

Distal and non-distal ordered abelian groups

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Indiscernible sequences

Convention: \mathbb{M} is a first order structure, possibly highly saturated.

Definition

A sequence $(a_i)_{i \in I}$ from \mathbb{M}^n is **A -indiscernible** if for all $i_1 < \dots < i_m$ and $j_1 < \dots < j_m$ from I we have $a_{i_1} \cdots a_{i_m} \equiv_A a_{j_1} \cdots a_{j_m}$

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Example

- A strictly increasing sequence (q_n) in $(\mathbb{Q}, <)$
- A sequence (a_n) of algebraically independent numbers in $(\mathbb{C}; 0, 1, +, -, \cdot)$, for example, $(\exp(\sqrt{p_n}))$, where $p_n = n$ th prime.

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The non-independence property (NIP)

Definition (for monster model \mathbb{M})

\mathbb{M} is NIP iff for every formula $\varphi(x, y)$, every indiscernible sequence $(a_i)_{i \in I}$ from $\mathbb{M}^{|x|}$ and every $b \in \mathbb{M}^{|y|}$, there is $\epsilon \in \{0, 1\}$ such that eventually $\models \varphi(a_i, b)^\epsilon$.

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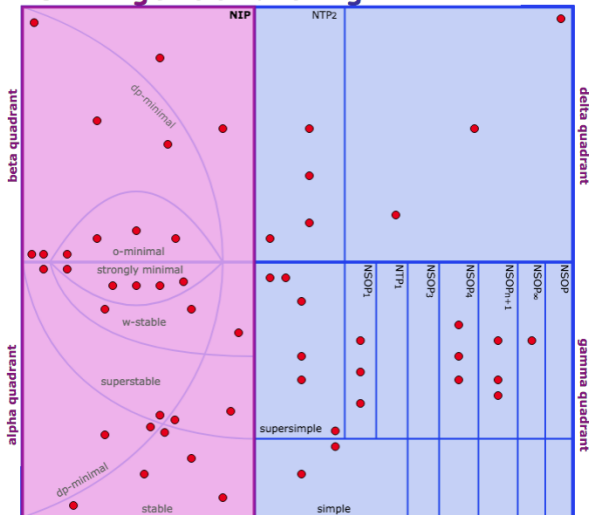
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More examples of NIP theories

forking and dividing



Questions? Suggestions? Corrections? email mgconant@nd.edu

[References](#) [Update Log](#)

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Map of the Universe

Nice Properties of Theories

ω -stable	superstable	stable (NOP)	
strongly minimal	o-minimal	dp-minimal	
NIP	supersimple	simple (NTP)	
NSOP ₁	NTP ₁	NTP ₂	NSOP
NSOP ₃	NSOP ₄	NSOP _{n+1}	NSOP _{∞}

Click a property above to highlight region and display details. Or click the map for specific region information.

Reset

NIP (dependent)

Examples

- $(\mathbb{Q}^n, <, \dots, <_n)$
- $(\mathbb{T}, +, \cdot, 0, 1, \theta, \leq)$

Contains:

- dp-minimal
- o-minimal
- strongly minimal
- stable
- superstable
- ω -stable

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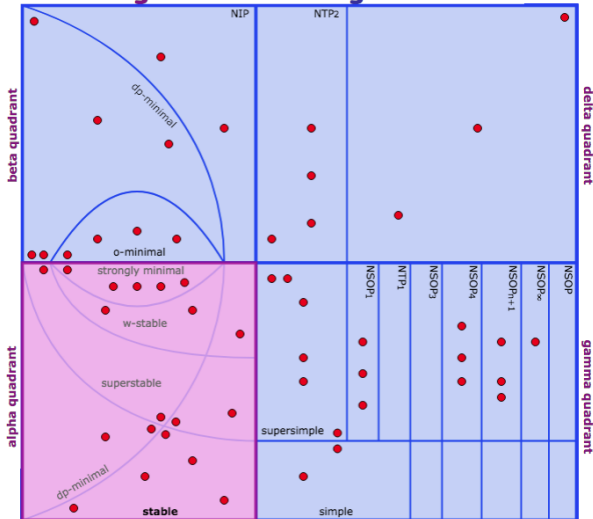
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Examples of stable structures

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stable (NOP)

Examples

- infinitely refining equivalence relations
- a strictly stable superflat graph
- infinitely cross-cutting equivalence relations
- DCF_p
- free group on $n > 1$ generators
- SCF_n^p
- $(\mathbb{Z}^n, +, 0)$

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- $G = \bigoplus_{q \in \mathbb{Q}} \mathbb{Z}_{(2)} \epsilon_q$ with the lexicographic order.
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Theorem (Cluckers, Halupczok)

In the theory of ordered abelian groups, each L_{qe} -formula $\psi(\bar{x}, \bar{\eta})$, where \bar{x} are home sort variables, and $\bar{\eta}$ are auxiliary sort variables, is equivalent to an L_{qe} -formula $\phi(\bar{x}, \bar{\eta})$ in family union form:

$$\phi(\bar{x}, \bar{\eta}) = \bigvee_{i=1}^k \exists \bar{\theta} (\xi_i(\bar{\eta}, \bar{\theta}) \wedge \psi_i(\bar{x}, \bar{\theta})),$$

where $\bar{\theta}$ are auxiliary sort variables, the formulas $\xi_i(\bar{\eta}, \bar{\theta})$ live purely in the auxiliary sorts, each $\psi_i(\bar{x}, \bar{\theta})$ is a conjunction of literals (i.e., of atoms and negated atoms), and for any ordered abelian group G and any $\bar{\beta}$ in the auxiliary sort of G corresponding to $\bar{\eta}$, the $L_{qe}(G)$ -formulas $\{\xi_i(\bar{\beta}, \bar{\alpha}) \wedge \psi_i(\bar{x}, \bar{\alpha}) : 1 \leq i \leq k, \bar{\alpha} \in \text{auxiliary sorts of } G\}$ are pairwise inconsistent.

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Suppose G is an ordered abelian group such that S_p is finite for all primes p . Then the following are equivalent:

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