

A NON-AUTOMATIC (!) APPLICATION OF GOSPER'S ALGORITHM EVALUATES A DETERMINANT FROM TILING ENUMERATION

M. CIUCU AND C. KRATTENTHALER

ABSTRACT. We evaluate the determinant $\det_{1 \leq i, j \leq n} \times \left(\binom{x+y+j}{x-i+2j} - \binom{x+y+j}{x+i+2j} \right)$, which gives the number of lozenge tilings of a hexagon with cut off corners. A particularly interesting feature of this evaluation is that it requires the proof of a certain hypergeometric identity which we accomplish by using Gosper's algorithm in a nonautomatic fashion.

The purpose of this paper is to provide a direct evaluation of the determinant

$$(1) \quad \det_{1 \leq i, j \leq n} \left(\binom{x+y+j}{x-i+2j} - \binom{x+y+j}{x+i+2j} \right).$$

This determinant arises in our study [4] on the enumeration of lozenge tilings of hexagons with cut off corners. For example, consider a hexagon with side lengths $x+n$, n , y , $x+n$, n , y , in cyclic order, and angles of 120° of which two adjacent corners are cut off as in Figure 1(a).¹ Figure 1(b) shows a lozenge tiling of this region, by which we mean a tiling by unit rhombi with angles of 60° and 120° , referred to as lozenges. The number of these lozenge tilings is given by the determinant (1). This is seen by converting the lozenge tilings into families (P_1, P_2, \dots, P_n) of nonintersecting lattice paths consisting of

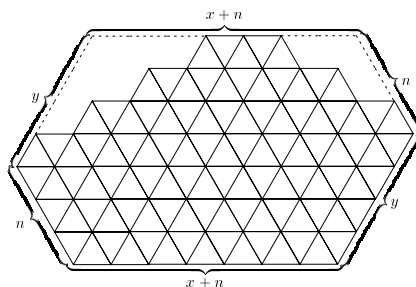
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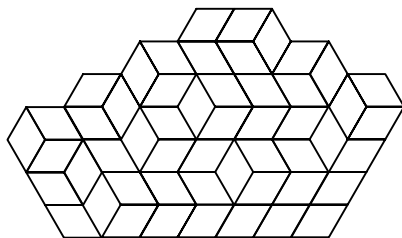
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¹ To be precise, from the top-left corner we cut off a (reversed) staircase of the form $(y-1, y-2, \dots, 1)$, meaning that the cut-off staircase consists of $y-1$ rhombi in the first row, $y-2$ rhombi in the second row, etc., and from the top-right corner we cut off a staircase of the form $(n-1, n-2, \dots, 1)$.



(a) A hexagon with cut off corners.



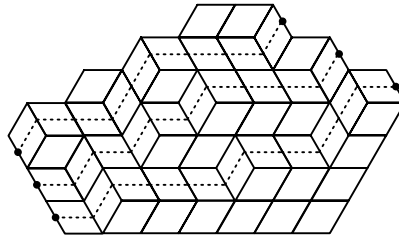
(b) A lozenge tiling of the hexagon with cut off corners.

Figure 1.

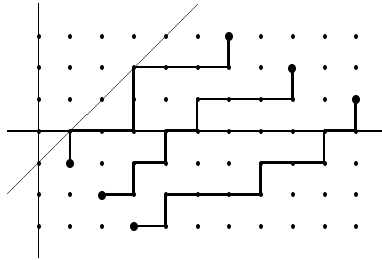
positive unit steps, where the path P_i runs from $(i, -i)$ to $(x+2i, y-i)$, $i = 1, 2, \dots, n$ and does not cross the diagonal $y = x - 1$, (see Figure 2), and then applying the main theorem of nonintersecting lattice paths [18, Lemma 1], [8], [23, Theorem 1.2], (see [4] for details and background; there is also another case in [4] in which the determinant (1) provides the solution).

