

A complexity theorem for the Novelli–Pak–Stoyanovskii-algorithm

Robin Sulzgruber
joint work with Christoph Neumann

Universität Wien

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Outline

- 1 The Novelli–Pak–Stoyanovskii-algorithm
- 2 Complexity
- 3 The theorem

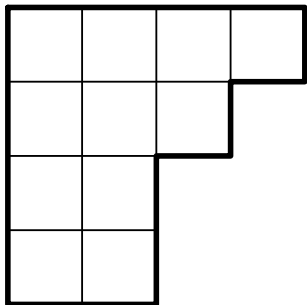
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2 Complexity

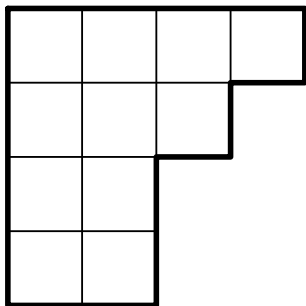
3 The theorem

Partitions



The partition $4 + 3 + 2 + 2$.

Partitions



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A partition λ of n is a left-justified array of cells with λ_i cells in the i -th row, where $\lambda_1 \geq \lambda_2 \geq \dots$ and $\sum \lambda_i = n$.

Tabloids

| | | | |
|---|----|---|---|
| 2 | 9 | 5 | 6 |
| 1 | 4 | 3 | |
| 7 | 11 | | |
| 8 | 10 | | |

A tabloid.

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Let $T(\lambda)$ denote the set of all tabloids of shape λ .

Standard Young tableaux

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A tabloid $T \in T(\lambda)$ is called standard Young tableau if T increases from left to right and from top to bottom.

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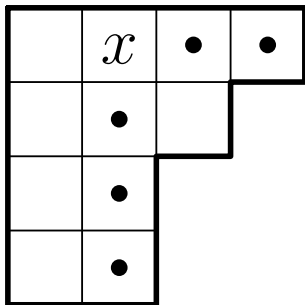
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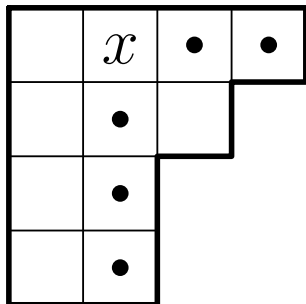
Let $\text{SYT}(\lambda)$ denote the set of all standard young tableaux of shape lambda.

Hooks



We have $h_\lambda(x) = 6$.

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The hook-length $h_\lambda(x)$ of a cell $x \in \lambda$ is the number of cells directly below or to right of x plus one.

How many SYT are there?

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For any partition λ we have

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They show that

$$n! = \#\text{SYT}(\lambda) \cdot \prod_{x \in \lambda} h_\lambda(x)$$

by assigning to each tabloid $T \in \mathbb{T}(\lambda)$ a pair of a SYT W and a hook function H .

$$T \mapsto (W, H)$$

Ordering the cells

We define an order \prec_U on the cells of λ where $U \in \text{SYT}(\lambda)$ as

$$x \prec_U y :\Leftrightarrow U(x) < U(y).$$

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 $x = (1, 2)$.

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Let $s = T(x)$. Exchange s with the least entry among its right or bottom neighbours.

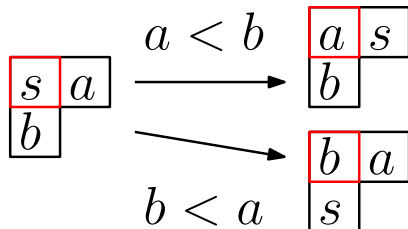
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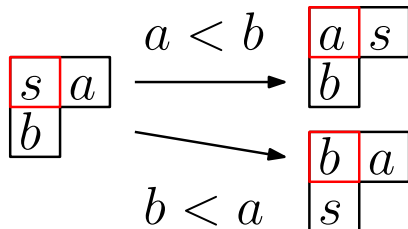
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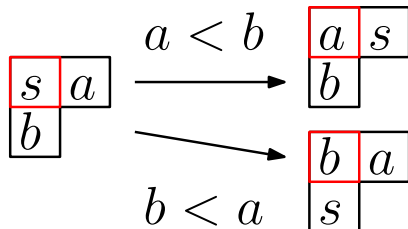
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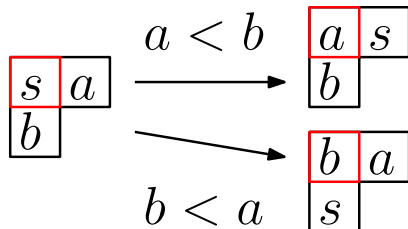
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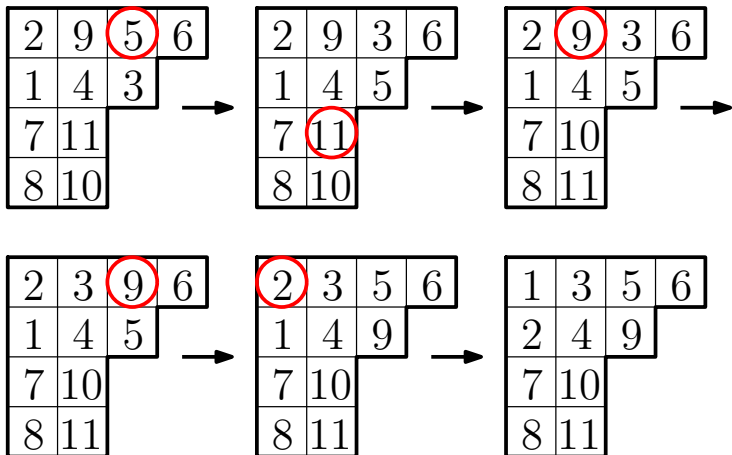
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An example



