

Compatibility fans realizing graph associahedra

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joint work with **Vincent Pilaud** (CNRS, LIX Polytechnique)

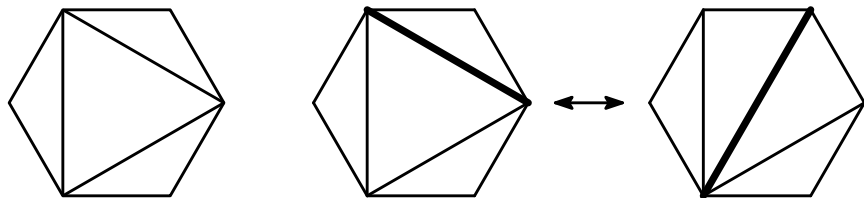
September 7th, 2015

The flip operation

Flip graph on the triangulations of the polygon:

Vertices: *triangulations*

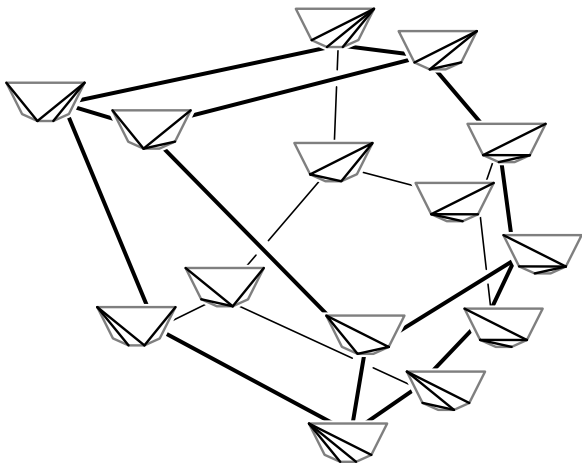
Edges: *flips*



$(n + 3)$ -gon $\Rightarrow n$ diagonals \Rightarrow the flip graph is n -regular.

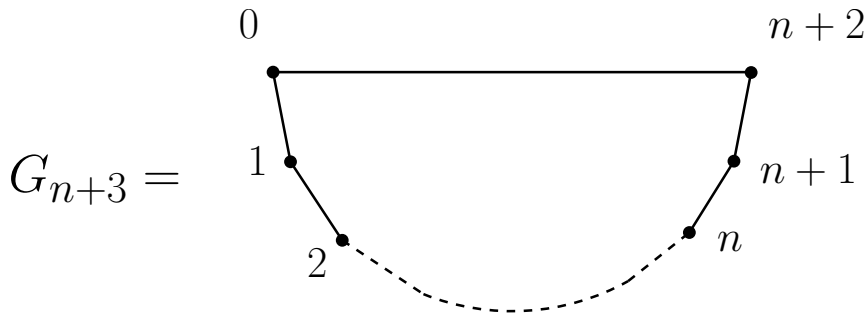
Definition

An *associahedron* is a polytope whose graph is the flip graph of triangulations of a convex polygon.



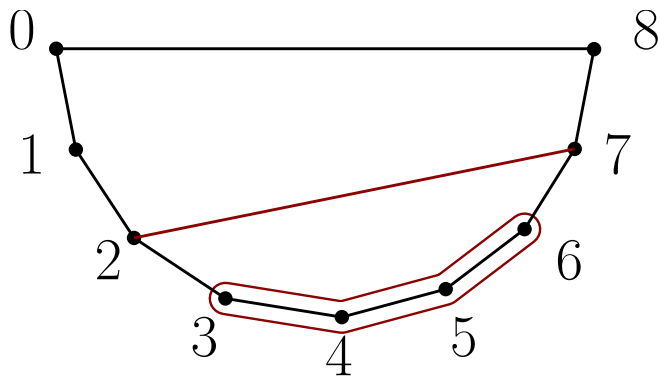
Faces \leftrightarrow dissections of the polygon

Useful configuration (Loday's)



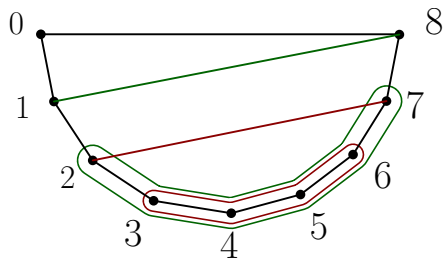
Graph point of view

$\{\text{diagonals of } G_{n+3}\} \longleftrightarrow \{\text{strict subpaths of the path } [n+1]\}$

