# A dichotomy for finitely generated subgroups of word hyperbolic groups

Goulnara N. Arzhantseva

ABSTRACT. Given L > 0 elements in a word hyperbolic group G, there exists a number M = M(G, L) > 0 such that at least one of the assertions is true: (i) these elements generate a free and quasiconvex subgroup of G; (ii) they are Nielsen equivalent to a system of L elements containing an element of length at most M up to conjugation in G. The constant M is given explicitly. The result is generalized to groups acting by isometries on Gromov hyperbolic spaces. For proof we use a graph method to represent finitely generated subgroups of a group.

### 1. Introduction

Let H be a subgroup of a word hyperbolic group G. It is known that either H is elementary (that is, it contains a cyclic subgroup of finite index) or H contains a non-abelian free subgroup of rank two. In the case G is torsion-free, there are, up to conjugacy, finitely many Nielsen equivalence classes of non-free subgroups of G generated by two elements [3].

Our main result gives a sufficient condition for H to be free and quasiconvex in G. It is an improvement of a result due to Gromov [4, 5.3.A].

THEOREM 1. For any  $\delta \geq 0$  and an integer L > 0 there exists a number  $M = M(\delta, L) > 0$  with the following property.

Let G be a  $\delta$ -hyperbolic group with respect to a finite generating set  $\mathcal{X}$  and H be a subgroup of G generated by  $h_1, \dots, h_L$ . Then at least one of the following assertions is true.

- (i) H is free on  $h_1, \dots, h_L$  and quasiconvex in G;
- (ii) The tuple (h<sub>1</sub>,...,h<sub>L</sub>) is Nielsen equivalent to an L-tuple (h'<sub>1</sub>,...,h'<sub>L</sub>) with h'<sub>1</sub> conjugate to an element in G of word length at most M with respect to X.

 $<sup>2000\</sup> Mathematics\ Subject\ Classification.\ 20F67; 20E07.$ 

Key words and phrases. Word hyperbolic groups, quasigeodesics, Nielsen equivalence.

The work has been supported in part by the Swiss National Science Foundation, No. PP002-68627.

The constant  $M = M(\delta, L)$  can be calculated explicitly. Then, as an immediate consequence of Theorem 1 we obtain Gromov's result. Let

$$\mathsf{Rad}\ H = 1/2 \inf_{h \in H \setminus \{\mathsf{id}\}} |[h]|,$$

where  $|[h]| = \inf_{h'} |h'|$  for all  $h' \in G$  conjugate to h and |h| is the word length of h. This is called the injectivity radius of H.

THEOREM 2. (cf. [4, 5.3.A]) Let G be a  $\delta$ -hyperbolic group and H be a subgroup of G generated by L elements. If

Rad 
$$H \ge \frac{1}{2} \left( \frac{14400(\delta+1)}{\ln 2} L^2 \right)^{1+1/2}$$

then H is free and quasiconvex<sup>1</sup> in G.

The following result is a natural generalization of Theorem 1 to groups acting on hyperbolic spaces.

THEOREM 3. For any  $\delta \geq 0$  and an integer L > 0 there exists a number  $M = M(\delta, L) > 0$  with the following property.

Let G be a group on L generators  $g_1, \dots, g_L$  acting on a  $\delta$ -hyperbolic space (X, d) by isometries. Then at least one of the following assertions is true.

- (i) G is free on  $g_1, \dots, g_L$  and for every  $x \in X$  the map  $G \to X$  which assigns to each  $g \in G$  the element  $gx \in X$  is a quasi-isometric embedding;
- (ii) The tuple  $(g_1, \ldots, g_L)$  is Nielsen equivalent to an L-tuple  $(g'_1, \ldots, g'_L)$  with  $d(g'_1y, y) < M$  for some  $y \in X$ .

The constant  $M = M(\delta, L)$  can be calculated explicitly.

For proof of Theorem 1 a representation of finitely generated subgroups of a group by labelled graphs is used [1]. The technique is of independent interest. In particular, transformations of a labelled graph defined below can be viewed as a generalization of free reductions and Nielsen reductions of tuples of group elements.

Theorem 3 is obtained with essentially the same methods, except that, in our arguments, instead of the word length metric, we refer to the metric on G induced from the group action.

I was informed that Kapovich and Weidmann showed both results (without an explicit estimate on the constant) independently of my work and by different methods [6].

Acknowledgments This work has been done during my visit to University Luis Pasteur, Strasbourg, May 2000. I thank Thomas Delzant for suggesting the problem and hospitality. I also thank the referee for useful comments.

