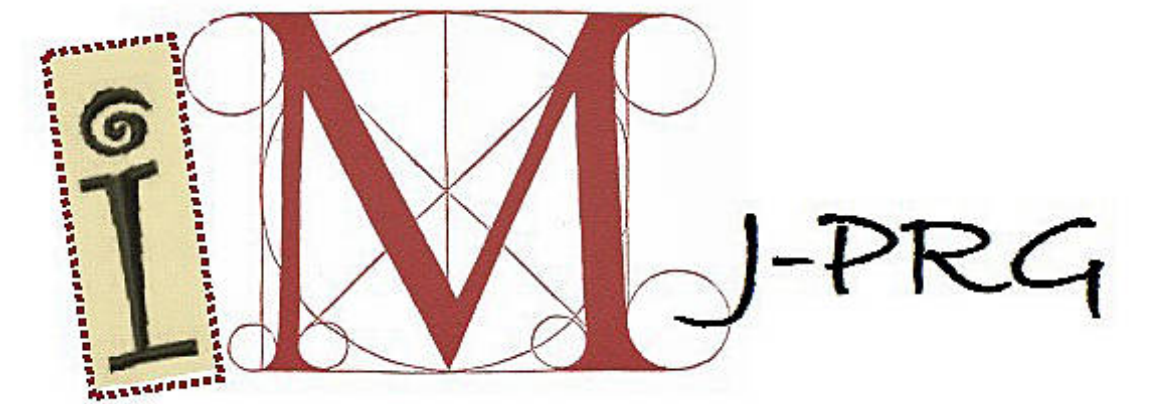


Monotone path polytopes of the hypersimplices $\Delta(n, 2)$

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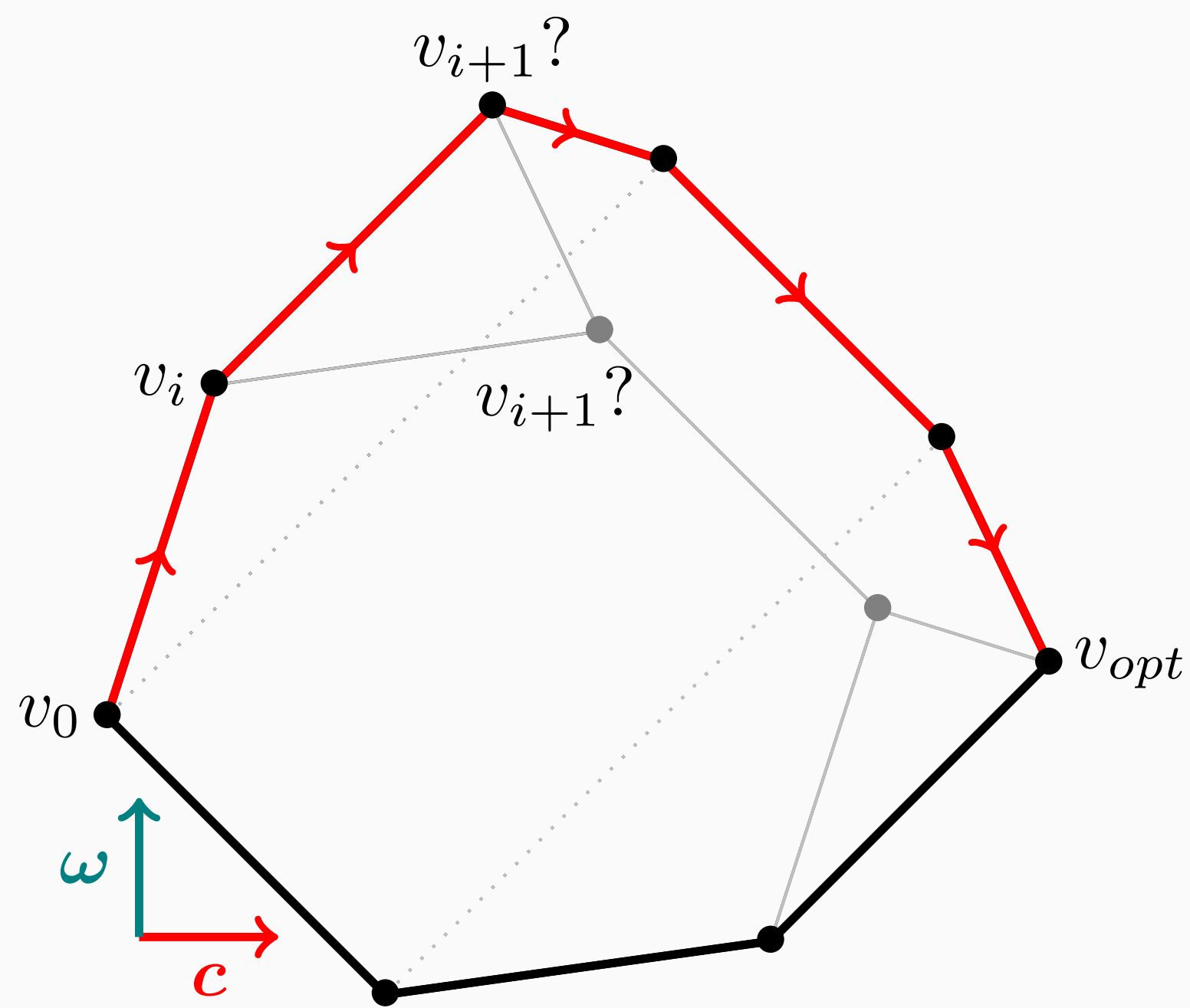
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Shadow vertex rule