By Theorem 1 below, the number of the lozenge tilings of the preceding paragraph is given by a closed form expression. The proof of Theorem 1 that we present in this paper² is primarily based on hypergeometric series identities. A remarkable aspect is that it contains an instance of a *nonautomatic* application of Gosper's algorithm [9] (see also [10 Section 5.7], [20 Section II.7]), see Step 3 of proof of Theorem 1. This is noteworthy because Gosper invented his algorithm to *automate* summation, so that a nonautomatic application must be almost considered as a misuse. But clearly, and more seriously, the fact that Gosper's algorithm is also useful in "computer-free territory"

² An alternative proof is presented in [4], in which a combinatorial argument is used to convert the determinant (1) into a different determinant that was already known from [12, Theorem 10].



(a) The path family corresponding to a lozenge tiling.



(b) The paths made orthogonal.

FIGURE 2.

only adds to its value. (The only other instance of a nonautomatic application of Gosper's algorithm that we are aware of appears in [19]. However, the purpose of use there is different. Roughly speaking, we use it to prove a *positive* result, namely to verify the truth of an identity between certain hypergeometric series, (see (14)). In contrast, Petkovšek and Wilf use it to prove a *negative* result, namely that a certain binomial sum cannot be expressed in terms of closed form expressions.)

Theorem 1. *Let n be a positive integer, and let x and y be nonnegative integers. Then the following determinant evaluation holds:*

$$\begin{aligned}
 (2) \quad & \det_{1 \leq i, j \leq n} \left(\binom{x+y+j}{x-i+2j} - \binom{x+y+j}{x+i+2j} \right) \\
 &= \prod_{j=1}^n \frac{(j-1)! (x+y+2j)! (x-y+2j+1)_j (x+2y+3j+1)_{n-j}}{(x+n+2j)! (y+n-j)!},
 \end{aligned}$$

where the shifted factorial $(a)_k$ is defined by $(a)_k := a(a+1)\cdots(a+k-1)$, $k \geq 1$, and $(a)_0 := 1$.

Remark. We formulate Theorem 1 only for integral x and y . But in fact, with a generalized definition of factorials and binomials, cf. [10, Section 5.5, (5.96), (5.100)], Theorem 1 would also make sense and be true for complex x and y .

Proof. We prove the determinant evaluation by “identification of factors,” a method that is also applied successfully in [2, 3, 5, 6, 7, 11–14, 16, 17, 21], (see in particular the tutorial description in [15, Section 2.4] or [13, Section 2]).

First of all, we take appropriate factors out of the determinant. To be precise, we take $(x + y + j)! / ((x + n + 2j)! (y + n - j)!)$ out of the j th column of the determinant in (2), $j = 1, 2, \dots, n$. Thus we obtain

$$(3) \quad \prod_{i=1}^n \frac{(x + y + j)!}{(x + n + 2j)! (y + n - j)!} \times \det_{1 \leq i, j \leq n} ((x + 2j - i + 1)_{n+i} (y + i - j + 1)_{n-i} - (x + 2j + i + 1)_{n-i} (y - i - j + 1)_{n+i})$$

for the determinant in (2). Let us denote the determinant in (3) by $D_n(x, y)$. Comparison of (2) and (3) yields that (2) will be proved once we are able to establish the determinant evaluation

$$(4) \quad \begin{aligned} D_n(x, y) &= \det_{1 \leq i, j \leq n} ((x + 2j - i + 1)_{n+i} (y + i - j + 1)_{n-i} - (x + 2j + i + 1)_{n-i} (y - i - j + 1)_{n+i}) \\ &= \prod_{j=1}^n (j - 1)! (x + y + j + 1)_j (x - y + 2j + 1)_j \\ &\quad \times (x + 2y + 3j + 1)_{n-j}. \end{aligned}$$