### 2. Auxiliary information

Let G be a group and  $\mathcal{X}$  be a finite set of generators for G. We fix both for the rest of the paper. All words are assumed to be in the alphabet  $\mathcal{X}^{\pm 1}$ . We shall make no essential distinction between words and elements of G. If w and v are words then the notation  $w =_G v$  means that they represent the same group element.

<sup>1</sup>This lower bound is not optimal. For example, one can take  $\frac{1}{2} \left( \frac{14400(\delta+1)}{\ln 2} L^2 \right)^{1+\varepsilon}$  with  $\varepsilon > 0$  or a better constant satisfying condition (4) below.

**2.1. Graphs representing subgroups.** Let  $\Gamma$  be a graph. By an edge of  $\Gamma$  we mean a directed edge, i.e., an edge of  $\Gamma$  in the usual sense with any of its two possible directions. If e is an edge of  $\Gamma$  then  $e^{-1}$  denotes the edge with the opposite direction. A map  $\psi$  from the edges of  $\Gamma$  to  $\mathcal{X}^{\pm 1}$  is called a *labelling function* on  $\Gamma$  if it satisfies  $\psi(e^{-1}) = (\psi(e))^{-1}$  for any edge e. By the label  $\psi(p)$  of a path  $p = e_1 e_2 \dots e_k$  of length k in  $\Gamma$  we mean the word  $\psi(e_1)\psi(e_2)\dots\psi(e_k)$ . The label of a path of length 0 (which by definition is identified with a vertex of  $\Gamma$ ) is the empty word.

A labelled graph is a finite connected graph  $\Gamma$  with a labelling function  $\psi$  and a distinguished vertex O. Any labelled graph  $\Gamma$  represents a subgroup  $H(\Gamma)$  of a free group  $F = F(\mathcal{X})$ , which is the image of the fundamental group  $\pi_1(\Gamma, O)$  under the homomorphism induced by  $\psi$ . In other words,  $x \in H(\Gamma)$  if and only if x may be represented by a word which can be read on a circuit at O.

It is easy to see that any finitely generated subgroup  $H \leq F$  may be represented by a labelled graph. To do this, we first take words  $h_1, h_2, \dots, h_k$  in the alphabet  $\mathcal{X}^{\pm 1}$  that represent generators of H. Next we take a rose of k circles attached to a point O and make each of the circles a circuit labelled  $h_i$ ,  $1 \leq i \leq k$ . For the resulting labelled graph  $\Gamma$ , we obviously have  $H(\Gamma) = H$ .

We define two types of transformations of a labelled graph  $\Gamma$ , which preserve the subgroup  $H(\Gamma)$  and which we call *reductions*. A transformation of the first type is identification of two edges with the same label and the same initial vertex. A transformation of the second type is removal of a vertex of degree 1 other than O, together with the incident edge.

A labelled graph  $\Gamma$  is said to be *reduced* if it admits no reductions, that is, it has no pair of edges with the same label and initial vertex and no vertices of degree 1 with the possible exception of the distinguished vertex O.

Starting from a labelled graph  $\Gamma$  with  $H(\Gamma) = H$  and performing all possible reductions, we reach a reduced labelled graph which represents the subgroup H. It is known [10, 8] that a reduced labelled graph representing a subgroup  $H \leq F$  is unique up to graph isomorphism (that is, it does not depend on the order of reductions, the choice of the initial graph  $\Gamma$ , and the choice of generators for H).

If  $\Gamma$  is a reduced labelled graph then it is easy to see that a reduced word w represents an element of  $H(\Gamma)$  if and only if w is the label of a reduced circuit at O in  $\Gamma$ . It follows in particular that the label of a path p in  $\Gamma$  starting at O represents an element of  $H(\Gamma)$  only if O is also the terminal vertex of p.

A finitely generated subgroup H of G can also be presented by a labelled graph  $\Gamma$ . It suffices to consider the graph obtained from a lift of the subgroup generators under the natural homomorphism  $F \to G$ . However the reduced form of  $\Gamma$  is not unique in this case. For  $\Gamma$  representing  $H \leq G$  we introduce a transformation of the third type as well, see [1].

Denote by  $p_{-}(p_{+})$  the initial (terminal) vertex of a path p. By an *arc* we mean a path p all of whose vertices except  $p_{-}$  and  $p_{+}$  have degree 2 and are distinct from the distinguished vertex O.

- (arc reduction) Let vertices  $O_1$  and  $O_2$  in  $\Gamma$  be joined by a path p so that  $\psi(p) \equiv w$  and the word w is equal to some word v in G. Let q' be an arc in  $\Gamma$  which is a subpath of p. First let us add to  $\Gamma$  a new graph formed by a single arc q with label  $\psi(q) \equiv v$  such that  $O_1 = q_-$  and  $O_2 = q_+$ 

are the only common points of the new graph and  $\Gamma$ . Then let us remove from  $\Gamma$  all edges and vertices of q' except  $q'_{-}$  and  $q'_{+}$ .