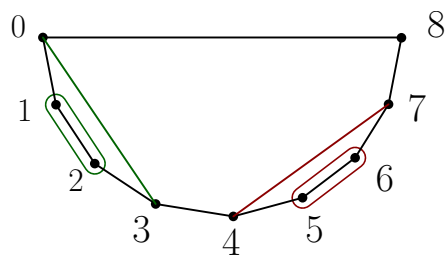


Non-crossing diagonals

Two ways to be non-crossing in Loday's configuration:



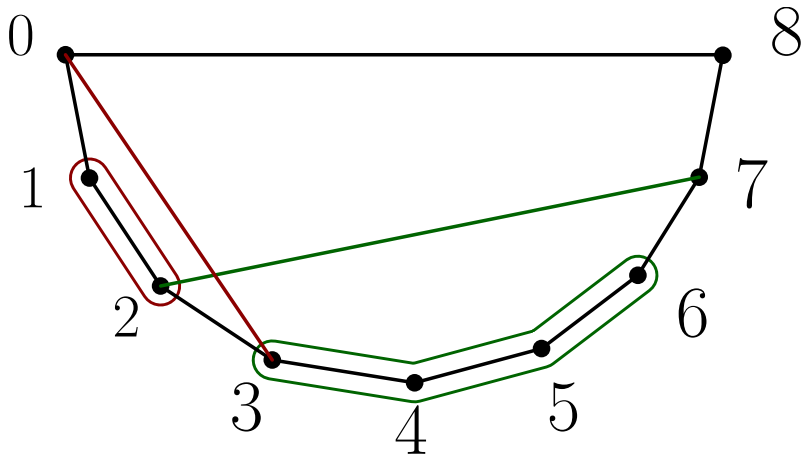
nested subpaths



non-adjacent subpaths

Caution with the second case:

The right condition is indeed *non-adjacent*, disjoint is not enough!



Now do it on graphs

$G = (V, E)$ a (connected) graph.

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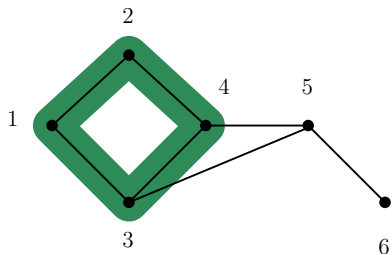
- A **tube** of G is a proper subset $t \subseteq V$ inducing a connected subgraph of G ;
- t and t' are **compatible** if they are nested or non-adjacent;

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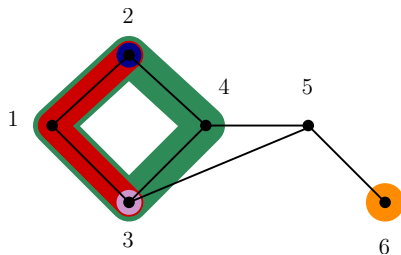
Definition

- A **tube** of G is a proper subset $t \subseteq V$ inducing a connected subgraph of G ;
- t and t' are **compatible** if they are nested or non-adjacent;
- A **tubing** on G is a set of pairwise compatible tubes of G .



A tube

(generalizes a diagonal)



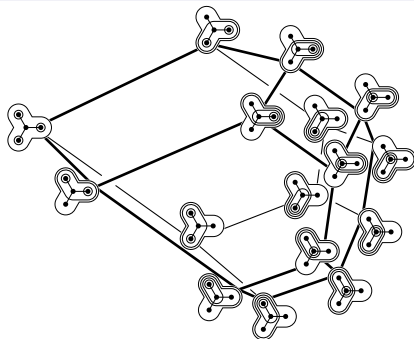
A maximal tubing

(generalizes a triangulation)

Graph associahedra

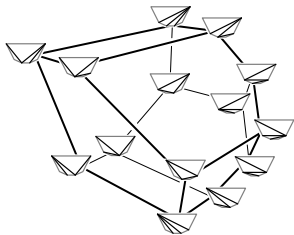
Theorem (Carr-Devadoss '06)

*There exists a polytope \mathbf{Asso}_G , the **graph associahedron** of G , realizing the complex of tubings on G .*

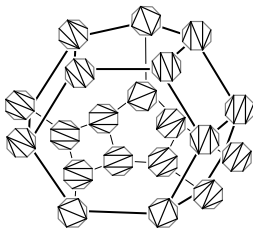


Faces \leftrightarrow tubings of G .

