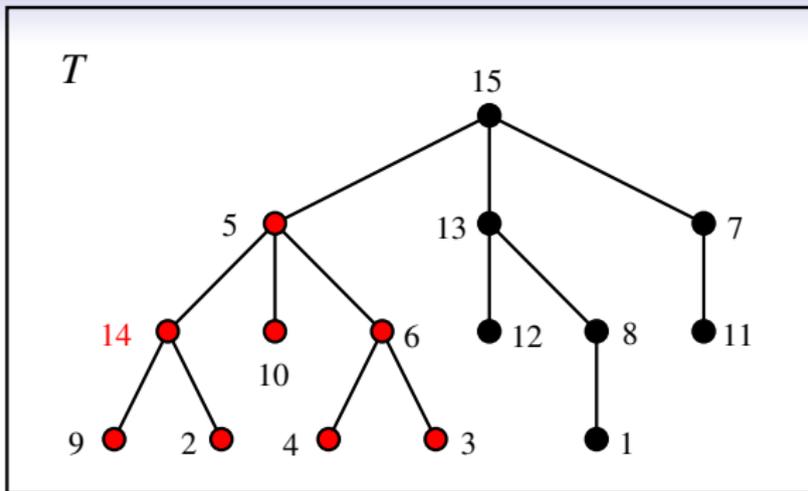
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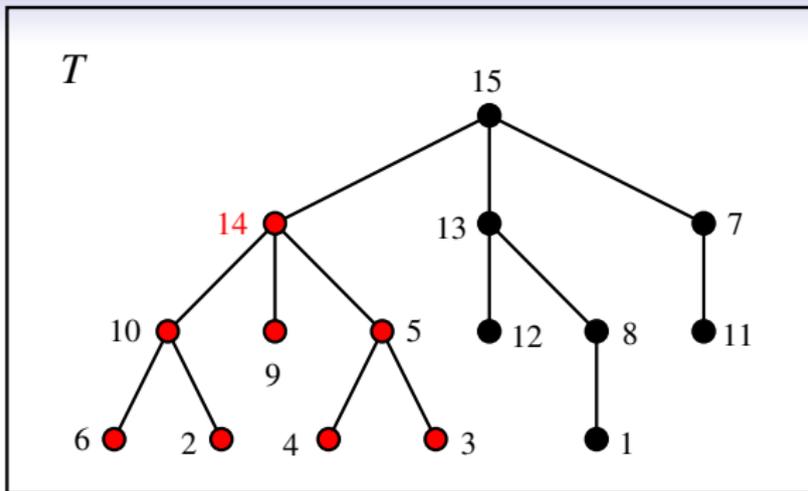
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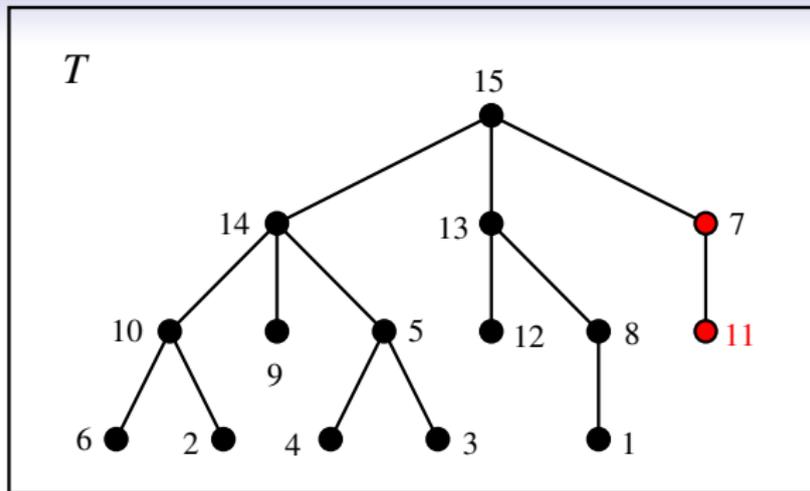
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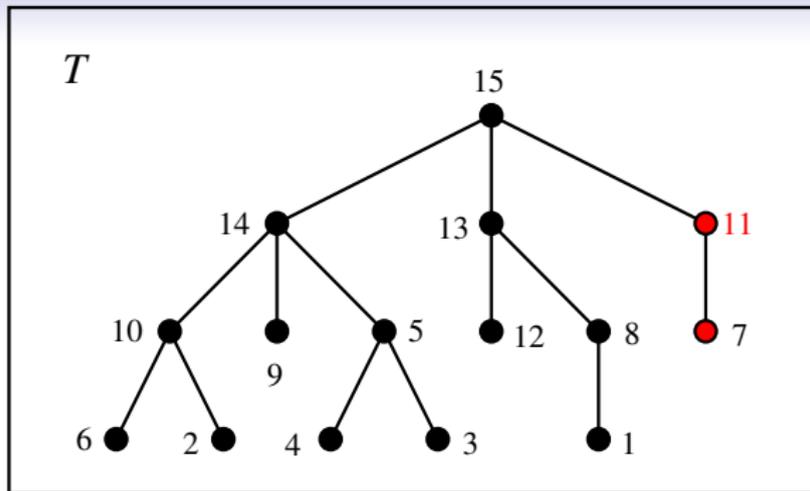
