

# Global Optimization in the COCONUT project



# Outline of Algorithm API Design Inference Engines Examples

Hermann Schichl,
Arnold Neumaier, Eric Monfroy,
and the COCONUT partners



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Project funded by the Future and Emerging Technologies arm of the IST Programme FET-Open scheme



### The COCONUT project



- European Union research and development project
- Partners from six European universities:
   Nantes, Lausanne, Vienna
   Louvain-la-Neuve, Coimbra, Darmstadt
   and an industrial partner:
   ILOG
- Aimed at the integration of the existing approaches to continuous global optimization and constraint satisfaction
- December 2000 November 2003







### **Basic Modular Setup**



**Strategy Eng.** 

Main part of A lgorithm, makes decisions

### Inference Eng.

P roblem structure analysis, local optimization, constraint satisfaction, interval analysis, Linear relaxation, convex opt., bisection,...

### **Reports**

P roduces files (A M P L), human readable output, checkpointing

### Management

P roblem management, R esources, P arallelization



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### Modular API design



- The API is designed to make the development of the various module types independent of each other and independent of the internal model representation.
- A collection of C++ classes.
- Uses FILIB++ and MTL.
- Supports dynamic linking.







### Modular API design (cont.d)



- All inference engines are subclasses of one class, so they have the same basic structure.
- Communication with the strategy engine by a database-like communication.
- The API implementation (w/o inference engines) consists of 44000 lines of C++ code and a few perl scripts, organized into 128 files, occupying 1.3 MB of disk space.

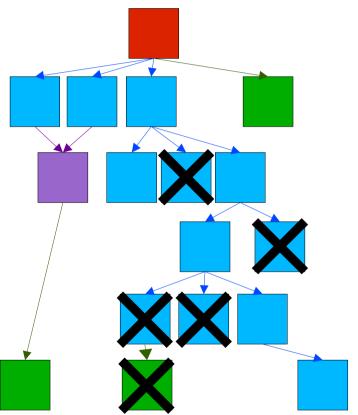






### Search graph





- Starts at the original model
- Contains relaxations
- and splits.
- It is not a tree since it might also contain glueings.
- Some of the nodes will be terminal, since they can be solved completely.









### **Model Relations**



min 
$$c^{\top}x$$
  
s.t.  $b^{\top}x \le b_0$   
 $a(x) \le 0$   
 $x \in [x]$ 

Relaxation

min 
$$c^{\mathsf{T}}x$$
  
s.t.  $b^{\mathsf{T}}x \le b_0$   
 $x \in [x]$ 

Reduction e.g. Add cut, prune box

min 
$$c^{\top}x$$
  
s.t.  $b^{\top}x \le b_0$   
 $d^{\top}x \le d_0$   
 $a(x) \le 0$   
 $x \in [x]$ 