For the proof of (4) we proceed in several steps. An outline is as follows. In the first step we show that $\prod_{j=1}^n (x - y + 2j + 1)_j$ is a factor of $D_n(x, y)$ as a polynomial in x and y . In the second step we show that $\prod_{j=1}^n (x + y + j + 1)_j$ is a factor of $D_n(x, y)$, and in the third step we show that $\prod_{j=1}^n (x + 2y + 3j + 1)_{n-j}$ is a factor of $D_n(x, y)$. Then, in the fourth step we determine the maximal degree of $D_n(x, y)$ as a polynomial in x , and the maximal degree as a polynomial in y , which

turns out to be $n(3n + 1)/2$ in both cases. On the other hand, the degree in x , and also in y , of the product on the righthand side of (4), which by the first three steps divides $D_n(x, y)$, is exactly $n(3n + 1)/2$. Therefore we are forced to conclude that

$$(5) \quad D_n(x, y) = C(n) \prod_{j=1}^n (x - y + 2j + 1)_j \times (x + y + j + 1)_j (x + 2y + 3j + 1)_{n-j},$$

where $C(n)$ is a constant independent of x and y . Finally, in the fifth step, we determine the constant $C(n)$, which turns out to equal $\prod_{j=1}^n (j - 1)!$. Clearly, this would finish the proof of (4), and thus of (2), as we already noted.

Step 1. $\prod_{j=1}^n (x - y + 2j + 1)_j$ is a factor of $D_n(x, y)$. Let us concentrate on a typical factor $(x - y + 2j + l)$, $1 \leq j \leq n$, $1 \leq l \leq j$. We claim that for each such factor there is a linear combination of the columns that vanishes if the factor vanishes. More precisely, we claim that for any j, l with $1 \leq j \leq n$, $1 \leq l \leq j$, the following holds

$$(6) \quad \sum_{s=l}^{\lfloor \frac{j+l}{2} \rfloor} \frac{(j-l)(j+l-2s+1)_{s-l}}{(j-s)(s-l)!} \times \frac{(x+2j+l+n-s+1)_{s-l}(x+n+2s+1)_{j+l-2s}}{(2x+2j+l+s+1)_{j-s}} \times (\text{column } s \text{ of } D_n(x, x+2j+l)) + (\text{column } j \text{ of } D_n(x, x+2j+l)) = 0.$$

To avoid confusion, for $j = l$ it is understood by convention that the sum in (6) vanishes.

In order to verify (6), we have to check

$$\begin{aligned}
 (7) \quad & \sum_{s=l}^{\lfloor \frac{j+l}{2} \rfloor} \frac{(j-l)(j+l-2s+1)_{s-l}}{(j-s)(s-l)!} \\
 & \cdot \frac{(x+2j+l+n-s+1)_{s-l}(x+n+2s+1)_{j+l-2s}}{(2x+2j+l+s+1)_{j-s}} \\
 & \cdot ((x+i+2j+l-s+1)_{n-i}(x-i+2s+1)_{n+i} \\
 & \quad - (x-i+2j+l-s+1)_{n+i}(x+i+2s+1)_{n-i}) \\
 & \quad + (x-i+2j+1)_{n+i}(x+i+j+l+1)_{n-i} \\
 & \quad - (x+i+2j+1)_{n-i}(x-i+j+l+1)_{n+i} = 0,
 \end{aligned}$$

which is (6) restricted to the i th row. The exceptional case $j = l$ can be treated immediately. By assumption, the sum in (7) vanishes for $j = l$, and, by inspection, also the other two expressions in (7) vanish for $j = l$. So it remains to establish (7) for $j > l$. In terms of the standard hypergeometric notation

$${}_rF_s \left[\begin{matrix} a_1, \dots, a_r \\ b_1, \dots, b_s \end{matrix}; z \right] = \sum_{k=0}^{\infty} \frac{(a_1)_k \cdots (a_r)_k}{k! (b_1)_k \cdots (b_s)_k} z^k,$$