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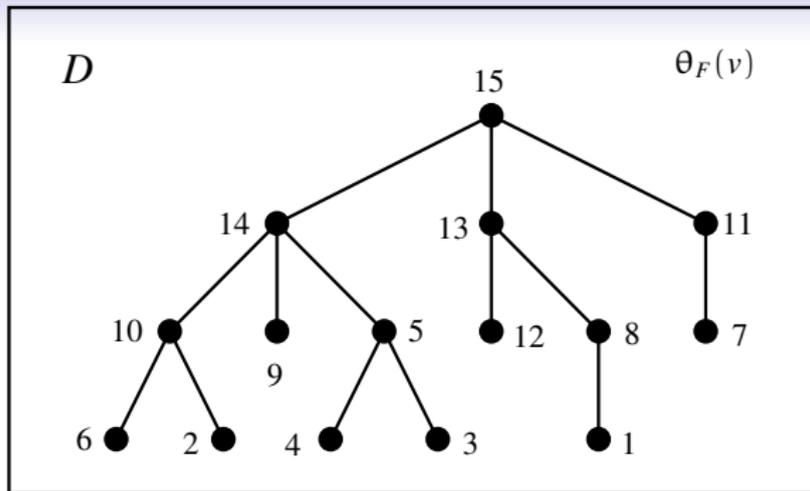
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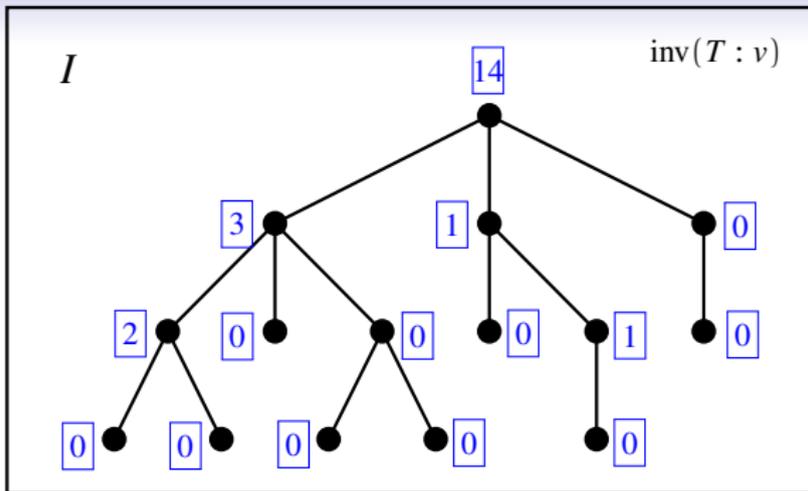
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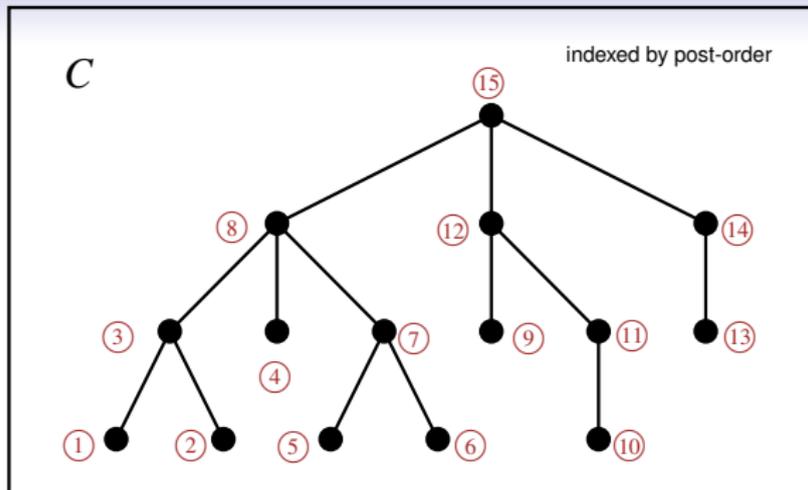
end do