**Split** 

$$[x'] \cup [x''] = [x]$$

min 
$$c^{\top}x$$
  
s.t.  $b^{\top}x \le b_0$   
 $a(x) \le 0$   
 $x \in [x'] \subset [x]$ 

min 
$$c^{\mathsf{T}}x$$
  
s.t.  $b^{\mathsf{T}}x \le b_0$   
 $a(x) \le 0$   
 $x \in [x''] \subset [x]$ 







### **Model Reductions**



$$\min \quad c^{\top}x$$

s.t. 
$$b^{\top}x \le b_0$$
  
 $a(x) \le 0$   
 $x \in [x]$ 

#### Relaxation

min 
$$c^{\top}x$$
  
s.t.  $b^{\top}x \le b_0$   
 $x \in [x]$ 

#### **Split**

$$[x'] \cup [x''] = [x]$$

min 
$$c^{\top}x$$
  
s.t.  $b^{\top}x \le b_0$   
 $a(x) \le 0$   
 $x \in [x'] \subset [x]$ 

min 
$$c^{\mathsf{T}}x$$

s.t. 
$$b^{\top}x \leq b_0$$
  
 $a(x) \leq 0$ 

$$x \in [x''] \subset [x]$$







# **Model Glueing**



min 
$$c^{\top}x$$
  
s.t.  $b^{\top}x \le b_0$   
 $a(x) \le 0$   
 $x \in [x]$ 

### Glue, Unsplit

$$[x'] \cup [x''] = [x]$$

min 
$$c^{\top}x$$
  
s.t.  $b^{\top}x \le b_0$   
 $a(x) \le 0$   
 $x \in [x'] \subset [x]$ 

min 
$$c^{\top}x$$
  
s.t.  $b^{\top}x \le b_0$   
 $a(x) \le 0$   
 $x \in [x''] \subset [x]$ 







### **Internal Representation**



- Models are organized in the search graph, represented by a Directed Acyclic Graph (DAG).
- For every model in the search graph the following information is stored:
  - Every equation/inequality is assigned a number of annotations describing its properties (e.g. linear, quadr., convex, separable, redundant, ...).
  - Additional local information (e.g. local optima, active constraints, Lagrangian multipliers,...) is added.
  - A description of the relation between the problem and its parent is provided.







### Search graph implementation



- The DAG is implemented using the VGTL, a library following the generic programming spirit of the C++ STL.
- There are two types of nodes:
  - Full nodes contain complete descriptions of models,
  - Delta nodes contain only the changes to the parent model in order to save storage capacity
- The search graph has a focus pointing to the model which is worked upon. This model is copied into an enhanced structure – the work node. A reference to this work node is passed to the inference engines.
- The graph itself can be analyzed by search inspectors.





# Internal Mathematical Representation



The internal mathematical representation of a problem is

min 
$$f_{lin}(x) + f_{quad}(x) + f_{sep}(x) + f_{0}(x)$$
  
s.t.  $G_{lin}(x) + G_{quad}(x) + G_{sep}(x) + G_{0}(x) \in S_{c}$   
 $x \in S_{v}$ 

where (currently) the sets S are boxes.

- The algorithmic representation is in graph form using not a tree (or forest) as usual but a directed acyclic graph.
- Variables appearing left of an assignment are substituted out.



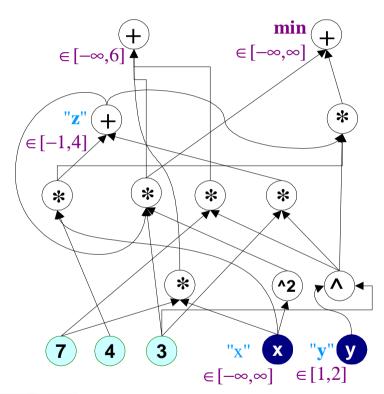


# **Directed Acyclic Graph (DAG)**



#### **Constraints**

**Objective** 



DAG representation of the model

min 
$$3x^2z + 4xy^3z$$
  
s.t.  $z = 4x + 3y^3$   
 $7x + 3x^2z + 7y^3 \le 6$   
 $y \in [1, 2], z \in [-1, 4]$ 

- similar to computational tree
- every node is an expression
- a node may have more than one parent
- Constants and variables are sources, objective and constraints are sinks







### **Expressions**



- Every vertex represents a function F mapping a vector  $x \in \mathbb{R}^n$  to a value  $F(x) \in \mathbb{R}$ .
- Predefined functions include sum, product, max, min, elementary real functions (exp, log, pow, sqrt, ...)
- Variable indicator contains the indices of the variables this vertex depends on.
- Additional information is added (ranges, semantics, variable name, vertex number, ...)







#### **Evaluation of a DAG**



- Evaluation works similar to computation trees by performing a graph walk.
- Caching keeps evaluation work minimal.
- The whole model is stored in one graph.
- Defining short-cuts makes it possible to replace graph walks by evaluation functions. Short-cuts may be defined at every node.
- Additional elementary functions can easily be incorporated.