this means to check

$$\begin{aligned}
 (8) \quad & \frac{(x+i+2j+1)_{n-i}(x-i+2l+1)_{n+i+j-l}}{(2x+2j+2l+1)_{j-l}} \\
 & \times {}_4F_3 \left[\begin{matrix} -\frac{j}{2} + \frac{l}{2}, \frac{1}{2} - \frac{j}{2} + \frac{l}{2}, -i - 2j - x, 1 + 2j + 2l + 2x \\ 1 - j + l, \frac{1}{2} - \frac{i}{2} + l + \frac{x}{2}, 1 - \frac{i}{2} + l + \frac{x}{2} \end{matrix}; 1 \right] \\
 & - \frac{(x-i+2j+1)_{n+i}(x+i+2l+1)_{-i+j-l+n}}{(2x+2j+2l+1)_{j-l}} \\
 & \times {}_4F_3 \left[\begin{matrix} -\frac{j}{2} + \frac{l}{2}, \frac{1}{2} - \frac{j}{2} + \frac{l}{2}, i - 2j - x, 1 + 2j + 2l + 2x \\ 1 - j + l, \frac{1}{2} + \frac{i}{2} + l + \frac{x}{2}, 1 + \frac{i}{2} + l + \frac{x}{2} \end{matrix}; 1 \right] \\
 & + (x-i+2j+1)_{n+i}(x+i+j+l+1)_{n-i} \\
 & - (x+i+2j+1)_{n-i}(x-i+j+l+1)_{n+i} = 0.
 \end{aligned}$$

Both ${}_4F_3$ -series can be summed by means of a ${}_4F_3$ -summation which appears in a paper by Andrews and Burge [1, Lemma 1], (see [12, Lemma A3] for a simpler proof),

$${}_4F_3 \left[\begin{matrix} -\frac{N}{2}, \frac{1}{2} - \frac{N}{2}, -a, a + b \\ 1 - N, \frac{b}{2}, \frac{1}{2} + \frac{b}{2} \end{matrix}; 1 \right] = \frac{(a+b)_N}{(b)_N} + \frac{(-a)_N}{(b)_N},$$

where N is a positive integer. We have to apply the case where $N = j - l$. This is indeed a positive integer because of our assumption $j > l$. Some simplification then leads to (8).

This shows that $\prod_{j=1}^n (x - y + 2j + 1)_j$ divides $D_n(x, y)$.

Step 2. $\prod_{j=1}^n (x + y + j + 1)_j$ is a factor of $D_n(x, y)$. Let us concentrate on a typical factor $(x + y + j + l)$, $1 \leq j \leq n$, $1 \leq l \leq j$. We claim that for each such factor there is a linear combination of the columns that vanishes if the factor vanishes. More precisely, we claim that for any j, l with $1 \leq j \leq n$, $1 \leq l \leq j$ the following holds

$$(9) \quad \sum_{s=1+j-l}^j \left(-\frac{1}{4}\right)^{j-s} \binom{l-1}{s+l-j-1} \cdot \frac{(x+n+2s+1)_{2j-2s} (2x+3j+l+s+1)_{j-s}}{(x+j+s+\frac{1}{2})_{j-s} (x+j+l+s)_{j-s} (x+j+l-n+s)_{j-s}} \cdot (\text{column } s \text{ of } D_n(x, -x-j-l)) = 0.$$

In order to verify (9), we have to check

$$\sum_{s=1+j-l}^j \left(-\frac{1}{4}\right)^{j-s} \binom{l-1}{s+l-j-1} \cdot \frac{(x+n+2s+1)_{2j-2s} (2x+3j+l+s+1)_{j-s}}{(x+j+s+\frac{1}{2})_{j-s} (x+j+l+s)_{j-s} (x+j+l-n+s)_{j-s}} \cdot ((-x+i-j-l-s+1)_{n-i} (x-i+2s+1)_{n+i} - (-x-i-j-l-s+1)_{n+i} (x+i+2s+1)_{n-i}) = 0,$$

which is (9) restricted to the i -th row. Equivalently, using hypergeometric notation, this means to check