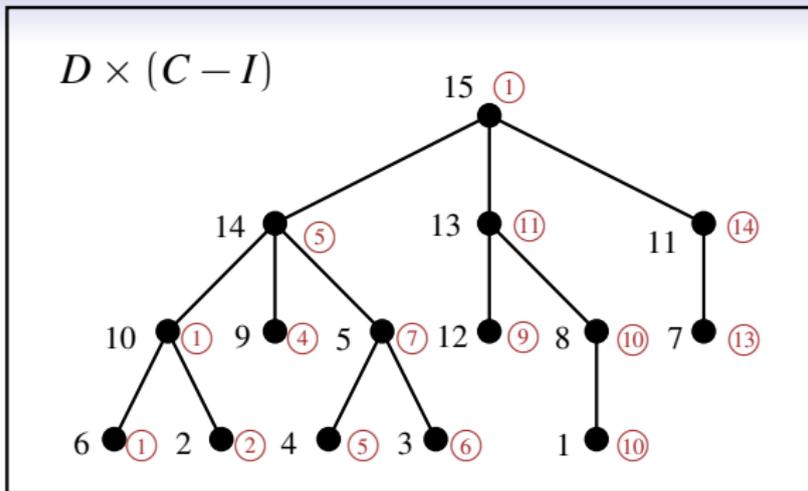
The decreasing tree is made after the process for every vertex. But we cannot remake the original tree T from only tree D . So, we need another tree induced from the unused information of T .