REMARK. We remove an arc whenever we add another one. Hence the transformations preserve the Euler characteristic of  $\Gamma$  which is the number of vertices minus the number of edges. Thus if the fundamental group  $\pi_1(\Gamma)$  of  $\Gamma$  is generated by L loops, then for any graph  $\Gamma'$  obtained from  $\Gamma$  by the transformations,  $\pi_1(\Gamma')$ is L-generated as well.

LEMMA 4. [1, Lemma 1] If a labelled graph  $\Gamma'$  is carried into a labelled graph  $\Gamma$  by transformations of types 1–3 or their inverses, then it represents the same subgroup of G as  $\Gamma$  does.

## 2.2. Nielsen equivalence.

DEFINITION 5. Let  $\mathcal{U} = (u_1, \ldots, u_k)$  and  $\tilde{\mathcal{U}} = (\tilde{u}_1, \ldots, \tilde{u}_k)$  be k-tuples of elements of G. They are Nielsen equivalent if  $\mathcal{U}$  can be carried into  $\tilde{\mathcal{U}}$  by a finitely many regular elementary Nielsen transformations defined as follows:

(t1) replace  $u_i$  by  $u_i^{-1}$  for some i;

(t2) replace  $u_i$  by  $u_i u_j$  for some  $i \neq j$ ;

In both cases  $u_t$  remains unchanged if  $t \neq i$ .

REMARK. If  $w_1, \dots, w_k$  are words we shall consider  $(w_1, \dots, w_k)$  as the k-tuple of group elements represented by these words.

Note that the regular elementary Nielsen transformations generate a group containing every permutation of the  $u_i$ . It is obvious that if a k-tuple  $\mathcal{U}$  is carried by a regular elementary Nielsen transformation into a k-tuple  $\tilde{\mathcal{U}}$ , then they generate the same subgroup of G. Proceeding by induction we see that Nielsen equivalent tuples of group elements generate the same subgroup of G. Moreover, in a free group any two bases of a finitely generated subgroup are Nielsen equivalent. Here by a *basis* of a subgroup we mean a set of group elements freely generating the subgroup. In general, this is not true for a non-free group G. However we have the following

LEMMA 6. Let H be a subgroup of G generated by elements represented by words  $h_1, \dots, h_k$ . Let  $\Gamma$  be a graph obtained from a rose  $\Omega$  of k circuits at a vertex O labelled by  $h_1, \dots, h_k$  by transformations of 1-3 types or their inverses. Then for any basis  $l_1, \dots, l_k$  of  $\pi_1(\Gamma, O)$  the k-tuple  $(\psi(l_1), \dots, \psi(l_k))$  is Nielsen equivalent to  $(h_1, \dots, h_k)$ .

PROOF. Let  $\Gamma$  be obtained from  $\Omega$  by transformations of types 1–3 or their inverses. Hence there is a sequence of graphs  $\Gamma_0 = \Omega, \ldots, \Gamma_j, \ldots, \Gamma_f = \Gamma$ , so that  $\Gamma_i$  is carried into  $\Gamma_{i+1}$  by only one transformation. If f = 0, i.e.  $\Gamma$  coincides with  $\Omega$ , the lemma obviously holds. Suppose that f > 0. We claim that for each basis  $s_1, \cdots, s_k$  of  $\pi_1(\Gamma_{j+1}, O)$  there exists a basis  $t_1, \cdots, t_k$  of  $\pi_1(\Gamma_j, O)$  such that  $(\psi(t_1), \ldots, \psi(t_k))$  is Nielsen equivalent to  $(\psi(s_1), \ldots, \psi(s_k))$ . This is obvious if the performed transformation is a reduction. For  $t_i$  we take a loop in  $\Gamma_j$  which were carried by the reduction into  $s_i$ ,  $1 \leq i \leq k$  (possibly  $t_i$  and  $s_i$  coincide). Let's now consider an arc reduction. Let q be the added arc and q' the removed arc. By definition, there is a path  $p \in \Gamma_j$  with the same endpoints as  $q \in \Gamma_{j+1}$  such that p contains q' as a subpath and  $\psi(p) =_G \psi(q)$  in G. Let's define  $t_i \in \Gamma_j$ . We start with  $s_i \in \Gamma_{j+1}$ . We replace each entry of q in  $s_i$  by p. In such a way we obtain loops  $t_i$  at O in  $\Gamma_j$ ,  $1 \le i \le k$ . Clearly, it will be a basis of  $\pi_1(\Gamma_j, O)$ . The k-tuple of elements represented by labels of these loops is Nielsen equivalent to  $(\psi(s_1), \ldots, \psi(s_k))$ . Indeed, we have  $\psi(s_i) =_G \psi(t_i)$  as  $\psi(p) =_G \psi(q)$  in G.