$$(10) \quad (-1)^l \frac{(-i-2j-x)_{n+i} (3+i+2j-2l+x)_{-2-i+2l+n} (2+4j+2x)_{l-1}}{4^{l-1} (1+2j+x)_{l-1} (\frac{3}{2}+2j-l+x)_{l-1} (1+2j-n+x)_{l-1}} \times {}_4F_3 \left[\begin{matrix} 1-l, \frac{3}{2}+2j-l+x, 1+2j+x, 1+i+2j+x \\ 2+4j+2x, 2+\frac{i}{2}+j-l+\frac{x}{2}, \frac{3}{2}+\frac{i}{2}+j-l+\frac{x}{2} \end{matrix} ; 1 \right] - (-1)^l \frac{(i-2j-x)_{n-i} (3-i+2j-2l+x)_{-2+i+2l+n} (2+4j+2x)_{l-1}}{4^{l-1} (1+2j+x)_{l-1} (\frac{3}{2}+2j-l+x)_{l-1} (1+2j-n+x)_{l-1}} \times {}_4F_3 \left[\begin{matrix} \frac{3}{2}+2j-l+x, 1+2j+x, 1-i+2j+x, 1-l \\ \frac{3}{2}-\frac{i}{2}+j-l+\frac{x}{2}, 2-\frac{i}{2}+j-l+\frac{x}{2}, 2+4j+2x \end{matrix} ; 1 \right] = 0.$$

In order to establish (10) we apply Bailey’s transformation for balanced ${}_4F_3$ -series, (see [22, (4.3.5.1)]),

$$\begin{aligned}
 & {}_4F_3 \left[\begin{matrix} a, b, c, -N \\ e, f, 1 + a + b + c - e - f - N \end{matrix}; 1 \right] = \frac{(e - a)_N (f - a)_N}{(e)_N (f)_N} \\
 & \times {}_4F_3 \left[\begin{matrix} -N, a, 1 + a + c - e - f - N, 1 + a + b - e - f - N \\ 1 + a + b + c - e - f - N, 1 + a - e - N, 1 + a - f - N \end{matrix}; 1 \right],
 \end{aligned}$$

where N is a nonnegative integer, to the second ${}_4F_3$ -series in (10). Thus it is converted into the first ${}_4F_3$ -series, and it is routine to check that also the remaining terms that go with the ${}_4F_3$ -series agree. So, the two terms on the lefthand side of (10) cancel each other, as desired.

This establishes that $\prod_{j=1}^n (x + y + j + 1)_j$ divides $D_n(x, y)$.

Step 3. $\prod_{i=1}^n (x + 2y + 3i + 1)_{n-i}$ is a factor of $D_n(x, y)$. This is the most difficult part of the proof of (4). Trials of finding linear combinations of columns that vanish resulted in extremely messy expressions. So we decided to work with linear combinations of rows this time. Still, the coefficients are not as “nice” as in Steps 1 and 2.

Let us concentrate on a typical factor $(x + 2y + 3i + l)$, $1 \leq i \leq n$, $1 \leq l \leq n - i$. We claim that for each such factor there is a linear combination of the rows that vanishes if the factor vanishes. More precisely, we claim that for any i, l with $1 \leq i \leq n$, $1 \leq l \leq n - i$ there holds

$$\begin{aligned}
 (11) \quad & \sum_{k=1}^{i+l} \frac{(k + i + l + 1)_{i+l-k}}{(i + l - k)!} P_l(2i, i + l - k) \\
 & \cdot (\text{row } k \text{ of } D_n(-2y - 3i - l, y)) = 0,
 \end{aligned}$$

where $P_l(e, f)$ is the polynomial

$$(12) \quad P_l(e, f) = \sum_{r=0}^{2l+1} a_r (e)_r (-f)_{2l+1-r},$$

with the expansion coefficients a_r given by

$$(13) \quad a_r = \langle x^r \rangle ((x^2 + x + 1)^{l-1} (2x + 1)(x + 2)(x - 1)).$$

Here, $\langle x^r \rangle g(x)$ denotes the coefficient of x^r in $g(x)$.