Label $\text{inv}(T : v)$ on vertex v . In order to distinguish it from other labels, we use the box (or blue color). And then, the trees D and I can produce the original tree T .



Label the vertices indexed by post-order which is indicated by circle (or brown color). The tree C is determined by only the *underlying graph*, that is, its tree structure. This is the reason why we define the method we draw the tree.



The plain labels are induced by D .

The circled labels are induced by C subtracted by I .

And then, we delete the tree structure and sort by plain #.

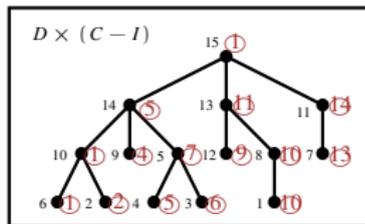
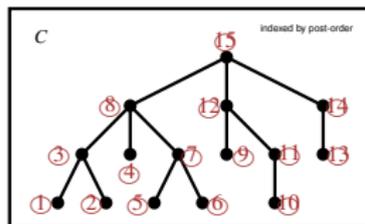
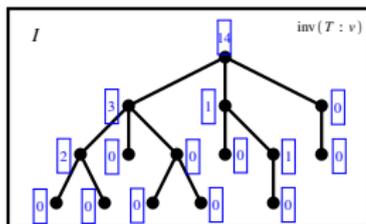
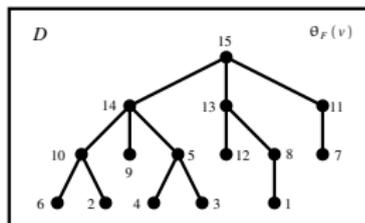
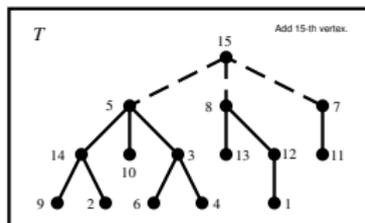
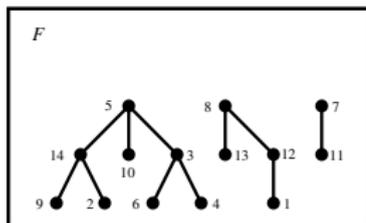
| | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|----|----|----|----|----|--|----|
| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | | 15 |
| ⑩ | ② | ⑥ | ⑤ | ⑦ | ① | ⑬ | ⑩ | ④ | ① | ⑭ | ⑨ | ⑪ | ⑤ | | ① |

Below the plain label 15, there is always circle label ①. So, we can omit it, and then second row (circle label) becomes a *parking function* P of length 14.

$$P = \text{⑩②⑥⑤⑦①⑬⑩④①⑭⑨⑪⑤}$$

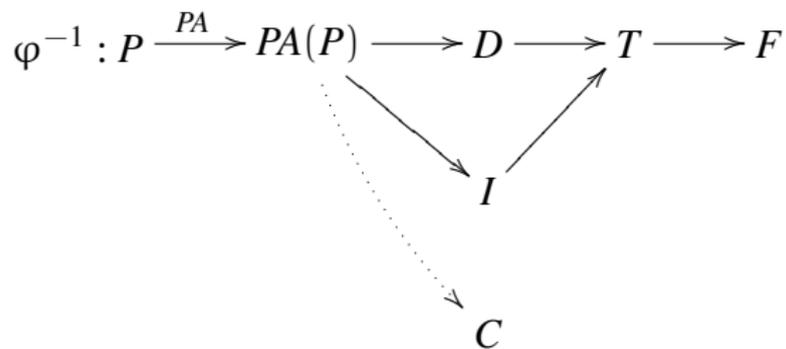
Because all labels of C are distinct in worst case which means every labels of I is all zero. Note that every permutation is a parking function.

Summary of the map φ



$$P = \begin{array}{cccccccccccccccc|c} 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 & 13 & 14 & & 15 \\ \hline 10 & 2 & 6 & 5 & 7 & 1 & 13 & 10 & 4 & 1 & 14 & 9 & 11 & 5 & & 1 \end{array}$$

Diagram of φ^{-1}



Inverse Map of φ

Let $P = (10, 2, 6, 5, 7, 1, 13, 10, 4, 1, 14, 9, 11, 5)$

After adding the (1) at the end, 15 cars is parked as following by the parking algorithm.

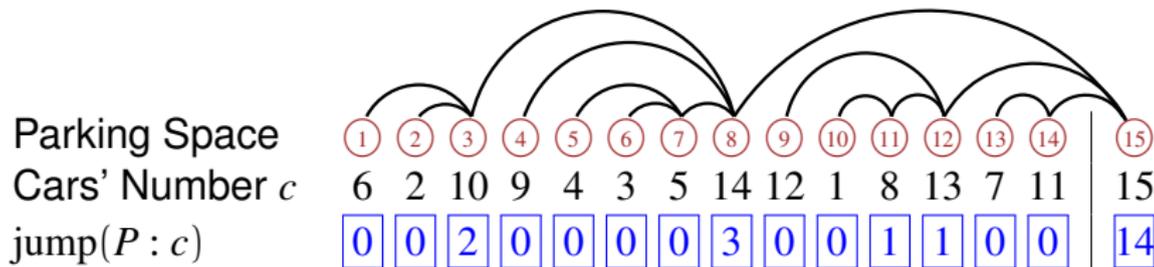
| | | | | | | | | | | | | | | | |
|----------------------|---|---|----|---|---|---|---|----|----|---|---|----|---|----|----|
| Parking Space | ① | ② | ③ | ④ | ⑤ | ⑥ | ⑦ | ⑧ | ⑨ | ⑩ | ⑪ | ⑫ | ⑬ | ⑭ | ⑮ |
| Cars' Number c | 6 | 2 | 10 | 9 | 4 | 3 | 5 | 14 | 12 | 1 | 8 | 13 | 7 | 11 | 15 |
| $\text{jump}(P : c)$ | 0 | 0 | 2 | 0 | 0 | 0 | 0 | 3 | 0 | 0 | 1 | 1 | 0 | 0 | 14 |