We need five exchanges.

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Complexity

Definition

Complexity

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We define the complexity of the Novelli–Pak–Stoyanovskii-algorithm with respect to the order \prec_U as the average number of exchanges, i.e.,

$$C(U) := \frac{1}{n!} \sum_{T \in \mathcal{T}(\lambda)} N_U(T),$$

where $N_U(T)$ is the number of exchanges needed to sort the tabloid T .

A conjecture

Conjecture (Krattenthaler, Müller)

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|---|---|---|---|---|---|
| 1 | 7 | 3 | 4 | 5 | 6 |
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Exchange numbers

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Let $a < b$ be entries. Denote by $m_U(a, b)$ the number of times a and b are exchanged while sorting all tabloids in $T(\lambda)$ with respect to the order \prec_U .

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Let a, b, c be entries such that $a < b < c$. Then

$$m_U(a, b) = m_U(a, c).$$

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Let a, b, c be entries such that $a < b < c$. Then

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Therefore, we just write $m_U(a)$.

Proof of lemma

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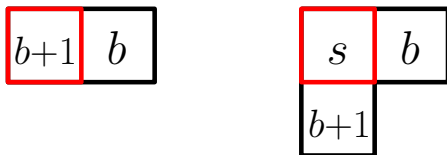
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All entries $s \notin \{b, b + 1\}$ are less than both b and $b + 1$ or larger than both. Thus, b and $b + 1$ behave similarly during sorting.



Both situations may appear only after b and $b + 1$ have dropped.

Distribution

Definition (Distribution vector)

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Let $U, W \in \text{SYT}(\lambda)$. Denote by $z_U(W)$ the number of tabloids $T \in \mathcal{T}(\lambda)$ such that sorting with respect to \prec_U transforms T into W . We call

$$\mathbf{z}_U := (z_U(W))_{W \in \text{SYT}(\lambda)}$$

the distribution vector of U .

Height

| | | | |
|---|----|---|---|
| 2 | 9 | 5 | 6 |
| 1 | 4 | 3 | |
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 $h'(3, T) = 3$.

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Let $x = (i, j)$ be a cell in λ . We define its height as
 $h'(x) := i + j - 2$.

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Let $T \in \mathcal{T}(\lambda)$ be a tabloid and a an entry, then we write
 $h'(a, T) := h'(T^{-1}(a))$.

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Suppose during sorting we exchange a and b where $a < b$. Then $h'(a, T_i) = h'(a, T_{i-1}) - 1$ and $h'(b, T_i) = h'(b, T_{i-1}) + 1$.

Initial and terminal height

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$$\sum_{W \in \text{SYT}(\lambda)} z_U(W) h'(b, W).$$

The recursion

Theorem

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Let \prec_U be given by $U \in \text{SYT}(\lambda)$. For all $1 \leq b \leq n$ we have

$$(n-b) m_U(b) = (n-1)! \sum_{x \in \lambda} h'(x) + \sum_{a=1}^{b-1} m_U(a) - \sum_{W \in \text{SYT}(\lambda)} z_U(W) h'(b, W).$$

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Corollary

The exchange numbers $m_U(a)$ only depend on the distribution vector \mathbf{z}_U .

The conjecture follows

We have $C(U) = \frac{1}{n!} \sum_{a=1}^n (n-a) m_U(a)$.

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Let \prec_U be the column-wise order. Due to the bijection of Novelli, Pak and Stoyanovskii we have

$$z_U(W) = \prod_{x \in \lambda} h_\lambda(x).$$

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Let \prec_V be the row-wise order, then $\mathbf{z}_U = \mathbf{z}_V$. The conjecture follows.

The end

Thanks for your attention!