By specializing (11) to the j th column, splitting the resulting sum into two parts in the obvious way, and then moving one sum to the righthand side, we see that, in order to verify (11), we have to check

$$\begin{aligned} & \sum_{k=1}^{i+l} \frac{(k+i+l+1)_{i+l-k}}{(i+l-k)!} P_l(2i, i+l-k) \\ & \quad \times (-2y-3i-l+2j-k+1)_{n+k} (y+k-j+1)_{n-k} \\ = & \sum_{k=1}^{i+l} \frac{(k+i+l+1)_{i+l-k}}{(i+l-k)!} P_l(2i, i+l-k) \\ & \quad \times (-2y-3i-l+2j+k+1)_{n-k} (y-k-j+1)_{n+k}, \end{aligned}$$

or, after adding one more term as first summand on both sides, equivalently,

$$\begin{aligned} & \sum_{k=0}^{i+l} \frac{(k+i+l+1)_{i+l-k}}{(i+l-k)!} P_l(2i, i+l-k) \\ & \quad \times (-2y-3i-l+2j-k+1)_{n+k} (y+k-j+1)_{n-k} \\ (14) \quad & = \sum_{k=0}^{i+l} \frac{(k+i+l+1)_{i+l-k}}{(i+l-k)!} P_l(2i, i+l-k) \\ & \quad \times (-2y-3i-l+2j+k+1)_{n-k} (y-k-j+1)_{n+k}. \end{aligned}$$

Empirically, we discovered that apparently both sums in (14) are indefinitely summable (“Gosper-summable”; (see [10, Section 5.7], [20, Section II.5])). It is exactly this fact which makes (14) tractable.

In the following we will show that the sums in (14) are equal, however, without exhibiting an explicit expression for the sums. Instead, what we will do is to read through Gosper’s algorithm [9], (see [10, Section 5.7], [20, Section II.5]), which is an algorithm that solves the problem of indefinite summation for hypergeometric sums. (For any fixed l , our sums in (14) belong to the category of hypergeometric sums.) In the course of reading through Gosper’s algorithm it will emerge that the sums on both sides of (14) must be equal.

Let us recall what Gosper’s algorithm does and how it works. Let $t(k)$ be a “hypergeometric term,” i.e., be a term such that the ratio

$t(k + 1)/t(k)$ is a rational function in k . Then the Gosper algorithm will find a hypergeometric term $T(k)$, (if it exists), satisfying

$$(15) \quad t(k) = T(k + 1) - T(k).$$

The upshot of this is that then the indefinite summation of the term $t(k)$ can be easily carried out,

$$(16) \quad \sum_{k=A}^B t(k) = T(B + 1) - T(A).$$

The term $T(k)$ is found in the following way. First, one finds polynomials $p(k)$, $q(k)$, and $r(k)$ such that

$$(17) \quad \frac{t(k + 1)}{t(k)} = \frac{p(k + 1)}{p(k)} \frac{q(k)}{r(k + 1)},$$

where $q(k)$ and $r(k)$ have the property that whenever $(k + \alpha) \mid q(k)$ and $(k + \beta) \mid r(k)$ then the difference $\alpha - \beta$ must not be a positive integer. Next, one finds a polynomial $s(k)$ satisfying the recurrence relation

$$(18) \quad p(k) = q(k)s(k + 1) - r(k)s(k)$$

for all k . The term $T(k)$ is then given by

$$(19) \quad T(k) = \frac{r(k)s(k)}{p(k)}t(k).$$

Now let us carry out this program with the summands in (14). First, let $t(k) = t_1(k)$, where $t_1(k)$ is the summand of the sum on the left hand side of (14),

$$t_1(k) = \frac{(k + i + l + 1)_{i+l-k}}{(i + l - k)!} P_l(2i, i + l - k) \times (-2y - 3i - l + 2j - k + 1)_{n+k} (y + k - j + 1)_{n-k}.$$