The claim is true for any  $1 \leq j \leq f$ . Thus, for any basis  $l_1, \dots, l_k$  of  $\pi_1(\Gamma, O)$ there exists a basis  $l'_1, \dots, l'_k$  of  $\pi_1(\Omega, O)$  such that  $(\psi(l_1), \dots, \psi(l_k))$  is Nielsen equivalent to  $(\psi(l'_1), \dots, \psi(l'_k))$ . As it was mentioned above the labels of any two bases of  $\pi_1(\Omega, O)$  represent Nielsen equivalent k-tuples. Hence  $(\psi(l'_1), \dots, \psi(l'_k))$  is Nielsen equivalent to  $(h_1, \dots, h_k)$ . This completes the proof.

**2.3. Word hyperbolic groups.** For a background on word hyperbolic groups we refer to [4, 5], and [2].

Let C(G) be the Cayley graph of G with respect to  $\mathcal{X}$ . It is a graph whose set of vertices is G and whose set of edges is  $G \times (\mathcal{X}^{\pm 1})$ . An edge (h, x) starts at the vertex  $h \in G$  and ends at the vertex hx. We consider an edge (h, x) of C(G) as labelled by the letter x. The label  $\varphi(\rho)$  of a path  $\rho = e_1 e_2 \dots e_n$  in C(G) is the word  $\varphi(e_1)\varphi(e_2)\dots\varphi(e_n)$  where  $\varphi(e_i)$  is the label of the edge  $e_i$ . We regard  $\varphi(\rho)$ as an element of G. We endow C(G) with a metric by assigning to each edge the metric of the unit segment [0, 1] and then defining the distance |g - h| to be the length of a shortest path between g and h. Thus C(G) becomes a geodesic metric space. Obviously, this metric on C(G) is invariant under the natural left action of G. For any  $g \in G$ , we write |g| for the length of a shortest path from the unit element to g. In particular,  $|g - h| = |g^{-1}h|$ .

DEFINITION 7. [4, 6.3.B] The *Gromov product* of points x and y of a metric space  $\mathcal{M}$  with respect to a point  $z \in \mathcal{M}$  is defined to be

$$(x|y)_z = \frac{1}{2}(|x-z| + |y-z| - |y-x|)$$

where |x - y| denotes the distance between x and y.

A geodesic metric space  $\mathcal{M}$  is called  $\delta$ -hyperbolic, for  $\delta \geq 0$ , if

$$(x|y)_w \ge \min\{(x|z)_w, (z|y)_w\} - \delta$$

for any  $x, y, z, w \in \mathcal{M}$ .

A group G is  $\delta$ -hyperbolic with respect to a finite generating set  $\mathcal{X}$  if the Cayley graph C(G) with respect to  $\mathcal{X}$  is a  $\delta$ -hyperbolic space. The group G is called *word* hyperbolic if it is  $\delta$ -hyperbolic for some  $\delta \geq 0$  and  $\mathcal{X}$ .

It turns out that the word hyperbolicity of a group is independent of the finite generating set chosen [4, 2.3.E].

LEMMA 8. [9, Lemma 21] Let  $c \ge 14\delta$  and  $c_1 > 12(c+\delta)$ , and suppose that a geodesic n-gon  $[x_1, \ldots, x_n]$  in a  $\delta$ -hyperbolic metric space satisfies the conditions  $|x_{i-1} - x_i| > c_1$  for  $i = 2, \ldots, n$  and  $(x_{i-2}|x_i)_{x_{i-1}} < c$  for  $i = 3, \ldots, n$ . Then the polygonal line  $\rho = [x_1, x_2] \cup [x_2, x_3] \cup \cdots \cup [x_{n-1}, x_n]$  is contained in the 2cneighbourhood of the side  $\tau = [x_n, x_1]$ , and the side  $\tau$  is contained in the 14 $\delta$ neighbourhood of  $\rho$ .

COROLLARY 9. Under the assumptions of the previous lemma there is a constant  $\lambda = \lambda(c, c_1) > 0$  such that

 $|\tau| \ge \lambda \|\rho\|,$ 

where  $\|\rho\|$  is the length of the path  $\rho$ .