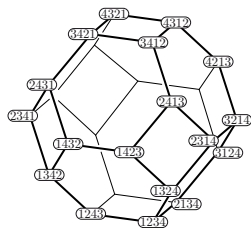
Some classical polytopes...



The associahedron

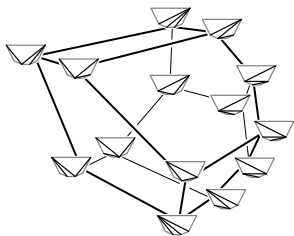


The cyclohedron

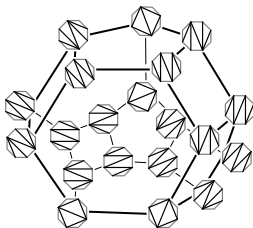


The permutahedron

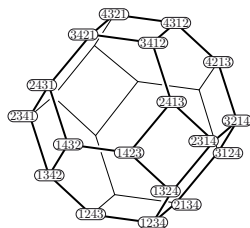
...can be seen as graph associahedra



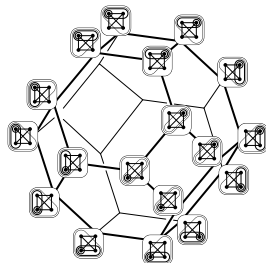
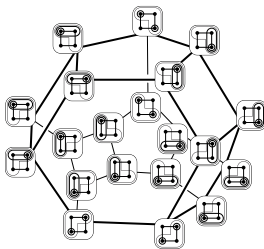
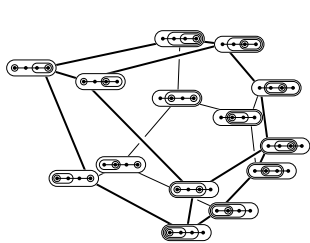
The associahedron



The cyclohedron



The permutahedron



Many different associahedra

Hohlweg-Lange [HL]: $O(2^n)$

Ceballos-Santos-Ziegler [CSZ] (Santos): $O(\text{Cat}(n))$

[HL] \cap [CSZ] = Chapoton-Fomin-Zelevinsky [CFZ] (type A): 1

Few graph associahedra

Carr-Devadoss [CD]: $1 \subset$ Postnikov [P]: 1

Volodin [Vol]: ???

Probably many, but not explicit.

Many different associahedra

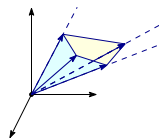
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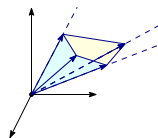
Fans

Polyhedral Cone: positive span of finitely many vectors.

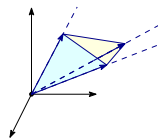


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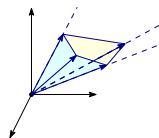


Simplicial Cone: positive span of independent vectors.

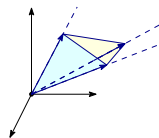


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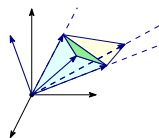
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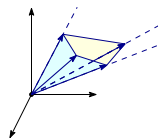


Fan = set of polyhedral cones intersecting properly.

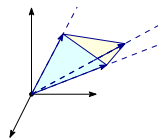


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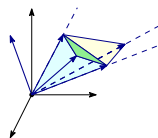
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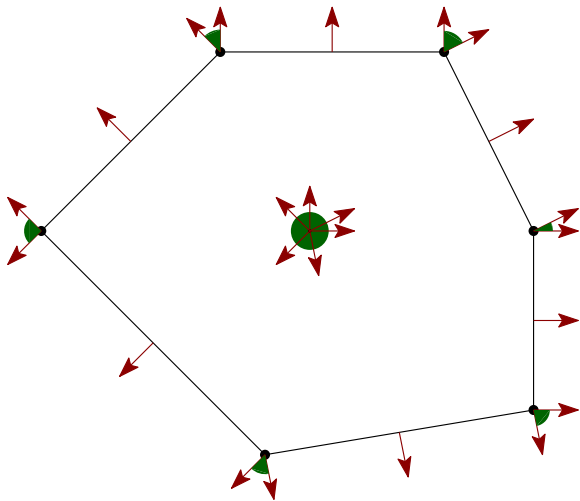


Simplicial Fan: fan whose cones all are simplicial.