### **Graphs and Evaluators**



- Generic Graph Library (Vienna Graph Template Library) in C++ to construct and manipulate DAGs, and forests (trees).
- Generic Programming approach with containers, walkers, function objects, and generic algorithms.
- For expression graphs (DAG or tree) special visitors are provided — (cached) forward and backward evaluators.
- Currently implemented Evaluators:
  - Real Function Values and Function Ranges
  - Gradients (Real, Interval)
  - Slopes
- In the near future Evaluators for:
  - Hessians (Real, Interval)
  - Second order Slopes



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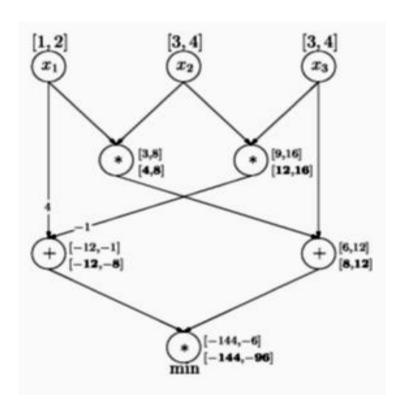


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# **Example**





- Interval evaluation and constraint propagation produce bounds on each node
- No reduction on the domain of the variables
- The bounds on intermediate nodes are improved compared to interval evaluation

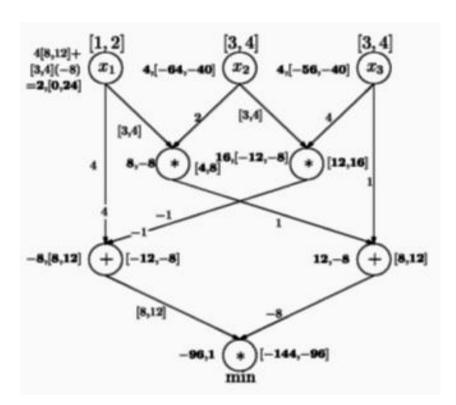






### **Example (ctd.)**





 Linear enclosures produced using slopes give redundant constraints, e.g.

$$24(x_1-2)-48(x_2-4)-32(x_3-4) \le 0$$



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### **Example (ctd.)**



Now constraint propagation leads to a reduction of the domain of the variables

$$x_2 \in [3.4, 4]$$
  
 $x_3 \in [3.4, 4]$ 

 With previously known techniques but without (expensive) higher order consistency, such a reduction would have required a split of the box.







### **Inference Engines**



- Corresponding to every type of problem change, a class of inference engines is designed:
  - Model analysis (e.g. find convex part)
  - Model reduction (e.g. pruning, fathoming)
  - Model relaxation (e.g. linear relaxation)
  - Model splitting (e.g. bisection)

- Model glueing (e.g. undo excessive splitting)
- Update local information (e.g. probing, local optimization)
- Check certificate (check correctness of calculation)
- Inference engines never change the model but calculate which changes may be made and are considered useful.







### Inference Engines: General features



- All inference modules only advertise changes.
- There is a fixed documentation structure defined.
  - Services Provided
  - Limits
  - Structure, Prerequisites of Input
  - Structure, Features of Output
  - Control Parameters
  - Termination Reason
- They produce a result where every possible change is listed together with a weight (the higher the weight the more important the change) and a certificate for correctness.
- They collect statistical data to support the strategy engine in making decisions.





# Inference Engines implemented as State of the Art



- Several state of the art techniques were implemented as inference engines:
  - STOP (starting point generator)
  - DONLP2-INTV (local optimizer)
  - Karush-John-Condition Generator
  - Point Verifier
  - Exclusion Box
  - Interval constraint propagation
  - Linear Relaxation
  - CPLEX (linear programming solver)
  - Basic Splitter
  - BCS (box covering solver)





### **Inference Engine: STOP**



- Heuristic Global Optimization Algorithm
- Combines Multi-Level-Coordinate-Search and Constraint Propagation
- Produces Starting Points for Local Optimization





### **Inference Engine: DONLP2-INTV**



- General purpose non-linear local optimizer for continuous variables
- SQP method
- Dense Linear Algebra
- Envelope uses standard evaluators, gradients are computed by automatic differentiation