We draw a edge between car c and the closest car on its right which is larger than c . If we consider 15 as a root, we can rebuild the tree structure and find trees C , D and I .

Inverse Map of φ

Let $P = (10, 2, 6, 5, 7, 1, 13, 10, 4, 1, 14, 9, 11, 5)$

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i -Leader in Forests, i -Lucky in PFs

Definition (i -Leader in Forests)

- 1 $v = i$ -leader in $F \Leftrightarrow \text{inv}(F; v) = i$
- 2 $\text{lead}_i(F) =$ the # of i -leaders in F

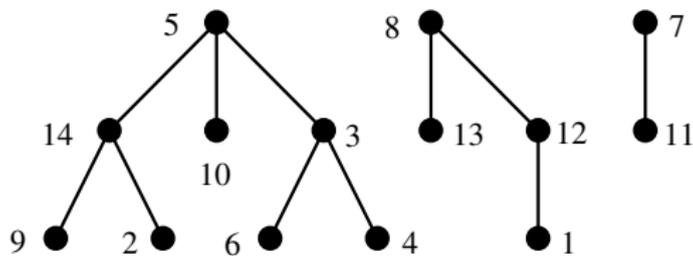
Definition (i -Lucky in Parking Functions)

- 1 $c = i$ -lucky in $P \Leftrightarrow \text{jump}(P; c) = i$
- 2 $\text{lucky}_i(P) =$ the # of i -luckys in P

We define

$$\begin{aligned} \mathbf{inv}(F) &= (\text{lead}_0(F), \text{lead}_1(F), \dots, \text{lead}_n(F)) \\ &= \text{type of inversions of } F \end{aligned}$$

$$\begin{aligned} \mathbf{jump}(P) &= (\text{lucky}_0(P), \text{lucky}_1(P), \dots, \text{lucky}_n(P)) \\ &= \text{type of jumps of } P \end{aligned}$$



$$\mathbf{inv}(F) = (10, 2, 1, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0)$$

$$\Downarrow \varphi$$

$$(10, 2, 6, 5, 7, 1, 13, 10, 4, 1, 14, 9, 11, 5)$$

$$\mathbf{jump}(P) = (10, 2, 1, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0)$$

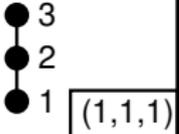
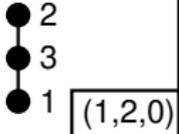
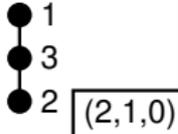
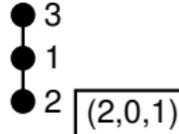
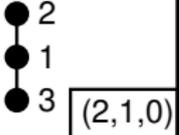
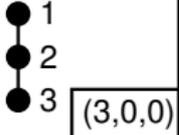
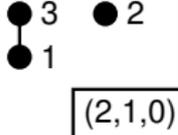
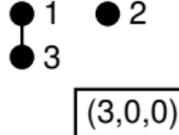
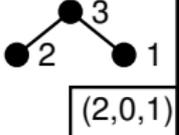
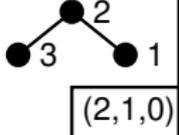
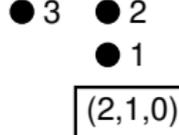
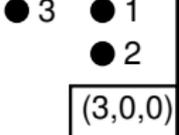
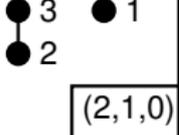
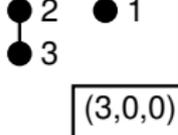
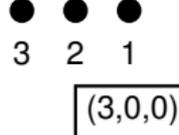
Theorem (S. 2008)