Complete Fan: fan whose cones cover the whole space.

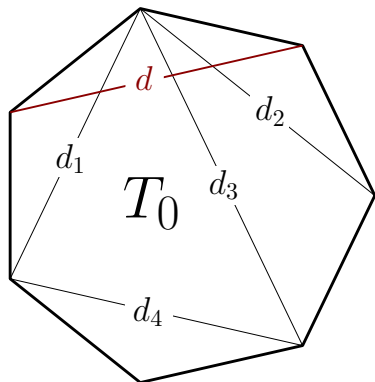
polytope \Rightarrow complete fan (*normal fan*).

simple polytope \Rightarrow complete simplicial fan.



Santos' construction for the fan

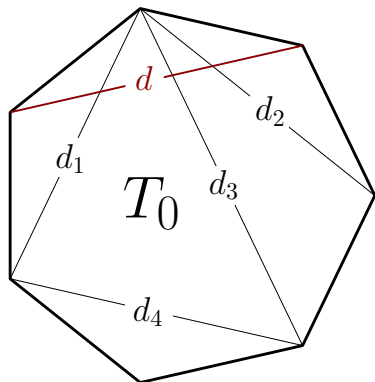
→ choose an initial triangulation T_0 of the polygon.



set $u_{d_i} = -e_i$

Santos' construction for the fan

→ choose an initial triangulation T_0 of the polygon.



set $u_{d_i} = -e_i$

→ for a diagonal $d \notin T_0$, define $u_d = (\mathbf{1}_{d \text{ crosses } d_i})_{d_i \in T_0}$.

→ for a triangulation T , define $C(T) = \text{cone}(u_d | d \in T)$.

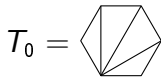
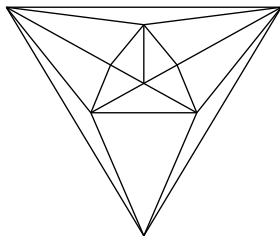
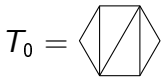
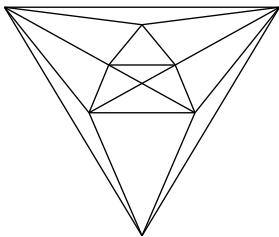
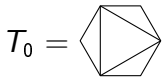
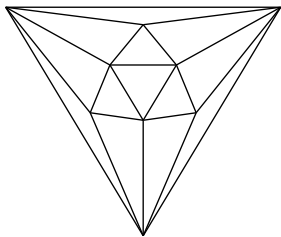
→ Define $\mathcal{F} = \{C(T) | T \text{ triangulation}\}$.

Theorem (Ceballos-Santos-Ziegler 13)

\mathcal{F} is a complete simplicial fan realizing the associahedron.

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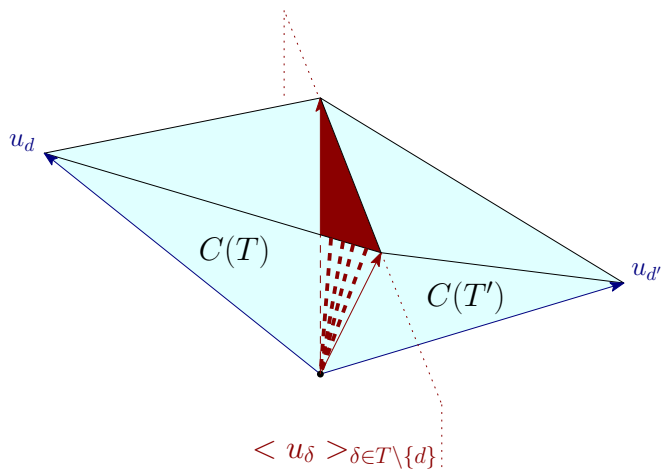


Idea of the proof

- The cone $C(T_0)$ is the negative orthant.
 - ⇒ full-dimensional and simplicial

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 - ⇒ full-dimensional and simplicial
- Local condition on flips $T \leftrightarrow T' = T \setminus \{d\} \cup \{d'\}$.