Linear program (P, c) : how to choose next vertex in simplex method?



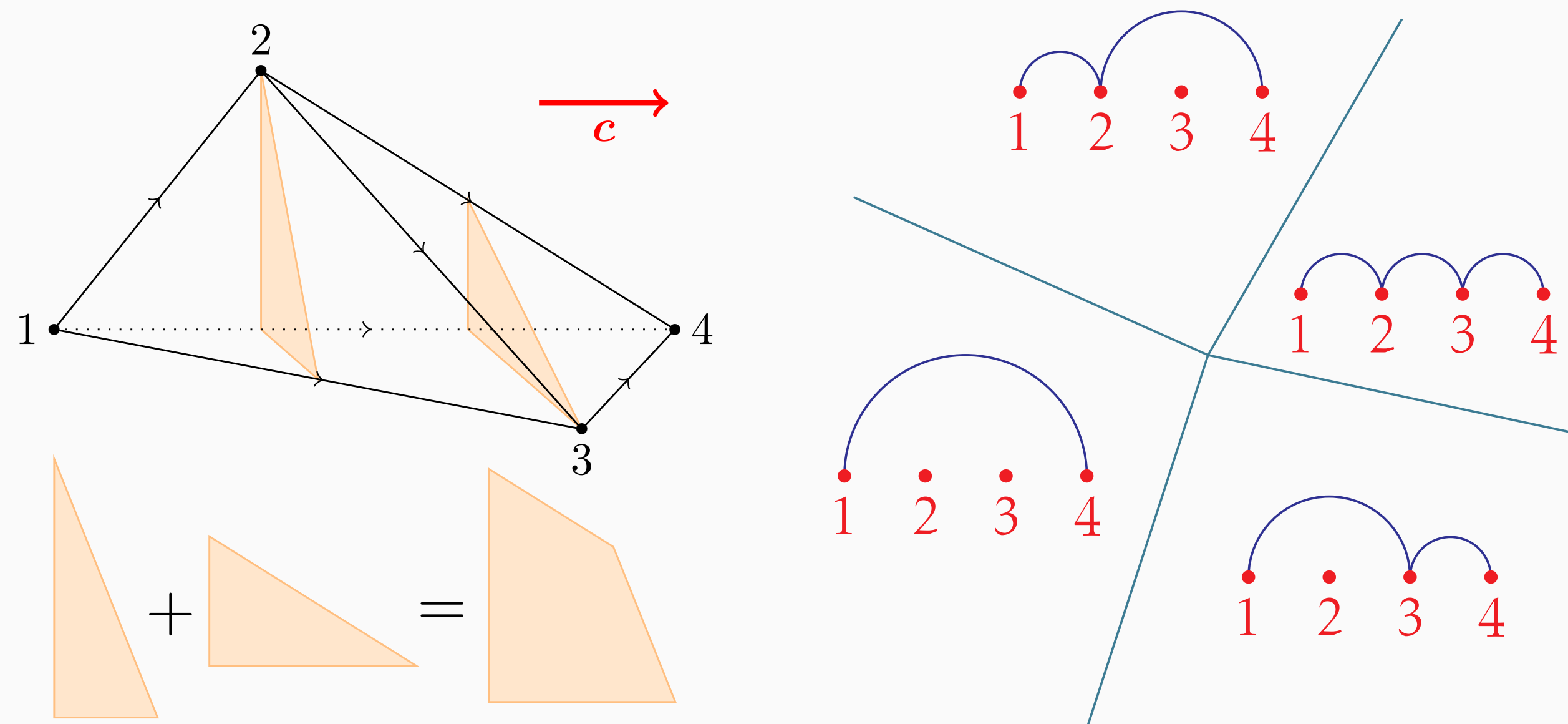
Shadow vertex: for ω , project in plane (c, ω) , take the neighbor with the best slope:

$$v_{i+1} = \operatorname{argmax} \left\{ \begin{array}{l} \langle \omega, u - v_i \rangle; \quad u \text{ improving} \\ \langle c, u - v_i \rangle; \quad \text{neighbor of } v_i \end{array} \right\}$$

Monotone path polytope of a polytope

Coherent monotone path: monotone path arising from shadow vertex rule

Monotone path fan: $\omega \sim \omega'$ iff same monotone path



Monotone path polytope $\Sigma_c(P)$: Polytope dual to monotone path fan
 \simeq Minkowski sum of section over (images of) vertices
 \simeq Fiber polytope $\Sigma_\pi(P, Q)$ for $\pi : x \mapsto \langle x, c \rangle$

Vertices of $\Sigma_c(P) \longleftrightarrow c$ -coherent monotone paths on P

Monotone path polytope of simplices

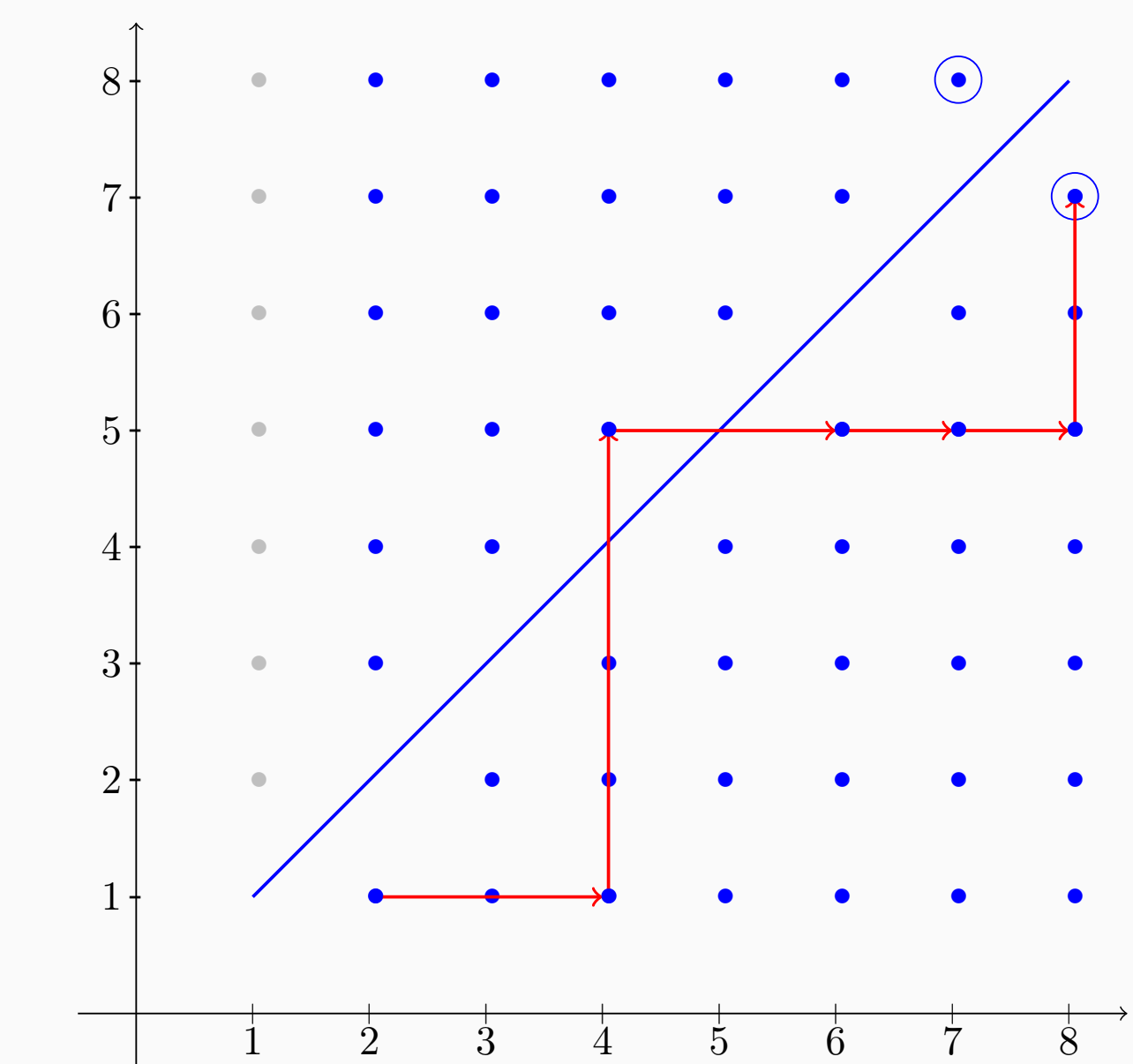
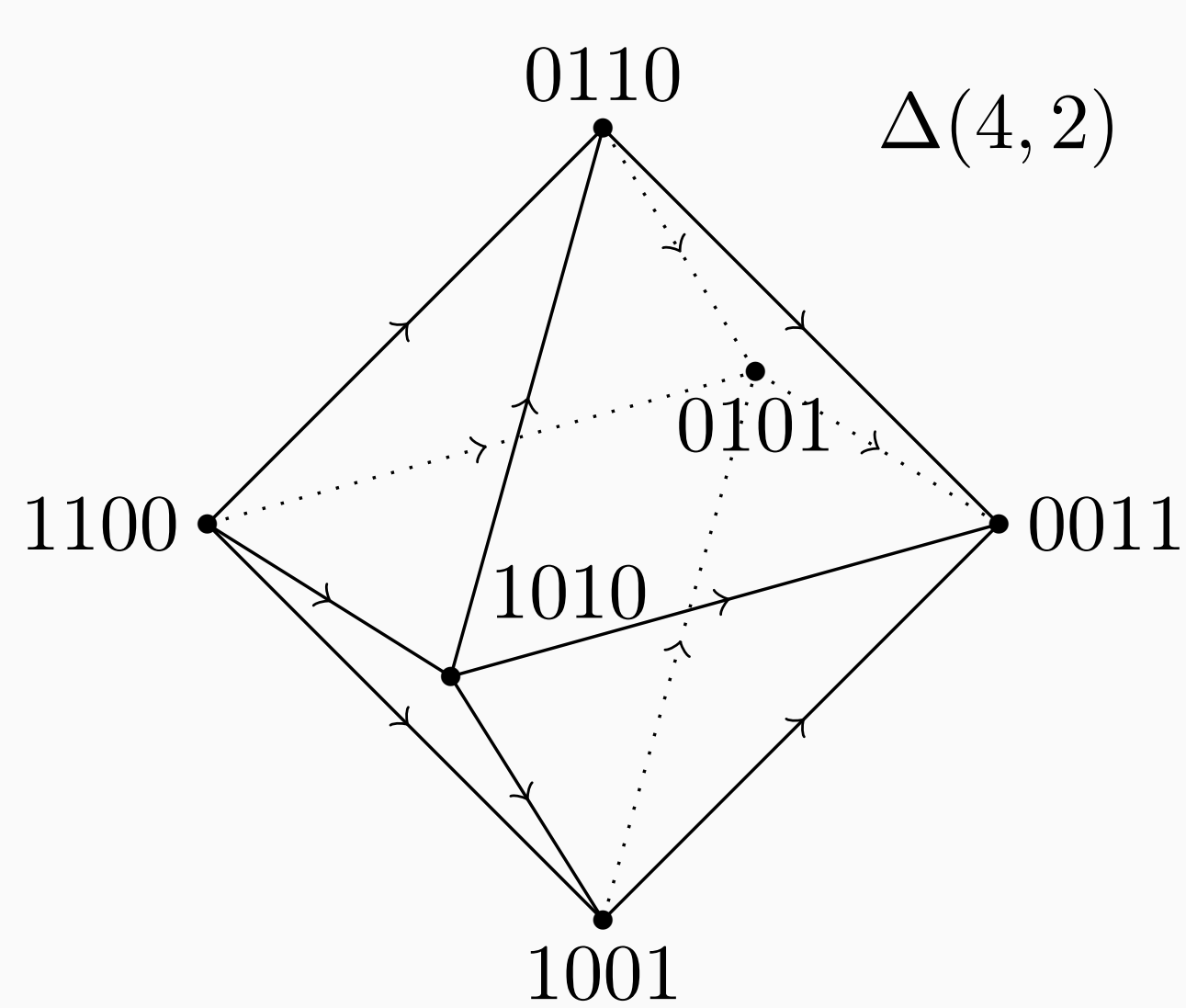
THM. [Billera, Sturmfels] for all c , $\Sigma_c(\Delta_{n+1}) \simeq \text{Cube}_{n-1}$