Then (17) holds with $p(k) = p_1(k)$, $q(k) = q_1(k)$, $r(k) = r_1(k)$, where $p_1(k) = P_l(2i, i + l - k)$, $q_1(k) = (i + l - k)(-2y - 3i - l + 2j - k)$, and

$r_1(k) = (i + l + k)(y - j + k)$. So, next we have to find a polynomial $s_1(k)$ satisfying the recurrence

$$(20) \quad P_l(2i, i + l - k) = (i + l - k)(-2y - 3i - l + 2j - k)s_1(k + 1) \\ - (i + l + k)(y - j + k)s_1(k).$$

For each specific instance of i and l this is just routine. However, we were not able to find an explicit formula for $s_1(k)$ in general. Fortunately, we do not need such an explicit expression. Assuming that we have found a polynomial $s_1(k)$ satisfying (20), by (16) and (19) we have

$$(21) \quad \sum_{k=0}^{i+l} \frac{(k + i + l + 1)_{i+l-k}}{(i + l - k)!} P_l(2i, i + l - k) \\ \times (-2y - 3i - l + 2j - k + 1)_{n+k} (y + k - j + 1)_{n-k} \\ = \frac{r_1(i + l + 1) s_1(i + l + 1)}{p_1(i + l + 1)} t_1(i + l + 1) - \frac{r_1(0) s_1(0)}{p_1(0)} t_1(0) \\ = -\frac{(i + l)_{i+l+1}}{(i + l)!} (-2y - 3i - l + 2j + 1)_n (y - j)_{n+1} s_1(0),$$

the last line being due to the fact that $t_1(i + l + 1) = 0$.

On the other hand, for $t(k) = t_2(k)$, where $t_2(k)$ is the summand of the sum on the righthand side of (14),

$$t_2(k) = \frac{(k + i + l + 1)_{i+l-k}}{(i + l - k)!} P_l(2i, i + l - k) \\ \times (-2y - 3i - l + 2j + k + 1)_{n-k} (y - k - j + 1)_{n+k},$$

we may choose $p(k) = p_2(k)$, $q(k) = q_2(k)$, $r(k) = r_2(k)$, where $p_2(k) = P_l(2i, i + l - k)$, $q_2(k) = (i + l - k)(y - j - k)$, and $r_2(k) = (i + l + k)(-2y - 3i - l + 2j + k)$. So here we have to find a polynomial $s_2(k)$ satisfying the recurrence

$$(22) \quad P_l(2i, i + l - k) = (i + l - k)(y - j - k)s_2(k + 1) \\ - (i + l + k)(-2y - 3i - l + 2j + k)s_2(k).$$

Again, this is just routine for each specific instance of i and l , but we do not know an explicit formula for $s_2(k)$ in general. Assuming that

we have found a polynomial $s_2(k)$ satisfying (22), by (16) and (19) we have

$$\begin{aligned}
 & \sum_{k=0}^{i+l} \frac{(k+i+l+1)_{i+l-k}}{(i+l-k)!} P_l(2i, i+l-k) \\
 (23) \quad & \times (-2y-3i-l+2j+k+1)_{n-k} (y-k-j+1)_{n+k} \\
 & = \frac{r_2(i+l+1) s_2(i+l+1)}{p_2(i+l+1)} t_2(i+l+1) - \frac{r_2(0) s_2(0)}{p_2(0)} t_2(0) \\
 & = -\frac{(i+l)_{i+l+1}}{(i+l)!} (-2y-3i-l+2j)_{n+1} (y-j+1)_n s_2(0),
 \end{aligned}$$

the last line being due to the fact that also $t_2(i+l+1) = 0$.