PROOF. By the lemma, there exist points  $v_1, \ldots, v_n$  on  $[x_n, x_1]$  such that  $|x_i - v_i| \leq 2c$ ,  $1 \leq i \leq n$ . The hypothesis of the lemma imply easily that  $v_{i-1}$  is located between  $v_i$  and  $v_{i-2}$ . So,

$$|\tau| = \sum_{i=2}^{n} |v_{i-1} - v_i| \ge \sum_{i=2}^{n} (|x_{i-1} - x_i| - 4c) = \|\rho\| - (n-1)4c.$$

Taking  $\lambda = c/c_1$  and using  $\|\rho\| > (n-1)c_1$  we obtain  $|\tau|/\|\rho\| \ge \lambda$ .

The following lemma is obvious as any side of a geodesic triangle in a  $\delta$ -hyperbolic space belongs to the  $4\delta$ -neighbourhood of the union of the other two sides [5, Ch.2, Pr. 21].

LEMMA 10. (cp. [4, § 7]) Let  $\rho$  be a path in a  $\delta$ -hyperbolic space. Then for every point A on a geodesic segment with the same endpoints as  $\rho$ 

$$\inf_{B \in \rho} |A - B| \le 4\delta \log_2 \|\rho\| + 1.$$

For a word w, the length ||w|| is the length of a path in C(G) labelled by w and |w| is the length of a geodesic in C(G) between the same endpoints. If ||w|| = |w| the word w is called a *geodesic word*. The notation  $w \equiv xy$  means that w can be decomposed, as a word, into a product of two words which represent elements  $x, y \in G$ . The following fact is known (see the proof of  $(P1, \delta) \Rightarrow (P2, 4\delta)$  in [5, Ch.2, Pr. 21]). If G is  $\delta$ -hyperbolic then

(H) for any two geodesic words u and v, if  $u \equiv u_1 u_2$ ,  $v \equiv v_1 v_2$  and  $|u_1| = |v_1| \le \frac{1}{2}(|u| + |v| - |u^{-1}v|)$  then  $|u_1^{-1}v_1| \le 4\delta$ .

Notice that if G is  $\delta$ -hyperbolic then it is also  $\delta'$ -hyperbolic for any  $\delta' > \delta$ . So we can always assume  $\delta \geq 1$ .

## 3. Proof of Theorem 1

From now on, we assume G to be word hyperbolic and fix a number  $\delta \geq 1$  such that (H) holds.

The following two lemmas are the main technical tools for the proof of our theorems.

LEMMA 11. Let  $\delta \ge 1$ ,  $K > 16\delta + 1$ . Let x be a word with  $|x| < K \log_2 ||x||$ and  $||x|| \ge 2^4$ . Then there exists a subword y of x such that

$$\frac{1}{2}||x|| \le ||y|| < ||x||$$
 and  $|y| < K \log_2 ||y||$ .

PROOF. We denote by  $\eta$  a path starting at the unit vertex of the Cayley graph C(G) and labelled by x. Let  $\rho$  be a geodesic path in C(G) with the same endpoints as  $\eta$  and z be the label of  $\rho$ . Note that ||z|| = |x|. We take a middle point A on  $\rho$  so that  $z = z_1 z_2$ , where A is a terminal vertex of a subpath labelled by  $z_1$  and  $||z_1|| = \left\lceil \frac{||z||}{2} \right\rceil$ .

Suppose that  $||z|| < 9\delta \log_2 ||x||$ . Then, for a desired subword y we take x without its terminal letter, i.e. ||y|| = ||x|| - 1. The assumption on ||x||, |x|, and K implies easily the needed inequalities on ||y|| and |y|.

The remaining case is  $||z|| \ge 9\delta \log_2 ||x||$ . By Lemma 10, there is a point B on  $\eta$ such that  $|A - B| \le 4\delta \log_2 ||\eta|| + 1 = 4\delta \log_2 ||x|| + 1$ . Then B gives a decomposition of  $\eta$  into two subpaths labelled by words  $x_1$  and  $x_2$  with  $x = x_1 x_2$  and

$$|x_i| \le \frac{K}{2} \log_2 ||x|| + 4\delta \log_2 ||x|| + 1 = \left(\frac{K}{2} + 4\delta\right) \log_2 ||x|| + 1.$$

The words  $x_1$  and  $x_2$  are nontrivial which easily follows from the assumption on ||z|| and the bound on |A-B|. Hence  $\frac{||x||}{2} \le ||x_i|| < ||x||$  for i = 1 or i = 2. Without loss of generality, we assume that  $||x_1|| \ge ||x_2||$ . Since  $||x|| \ge 2^4$  and  $K > 16\delta + 1$ we have

$$|x_i| \le \left(\frac{K}{2} + 4\delta\right) \log_2 ||x|| + 1 < K \log_2 \frac{||x||}{2} \le K \log_2 ||x_1||.$$

Thus, we can take  $x_1$  for a desired subword y of x.