### Inference Engine: Karush-John Conditions



- Generates the DAG representation of the Karush-John first order optimality conditions
- Every constraint (even two-sided) gets associated one Lagrange multiplier
- Constructed by symbolic differentiation of the DAG representation







### **Inference Engine: Point Verifier**



- Computes a uniqueness region around an approximate solution, in particular a verified point
- Uses Karush-John conditions
- Tries to maximize the uniqueness region by inflation





### **Inference Engine: Exclusion Box**



- Derives a large exclusion box and a tiny inclusion box such that the area between these two boxes does not contain a local optimizer.
- They are computed around an approximate local optimizer to get rid of the cluster effect.
- Does not focus on uniqueness.

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Uses slopes and H-matrix techniques.



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# Inference Engine: Constraint Propagation



- Performs the hull-consistency algorithm for constraint propagation.
- Reduces the possible range of the variables
- Might detect infeasibility of the problem





# **Inference Engine: Linear Relaxation**



- Computes a linear relaxation of the problem.
- Uses centered forms and slopes to compute the linear inequalities.
- Makes use of the DAG enhancements to improve the slopes.
- Either adds the linear relaxation as cuts or generates a full linear model.





### **Inference Engine: CPLEX**



- Solves linear problems.
- Interfaces the state-of-the-art commercial linear solver CPLEX.
- Extremely good performance.





### Inference Engine: Basic Split



- Provides splitting coordinates and split points.
- Computes a difficulty estimate for the variables involved.
- Suggests splits for the *n* most difficult variables.
- Uses exclusion box and solution information to improve the choice of split points.
- Cuts exclusion boxes out of the search area by careful choice of splitting coordinates.







### **Inference Engine: BCS**



- Covers the feasible area by boxes.
- Uses DMBC (dichotomous maintaining bound-consistency) and UCA6 (unionconservative approximation) in both basic and enhanced variants.
- Distinguishes between boxes in the interior and at the border of the feasible region.
- Uses the commercial ILOG Solver, or the constraint propagator provided by IRIN, but can work with any constraint propagator





# Contributions from the outside of the COCONUT project



We are happy that researchers and companies from outside the COCONUT project agreed to complement our efforts in integrating the known techniques:

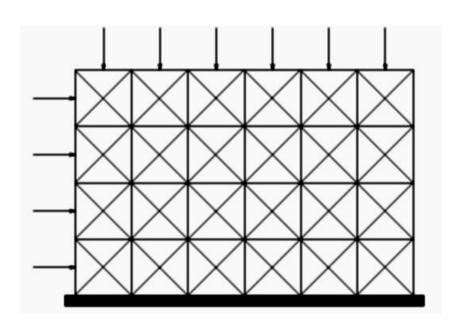
- Bernstein modules by J. Garloff & A. Smith (U. Konstanz)
- Verified lower bounds for convex relaxations by Ch. Jansson (TU Hamburg-Harburg)
- GAMS reader by the GAMS consortium
- Taylor arithmetic by G. Corliss (Marquette U.)
- Asymptotic arithmetic by K. Petras (U. Braunschweig)
- XPRESS commercial LP-solver (Dash Optimization)
- Hopefully additional contributions by you!





# Worst case finite element analysis





- Linear FEM equations become non-convex if material data is uncertain.
- Typical size of uncertainty is 10-20% in elasticity and cross-section area.
- Law requires the computation of the worst case.
- Industry relevant problems have some thousand variables.