Hypersimplex $\Delta(n, 2)$

Hypersimplex $\Delta(n, k) = \operatorname{conv} \{v \in \{0, 1\}^n : \sum v_i = k\}$; Along $c = (1, 2, \dots, n)$

$\Delta(n, 1) \simeq \Delta(n, n-1) \simeq \Delta_n$: simplex

Here, focus on $\Delta(n, 2)$



$2 \xrightarrow{1} 4, 1 \xrightarrow{4} 5, 4 \xrightarrow{5} 6, 6 \xrightarrow{5} 7, 7 \xrightarrow{5} 8, 5 \xrightarrow{8} 7$

Monotone paths on $\Delta(n, 2) \leftrightarrow$ lattice paths on $[1, n]^2$, start at $(2, 1)$, avoid (i, i) , end at $(n, n-1)$ or $(n-1, n)$

Notation: $i \xrightarrow{a} j$ when step $(i, a) \rightarrow (j, a)$ or $(a, i) \rightarrow (a, j)$ in path

Counting coherent monotone paths on $\Delta(n, 2)$

Induction (see right): $v_n := |\text{Vertices}(\Sigma_c(\Delta(n, 2)))| = t_n + q_n + c_n$

Where $\begin{pmatrix} t_{n+1} \\ q_{n+1} \\ c_{n+1} \end{pmatrix} = M \begin{pmatrix} t_n \\ q_n \\ c_n \end{pmatrix}$, with $M = \begin{pmatrix} 1 & 2 & 2 \\ 0 & 2 & 1 \\ 2 & 0 & 2 \end{pmatrix}$, $\operatorname{Sp}(M) = \{0, 1, 4\}$, $\begin{pmatrix} t_0 \\ q_0 \\ c_0 \end{pmatrix} = \begin{pmatrix} 3 \\ 1 \\ 4 \end{pmatrix}$

THM. Vertices $\Sigma_c(\Delta(n, 2))$, i.e. coh. mon. paths: $(1 \ 1 \ 1)M^n \begin{pmatrix} 3 \\ 1 \\ 4 \end{pmatrix}$

$$v_n = \frac{1}{3} (25 \times 4^{n-4} - 1)$$

Sorting by length: $V_n(z) := \sum_{\ell} v_{n, \ell} z^\ell = T_n(z) + Q_n(z) + C_n(z)$