COROLLARY 12. Let  $\delta \geq 1$ ,  $K > 16\delta + 1$ . Let  $D \geq 2^4$  and x be a word with  $|x| < K \log_2 ||x||$  and  $||x|| \ge D$ . Then there is a subword y of x such that

 $\frac{D}{2} \le \|y\| < D \quad and \quad |y| < K \log_2 D.$ 

PROOF. By Lemma 11, there is a subword y of x with

(1) 
$$||y|| \ge \frac{D}{2}$$
 and  $|y| < K \log_2 ||y||$ .

We take such a y of the minimal possible length. We have ||y|| < D for otherwise using the previous lemma we could find a subword y' of y satisfying (1) with ||y'|| < ||y'||| < ||y'|| < ||y||||y||. Hence  $|y| < K \log_2 ||y|| < K \log_2 D$ . 

LEMMA 13. Let  $\delta \geq 1$ . For any  $T > 2^{\delta+7}$  there are numbers  $\lambda = \lambda(T, \delta) > 0$ 

and  $D_1 = D_1(T, \delta) > 0$  with the following property: Let x be a word with  $||x|| \ge D_1$ . If  $|y| \ge 20 \left(\frac{\delta T}{\log_2 T}\right)^{1/2} \log_2 ||y||$  for any subword y of x with  $||y|| \ge T$  then  $|x| \ge \lambda ||x||$ .

PROOF. Set c = T,  $c_1 = 12(c+\delta)+2$ , and  $D_1 = 2^{\frac{c_1}{K}}$ , where  $K = 20 \left(\frac{\delta T}{\log_2 T}\right)^{1/2}$ . Take  $\lambda = \lambda(c, c_1)$  by Corollary 9.

Suppose the lemma does not hold. Then there is a word x with  $||x|| \geq D_1$ and  $|x| < \lambda \|x\|$  such that for any subword y of x with  $\|y\| \ge T$  the inequality  $|y| \ge K \log_2 \|y\| \text{ holds for } K = 20 \left(\frac{\delta T}{\log_2 T}\right)^{1/2}.$ 

We take any decomposition  $x \equiv x_1 x_2 \dots x_s$  where  $D_1 \leq ||x_i|| \leq 2D_1$  for  $1 \leq ||x_i|| \leq 2D_1$  $i \leq s$ . For each  $x_i$ , we choose a shortest word  $z_i$  representing the same element of G. It is easy to see that  $D_1 > T$ . Then, by our assumption we have  $||z_i|| = |x_i| \ge |z_i|$  $K\log_2 \|x_i\| \ge K\log_2 D_1 \ge c_1.$ 

Let  $\rho$  be a path in C(G) labelled with  $x_1 x_2 \dots x_s$ . Each  $z_i$  labels a geodesic path with the same endpoints as the subpath of  $\rho$  labelled with  $x_i$ . By Lemma 8 applied to the (s+1)-gon in C(G) formed by the endpoints of the paths labelled with  $x_i$ , for some *i* we have

(2) 
$$|z_i z_{i+1}| < ||z_i|| + ||z_{i+1}|| - 2c.$$

Let us decompose  $z_i$  and  $z_{i+1}$  so that  $z_i \equiv y_i z'_i$  and  $z_{i+1} \equiv z'_{i+1} y_{i+1}$  with  $||z'_i|| = ||z'_{i+1}|| = c$  for some words  $y_i, y_{i+1}$ . By (2) and (H),  $|z'_i z'_{i+1}| \leq 4\delta$ . By Lemma 10 we find a terminal subword  $x'_i$  of  $x_i$  and an initial subword  $x'_{i+1}$  of  $x_{i+1}$  such that  $|x'_i(z'_i)^{-1}| \leq 4\delta \log_2 ||x_i|| + 1$  and  $|(z'_{i+1})^{-1} x'_{i+1}| \leq 4\delta \log_2 ||x_{i+1}|| + 1$ . Thus,

$$|x_i'x_{i+1}'| \le 8\delta \log_2 2D_1 + 2 + 4\delta$$

Since  $||x_i'||, ||x_{i+1}'|| \ge c - 4\delta \log_2 2D_1 - 1$  we have

$$\|x_i'x_{i+1}'\| \ge 2c - 8\delta \log_2 2D_1 - 2.$$

It is easy to check that for  $T > 2^{\delta+7}$  we have  $20 \left(\frac{\delta T}{\log_2 T}\right)^{1/2} \ge 104\delta$  which implies  $c \ge 8\delta \left(1 + \frac{c_1}{K}\right) + 2$ . From the last inequality we deduce that  $2c - 8\delta \log_2 2D_1 - 2 \ge T$ .