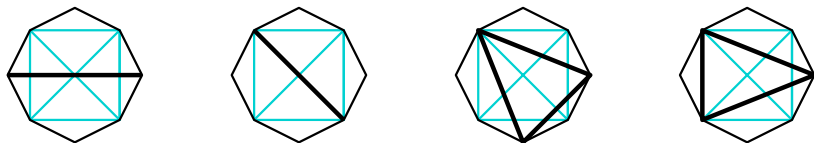
Checking local conditions

→ Formulation: $\alpha u_d + \alpha' u_{d'} + \sum_{\delta \in T \setminus \{d\}} \beta_\delta u_\delta = 0 \Rightarrow \alpha \cdot \alpha' > 0.$

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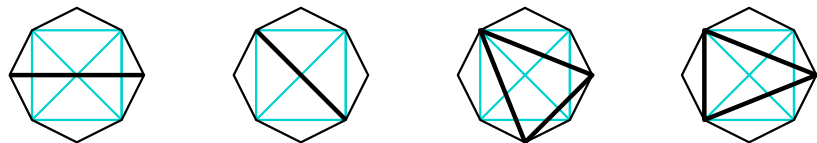
→ Reduction:



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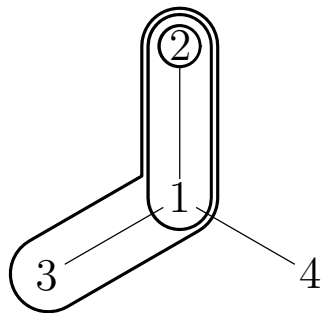
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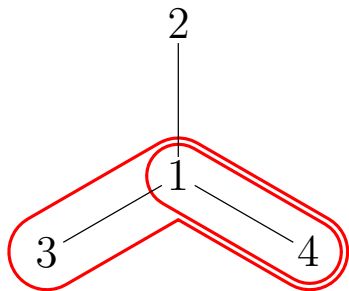
→ Finite number of linear dependences to check explicitly.



For graphs?



T_0



→ impossible to choose $-1, 0, 1$ coordinates.

The compatibility degree

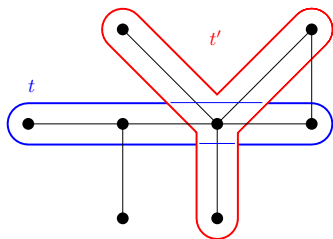
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$$(t \parallel t') = \begin{cases} -1 & \text{if } t = t', \\ \#(\text{neighbors of } t' \text{ in } t \setminus t') & \text{if } t' \not\subseteq t, \\ 0 & \text{otherwise.} \end{cases}$$

→ Counts compatibility obstructions.



$$(t \parallel t') = 2$$
$$(t' \parallel t) = 3$$

The result!

- Define $u_t = ((t \parallel t_1), \dots, (t \parallel t_n))$
- For a maximal tubing T , define $C(T) = \text{cone}(u_t | t \in T)$.
- Define $\mathcal{F}_G = \{C(T) | T \text{ triangulation}\}$.

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Theorem (M., Pilaud 15)

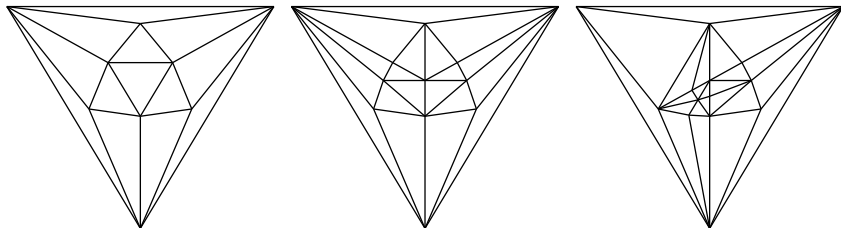
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THANK YOU FOR
YOUR PATIENT
LISTENING!