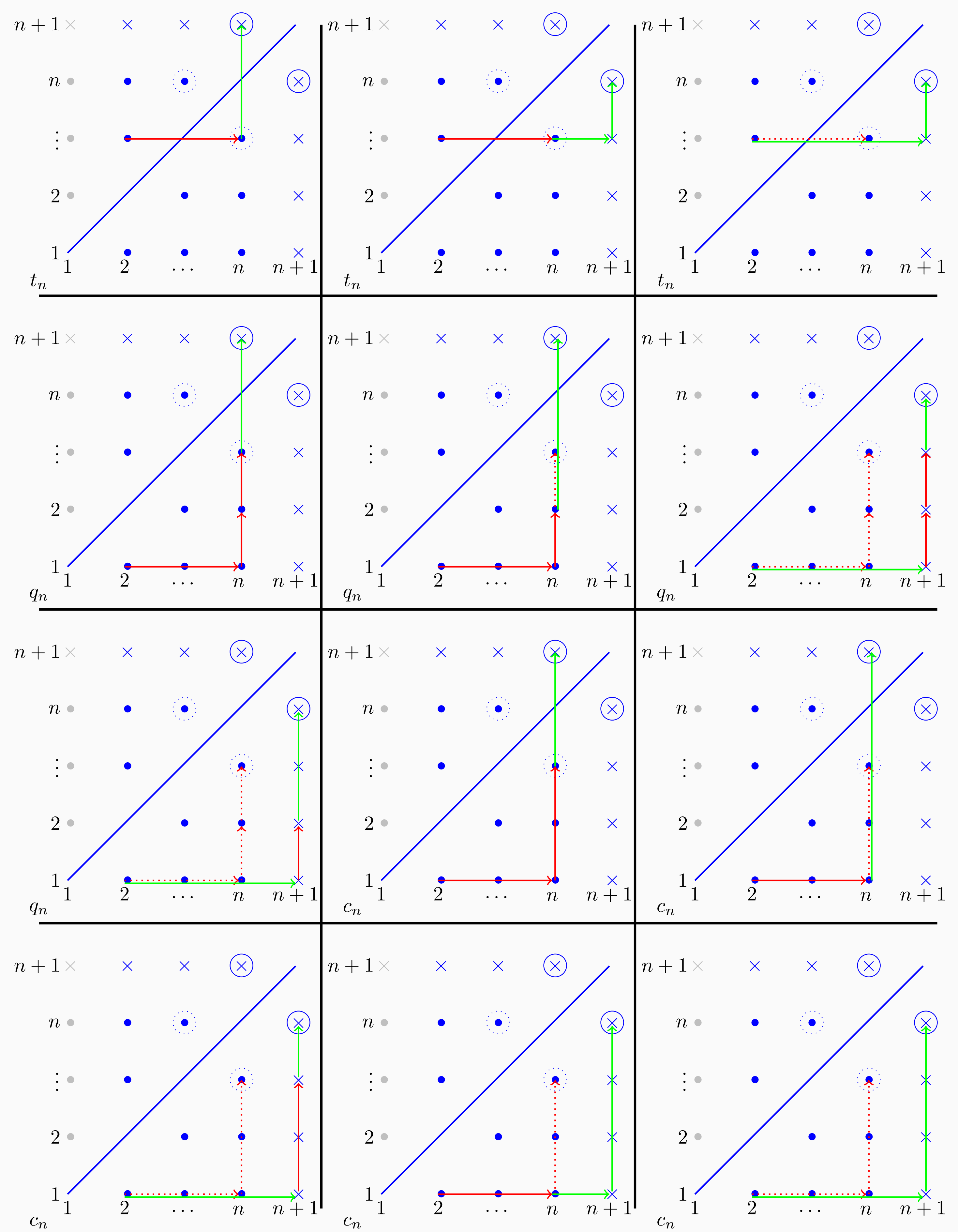
$\begin{pmatrix} T_{n+1} \\ Q_{n+1} \\ C_{n+1} \end{pmatrix} = M \begin{pmatrix} T_n \\ Q_n \\ C_n \end{pmatrix}$, with $M = \begin{pmatrix} z & 1+z & 1+z \\ 0 & 1+z & z \\ z^2+z & 0 & 1+z \end{pmatrix}$, $\begin{pmatrix} T_0 \\ Q_0 \\ C_0 \end{pmatrix} = \begin{pmatrix} z^3 + 2z^2 \\ z^3 \\ 2z^3 + 2z^2 \end{pmatrix}$