Now we prove that  $8\delta \log_2 2D_1 + 2 + 4\delta < K \log_2(2c - 8\delta \log_2 2D_1 - 2)$ . Indeed, since for  $K \ge 104\delta$  we have  $2c - 8\delta \left(1 + \frac{c_1}{K}\right) - 2 \ge c$ , it suffices to verify that

$$8\delta\left(1 + \frac{c_1}{K}\right) + 2 + 4\delta < K \log_2 c.$$

Using  $K \ge 104\delta$  and  $c_1 = 12(c+\delta)+2$  we obtain  $8\delta\left(1+\frac{c_1}{K}\right)+2+4\delta \le 96\delta\frac{c}{K}+13\delta+3$ . The latter is less or equal to  $112\delta\frac{c}{K}$  as c > K and  $\delta \ge 1$ . Now by the choice of c and K,  $112\delta\frac{c}{K} < K\log_2 c$ .

Thus we have found a subword  $y \equiv x'_i x'_{i+1}$  of x such that  $||y|| \ge T$  and  $|y| < K \log_2 ||y||$ . This contradicts the assumption.

PROOF OF THEOREM 1. Given  $\delta \ge 1$ , L > 0, and any  $K > 16\delta + 1$ , we choose  $D = D(\delta, L)$  by the inequality

(3) 
$$\frac{D}{\log_2 D} > 6KL$$

Let G be a  $\delta$ -hyperbolic group, H be a subgroup of G generated by L elements represented by words  $h_1, \dots, h_L$ . Let  $\Omega$  be a rose of L circuits at a vertex O labelled by  $h_1, \dots, h_L$ . Let  $\Gamma = \Gamma(H)$  be a graph representing H which is obtained from  $\Omega$  by transformations of types 1–3 and has the minimal possible number of edges. In particular,  $\Gamma$  is reduced and  $\pi_1(\Gamma)$  is L-generated. By Lemma 6, for any basis of  $\pi_1(\Gamma, O)$ , its image in G under the labelling function is Nielsen equivalent to the tuple  $(h_1, \dots, h_L)$ .

Suppose that H is not free on generators represented by  $h_1, \dots, h_L$ . Then there is a closed reduced path p in  $\Gamma$  starting at O labelled by a nontrivial word xrepresenting the identity element of G, i.e.  $x =_G 1$ . We take such a p of minimal length. There are two cases.

First suppose ||x|| < D. Then p contains a simple circuit  $\nu$  of length < D as a subpath (possibly,  $\nu = p$ ). Let  $\mu$  be any reduced path starting at O and ending at a vertex v on  $\nu$ . Then the label of  $\mu\nu\mu^{-1}$  represents an element  $h'_1 \in H$ . Obviously,  $h'_1$  is conjugate to an element of length less then D that is represented by the label of  $\nu$ . Moreover,  $h'_1$  can be included in a system of generators of H. Indeed, suppose that  $\nu = \nu_1 e \nu_2$ , where  $\nu_1$  starts at v and e is an edge of  $\nu$ . Then the tripod rooted at v consisting of tree branches  $\mu$ ,  $\nu_1$ , and  $\nu_2$ , can be included in a maximal tree spanning  $\Gamma$ . Hence  $h'_1$  is the label of one of L generators of  $\pi_1(\Gamma, O)$  given by this maximal tree. Thus, by Lemma 6,  $(h_1, \ldots, h_L)$  is Nielsen equivalent to an L-tuple  $(h'_1, \ldots, h'_L)$  and the conclusion (ii) of the theorem holds. The remaining case is  $||x|| \ge D$ . Since  $x =_G 1$ , we have  $0 = |x| < K \log_2 ||x||$ . By Corollary 12, there is a subword y of x with  $\frac{D}{2} \le ||y|| < D$  and  $|y| < K \log_2 D$ . We may assume that y labels a simple path  $\gamma$ . Otherwise,  $\gamma$  contains a simple circuit  $\nu$  of length < D as a subpath and we proceed as above. The number of arcs in  $\Gamma$  is less then 3L. Since  $||y|| \ge \frac{D}{2}$ , there is a subword u of y of length at least  $\frac{D}{6L}$  which labels an arc. Using a transformation of  $\Gamma$  of the third type, we remove this arc and add a new arc of length |y| with the same endpoints as  $\gamma$ . We label this arc by a shortest word representing the same group element as y. By (3) and the choice of y, the length of the new arc is less then  $\frac{D}{6L}$ . So, the number of edges in the obtained graph is less than one in  $\Gamma$ . This contradicts the choice of  $\Gamma$ .