We have

$$\sum_{F \in F_n} \mathbf{q}^{\text{inv}(F)} = \sum_{P \in PF_n} \mathbf{q}^{\text{jump}(P)}$$

where $\mathbf{q}^{\mathbf{v}} = q_0^{v_0} q_1^{v_1} \cdots q_n^{v_n}$.

► [here](#)

| | | | |
|--|--|--|---|
| <p>(1, 1, 1)</p>  | <p>(1, 1, 2)</p>  | <p>(1, 1, 3)</p>  | <p>(1, 2, 1)</p>  |
| <p>(1, 2, 2)</p>  | <p>(1, 2, 3)</p>  | <p>(1, 3, 1)</p>  | <p>(1, 3, 2)</p>  |
| <p>(2, 1, 1)</p>  | <p>(2, 1, 2)</p>  | <p>(2, 1, 3)</p>  | <p>(2, 2, 1)</p>  |
| <p>(2, 3, 1)</p>  | <p>(3, 1, 1)</p>  | <p>(3, 1, 2)</p>  | <p>(3, 2, 1)</p>  |

Question

How many forests with a given type of inversions are there?

There are $n!$ forests with type $(n, 0, \dots, 0)$.

There is only one forest with $(1, \dots, 1)$.

But I know nothing about general cases yet.

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Thank you for listening!

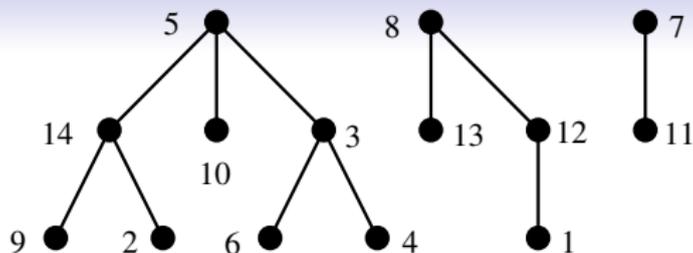
`hshin@math.univ-lyon1.fr`

Parking Algorithm

$PA(p_1, \dots, p_n)$

- $E_1 = \{1, \dots, n\}$
- for $i = 1, \dots, n$ do
 - $q_i = \min(E_1 \setminus \{p_i, \dots, n\})$
 - $E_{i+1} = E_i \setminus \{q_i\}$
- end do
- return (q_1, \dots, q_n)

◀ here



$$\text{tree}(F) = 3$$



(10, 2, 6, 5, 7, 1, 13, 10, 4, 1, 14, 9, 11, 5)

| | | | | | | | | | | | | | | |
|------------------|---|---|----|---|---|---|---|----|----|---|---|----|---|----|
| Parking Space | ① | ② | ③ | ④ | ⑤ | ⑥ | ⑦ | ⑧ | ⑨ | ⑩ | ⑪ | ⑫ | ⑬ | ⑭ |
| Cars' Number c | 6 | 2 | 10 | 9 | 4 | 3 | 5 | 14 | 12 | 1 | 8 | 13 | 7 | 11 |

$$\text{critical}(P) = 3$$

Theorem (S. 2008)

Moreover, we have

$$\sum_{F \in \mathcal{F}_n} \mathbf{q}^{\text{inv}(F)} c^{\text{tree}(F)} = \sum_{P \in \mathcal{P}\mathcal{F}_n} \mathbf{q}^{\text{jump}(P)} c^{\text{critical}(P)}$$

where $\mathbf{q}^{\mathbf{v}} = q_0^{v_0} q_1^{v_1} \cdots q_n^{v_n}$.

[◀ here](#)