Enumeration of polyominoes with fixed perimeter defect

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Gill Barequet, Yufei Zheng (Technion – Israel Institute of Technology) A polyomino is a set of unit cells in the grid, connected via edges.



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The size of a polyomino is the number of cells.

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Only 56 numbers are known (Jensen); The nature of the generating function is not known; The limit  $\lambda := \lim A(n+1)/A(n)$  exists (Madras); Bounds:  $\lambda > 4.0025$  (Barequet, Rote, Shalah),  $\lambda < 4.6496$  (Klarner, Rivest); Conjectured:  $\lambda \approx 4.0626$ .







We denote by A(n, p) the number of polyominoes with area n and perimeter p.



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Definition:

The (perimeter) defect of P is k := 2n + 2 - p.

A *perimeter cell*: a free cell that has an occupied neighbor.

An *excess* of a perimeter cell is the number of occupied neighbors minus 1.

The *total excess* of a polyomino is the sum of excesses over all perimeter cells.





- each free cell excess contributes its excess directly,
- each edge of a spanning tree contributes 2,
- each edge closing a cycle contributes 2.



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 $\begin{array}{l} k=0 \Rightarrow \\ e=f=0 \Rightarrow \\ \text{the polyomino is a "stick"} \Rightarrow \\ \text{for } n\geq 2 \text{, we have } A(n,2n+2)=1. \end{array}$ 

Small values of k (k = 2n + 2 - p, k = e + 2f) $k = 0 \Rightarrow$  $e = f = 0 \Rightarrow$ the polyomino is a "stick"  $\Rightarrow$ for  $n \geq 2$ , we have A(n, 2n+2) = 1.  $k = 1 \Rightarrow$  $e = 1, f = 0 \Rightarrow$ the polyomino is a "one-bend path"  $\Rightarrow$ for  $n \ge 3$ , we have A(n, 2n + 1) = 4(n - 2).

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A(n, 2n) = a quadratic polynomial.







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k = 3: We have 18 patterns, not all of them are polynomial. Upon addition of 18 generating functions:

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 $A(n, 2n - 1) = -\frac{2107}{2} - \frac{21}{2}(-1)^n + \left(\frac{1947}{4} + \frac{5}{4}(-1)^n\right)n - 89n^2 + \frac{13}{2}n^3$ 

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- 2. The number of patterns is finite.
- → A "correct" notion of pattern (to ensure partition).



We expect that these will be in the same class.



What about these?







































Vertical cut C: A maximal set of consecutive columns,  $C_i, \ldots, C_j$ , such that for  $C^* = \{C_{i-1}, \ldots, C_{j+1}\}$ we have  $P \cap C^* = \{i - 1, \ldots, j + 1\} \times A$ , where  $A \neq \emptyset$  and  $a, b \in A, a \neq b \Rightarrow |a - b| \ge 3$ .













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For the proof, we define three kinds of special cells:1. Excess cells (perimeter cells with non-zero excess),2. Bend cells (cells of P with at least one pair of non-opposite occupied neighbors),

3. Leaf cells (degree 1 in the dual).

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Claim 1: If a column that intersects a polyomino does not belong to a cut, then it or at least one of its neighbors contains a special cell.




























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1. Each excess cells contributes directly to e.

2. Each bend cell contributes directly to e or to f.

3.  $|V_1| \leq |V_3| + 2|V_4| + 2$ , and each cell of degree 3 or 4 is a bend cell.

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 $\Rightarrow$  The number of columns that do not belong to a cut, is bounded.

 $\Rightarrow$  All the reduced polyominoes lie in a square of a bounded size.

 $\Rightarrow$  There is a finite number of reduced polyominoes.

 $\Rightarrow$  There is a finite number of patterns classes.

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For each fiked k, the generating function of  $(A(n, 2n + 2 - k))_{n \ge k}$  is rational.

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(The maximum perimeter is 2n(d-1) + 2. In the formula k = e + 2f, f is the "circuit rank" – the number of edges that must be removed from the graph in order to obtain a tree. It is equal to |E| - |V| + 1.)

0 x-1 $(x-1)^2$ 1  $(x-1)^3$ 2  $(x-1)^4(x+1)^2$ 3  $(x-1)^5(x+1)^3$ 4  $(x-1)^6(x+1)^4(x^2+x+1)$ 5

CP

k