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# Worst case FEM structural analysis Promising result



Worst case analysis on the displacements u for a 20x20 wall in the non-linear system

$$A(x)u = b$$

- 1620 material parameters x with 16.4% uncertainty, 840 displacements u
- Traditional methods fail for 0.01% uncertainty
- Exploiting the special structure, within 30s on a 1.6 Ghz
   Pentium 4, without bisection we get

Uncertainty (%)	0.01	0.05	0.5	1	2.5	5	10	16.4
Overestimation	1.03	1.15	2.55	4.12	8.92	17.26	35.33	61.59

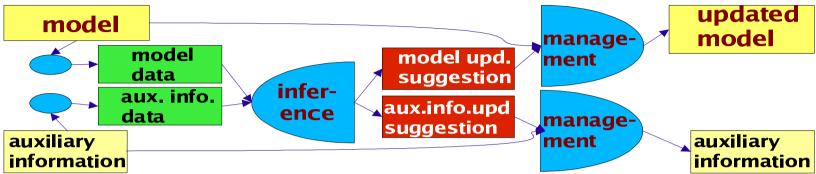






### **Basic algorithm design**





- This setup allows for highest flexibility and extensibility
  - the modules are split into inference engines (calculation) and management parts
  - additional modules for model handling are added
- The strategy engine decides which components are called in every algorithmic step of this type





### **Report Modules**



- This class of modules produces output.
  Various types of files and human readable output will have to be created.
- Examples:
  - Solution Report (humans, AMPL, GAMS)
  - Progress Information
  - Checkpointing
  - Debugging Information
  - Error Messages







### **Management Modules**



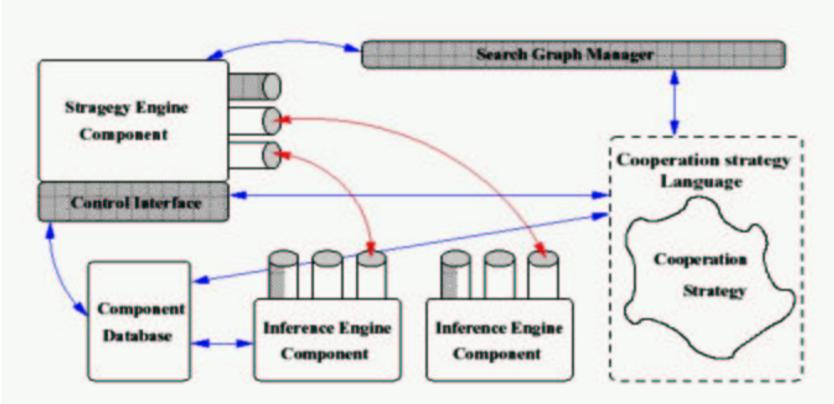
- Corresponding to every internal part of the program, a class of management modules is designed:
  - Model management
  - Data collection
  - Resource management
  - Initialization management
- Management modules never calculate anything. They just perform some of the changes which have been advertised by inference modules.





# **Strategy Engine**







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# **Strategy Engine (ctd.)**



- It is the core of the algorithm and consists of
  - The logic core ("search") which is essentially the main solution loop,
  - Special decision makers (very specialized inference engines) for determining the next action at every point in the algorithm.
- It calls the management modules, the report modules, and the inference engines in succession.
- It can be programmed using a simple strategy language (interpreted, Python based).
  - (Semi-)interactive and automatic solution process
  - Debugging and single-stepping of strategies
  - Object oriented, dynamically typed objects, garbage collected
  - Easily extendable







# **Strategy Engine (ctd.)**



- Manages the search graph via the search graph manager,
- Manages the search database via the database manager,
- Uses a component framework to communicate with the inference engines,
- Launches inference engines dynamically (on need) to avoid memory overload,
- Provides a management interface,
- Strategy engine is itself a component, so multilevel strategies are possible,
- Prepared for distributed and parallel computing, and distributed memory



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### **Extensibility**



- The strategy language makes it easy to change the strategy.
- A registration phase during initialization removes the need to recompile the program when new inference engines are added.
- Registration also allows us to balance scientific and commercial interests:
  - Free but reduced core version with open API specification
  - Free strategy engine with basic strategy
  - Advanced commercial components

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Extending the system by external contributers is made easy by this modular design.



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#### **Invitation**



# We hope that the community will contribute to this algorithmic framework.



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#### The End



# Thank you for your attention!

**COCONUT** Website:

http://www.mat.univie.ac.at/coconut



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