Suppose that H is free on  $h_1, \dots, h_L$  but not quasiconvex. We are going to find a constant  $T = T(\delta, L)$  such that  $(h_1, \dots, h_L)$  is Nielsen equivalent to an L-tuple containing an element conjugate to an element of length at most T. Take any  $T > 2^{\delta+7}$  satisfying

(4) 
$$\frac{T}{\log_2 T} > 6KL_t$$

where  $K = 20 \left(\frac{\delta T}{\log_2 T}\right)^{1/2}$ . Choose  $\lambda = \lambda(T, \delta) > 0$  and  $D_1 = D_1(T, \delta) > 0$  by Lemma 13.

Since *H* is supposed to be free but non-quasiconvex there are a reduced circuit at *O* in  $\Gamma$  labelled by a word *z* and a subword *x* of *z* such that either *x* represents the identity element of *G* or  $||x|| \ge D_1$  and  $|x| < \lambda ||x||$ .

We repeat the proof as above in the case when  $x =_G 1$  slightly modifying the subcase  $||x|| < D_1$ . Namely, if x labels a simple path then we identify the endpoints of this path removing an arc of this path. Since  $x =_G 1$  we obtain a new labelled graph representing H with less number of edges than one in  $\Gamma$ . This is a contradiction. If the path labelled by x is not simple it contains a circuit of length less than  $D_1$ .

In the second case, by Lemma 13, there is a subword y of x such that  $||y|| \ge T$ and  $|y| < K \log_2 ||y||$ . Using Corollary 12 we can assume that  $T/2 \le ||y|| < T$ . Then, arguing as above we conclude that either there is a circuit of length less than T in  $\Gamma$  or we can use a transformation of the third type reducing the number of edges in  $\Gamma$ .

We take  $M = \max\{D, T\}$  finishing the proof.

Thus,  $M = \max\{D, T\}$  with constants defined by (3) and (4). It is now a routine to check that Theorem 2 is a straightforward consequence of Theorem 1.

Theorem 3 can be shown by mimic arguments above where the word length is replaced by the length function induced by the group action. Namely, if  $x_0 \in X$  is arbitrary, then the length of an element  $g \in G$  is defined by  $\ell(g) := d(gx_0, x_0)$  and the left-invariant distance between group elements g and h is given by  $\ell(g^{-1}h)$ .

### References

- G.N.Arzhantseva and A.Yu.Ol'shanskii, The class of groups all of whose subgroups with lesser number of generators are free is generic, Mat. Zametki, 59(4)(1996), 489-496 (in Russian), English translation in Math. Notes, 59(4)(1996), 350-355.
- [2] M. Coornaert, T. Delzant, and A. Papadopoulos, Géométrie et théorie des groupes. Les groupes hyperboliques de Gromov, Lecture Notes in Mathematics, 1441, Springer-Verlag, 1990.

#### GOULNARA N. ARZHANTSEVA

- [3] T. Delzant, Sous-groupes deux générateurs des groupes hyperboliques, Group theory from a geometrical viewpoint (Trieste, 1990), 177–189, World Sci. Publishing, River Edge, NJ, 1991.
- M. Gromov, Hyperbolic groups, Essays in group theory, Math. Sci. Res. Inst. Publ., 8, Springer, New York-Berlin, 1987, 75–263.
- [5] E. Ghys and P. de la Harpe (editors), Sur les groupes hyperboliques d'après Mikhael Gromov, Progress in Mathematics, 83, Birkhauser, 1990.
- [6] I. Kapovich and R. Weidmann, Nielsen methods and groups acting on hyperbolic spaces, Geom. Dedicata 98 (2003), 95-121.
- [7] R. Lyndon and P. Schupp, Combinatorial group theory, Springer-Verlag, 1977.
- [8] S. W. Margolis and J. C. Meakin, Free inverse monoids and graph immersions, Int. J. Algebra and Comput., 3(1) (1993), 79–99.
- [9] A. Yu. Ol'shanskii, Periodic quotients of hyperbolic groups, Mat. Zbornik, 182(4)(1991), 543-567 (in Russian), English translation in Math. USSR Sbornik 72(2) (1992).
- [10] J. Stallings, Topology of finite graphs, Inv.Math., **71** (1983), 552–565.

Université de Genève, Section de mathématiques, 2-4, rue du Lièvre, CH-1211 Genève24

 $E\text{-}mail\ address:$  Goulnara.Arjantseva@math.unige.ch