THM. $v_{n, \ell}$ is a polynomial in n of degree $\ell - 2$

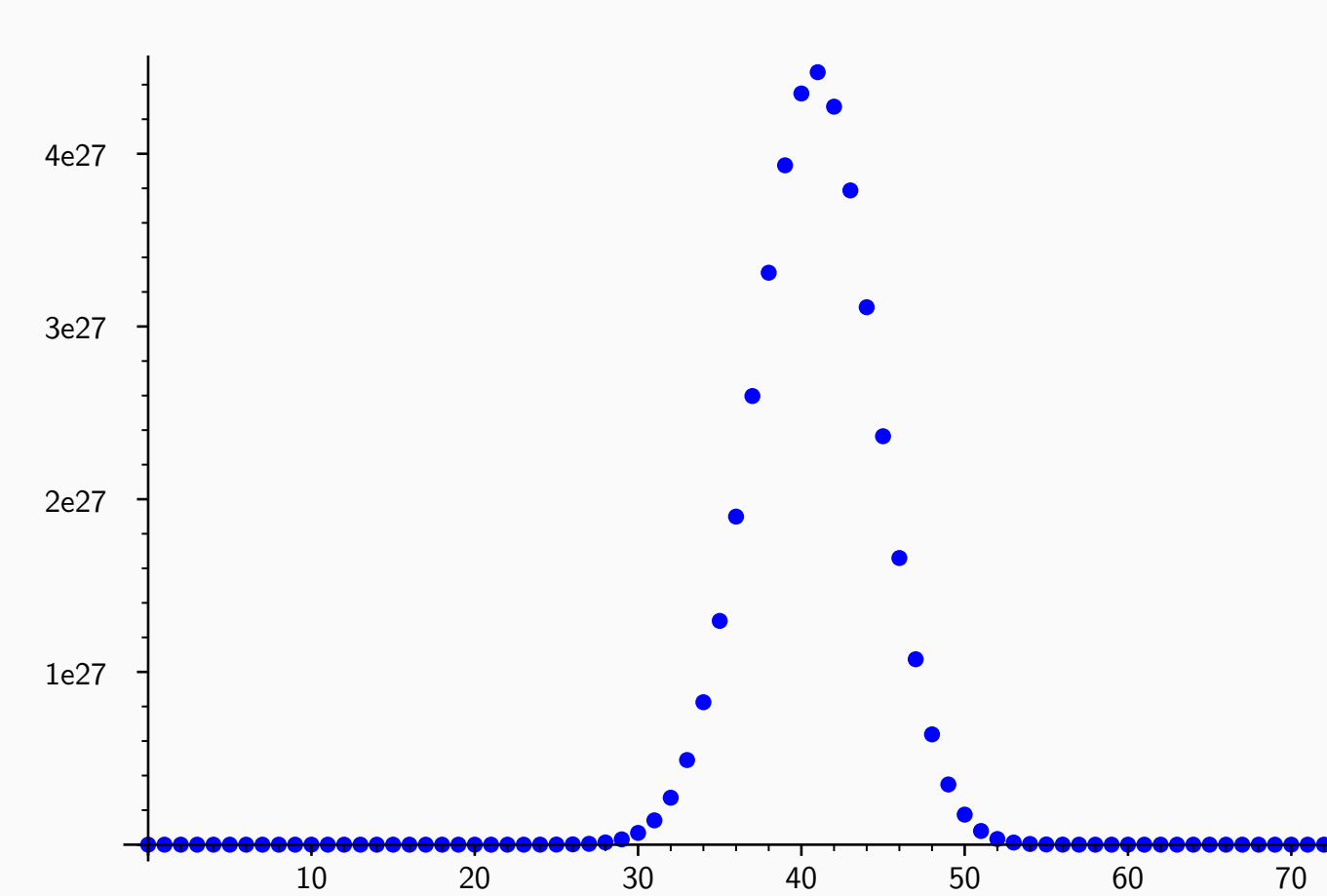
THM. Longest path: $\ell_{\max} = \lfloor \frac{3(n-1)}{2} \rfloor$, with $v_{n, \ell_{\max}} = \begin{cases} 1 & \text{if } n \text{ odd} \\ \lfloor \frac{3(n-1)}{2} \rfloor & \text{if } n \text{ even} \end{cases}$

Monotone paths on $\Delta(n, 2)$ - coherence

THM. Coherent iff when $i \xrightarrow{a} j$ precede $x \xrightarrow{z} y$ with $x < j$ then $j = z$ or $x = a$



Conjecture on log-concavity



CONJ. [De Loera] number of coherent monotone paths by length is log-concave for all polytopes
 i.e. for $\Delta(n, 2)$: $v_{n, \ell}$ log-concave

Here left: $v_{n, \ell}$ for $n = 50$
 With M , conjecture checked numerically up to $n = 150$