In order to relate $s_2(k)$ to $s_1(k)$, we make the following observation. We set $s_2(k) = \tilde{s}_2(-k+1)$, substitute this in the recurrence (22), then replace k by $-k$ and change the sign on both sides of (22). Thus we obtain for $\tilde{s}_2(k)$ the recurrence

$$\begin{aligned}
 (24) \quad & -P_l(2i, i+l+k) = (i+l-k)(-2y-3i-l+2j-k)\tilde{s}_2(k+1) \\
 & - (i+l+k)(y-j+k)\tilde{s}_2(k).
 \end{aligned}$$

This is almost the same recurrence as the recurrence (20) for $s_1(k)$! It is only the term on the left hand side which is different! But, in fact, there is no difference: We claim that:

Claim 1. *We have $P_l(e, e+2l-f) = -P_l(e, f)$.*

Claim 2. *There exists a unique solution for the recurrence (20).*

Let us for the moment assume that these claims have already been established. Then, because of Claim 1, the recurrences (20) and (24) are indeed the same. Furthermore, thanks to Claim 2, there does exist a unique solution for the recurrence (20), and so also for (24). Hence, the solutions must be the same, i.e., $s_1(k) = \tilde{s}_2(k)$, which means $s_1(k) = s_2(1-k)$. In particular, we have $s_1(1) = s_2(0)$. A further fact, which follows immediately from Claim 1 on replacing e by $2e$ and setting $f = e+l$, is that $P_l(2e, e+l) = 0$. Therefore, by setting $k = 0$ in (20), we obtain

$$0 = (i+l)(-2y-3i-l+2j)s_1(1) - (i+l)(y-j)s_1(0).$$

From this equation, and the previous observation that $s_1(1) = s_2(0)$, we infer

$$s_1(0) = \frac{(-2y - 3i - l + 2j)}{(y - j)} s_1(1) = \frac{(-2y - 3i - l + 2j)}{(y - j)} s_2(0).$$

Substitution of this relation in (21) gives

$$\begin{aligned} & \sum_{k=0}^{i+l} \frac{(k + i + l + 1)_{i+l-k}}{(i + l - k)!} P_l(2i, i + l - k) \\ & \quad \times (-2y - 3i - l + 2j - k + 1)_{n+k} (y + k - j + 1)_{n-k} \\ & = -\frac{(i + l)_{i+l+1}}{(i + l)!} (-2y - 3i - l + 2j)_{n+1} (y - j + 1)_n s_2(0). \end{aligned}$$

Comparison of this identity with (23) shows that indeed the sums on both sides of (14) are identical. This would prove (14).

So it remains to settle Claims 1 and 2.

We begin with Claim 1. By the definition (12) of $P_l(e, f)$, we have

$$\begin{aligned} P_l(e, e + 2l - f) &= \sum_{r=0}^{2l+1} a_r(e)_r (-e - 2l + f)_{2l+1-r} \\ &= \sum_{r=0}^{2l+1} a_r(e)_r (-1)^{r+1} (e + r - f)_{2l+1-r}, \end{aligned}$$

where the coefficients a_r are given by (13). Next we use the Chu-Vandermonde summation, (see e.g., [10, Section 5.1, (5.27)]) in the form

$$\sum_{s=0}^N \binom{N}{s} (x)_s (y)_{N-s} = (x + y)_N,$$

with $N = 2l + 1 - r$, $x = e + r$, and $y = -f$. Thus,

$$\begin{aligned} P_l(e, e + 2l - f) &= \sum_{r=0}^{2l+1} a_r(e)_r (-1)^{r+1} \\ & \quad \times \sum_{s=0}^{2l+1-r} \binom{2l + 1 - r}{s} (e + r)_s (-f)_{2l+1-r-s} \\ &= -\sum_{m=0}^{2l+1} (e)_m (-f)_{2l+1-m} \sum_{r=0}^m \binom{2l + 1 - r}{m - r} (-1)^r a_r. \end{aligned}$$

Therefore, Claim 1 will follow immediately, if we are able to show that

$$(25) \quad \sum_{r=0}^m \binom{2l+1-r}{m-r} (-1)^r a_r = a_m.$$

This can be readily done by using generating functions. The definition (13) of the coefficients a_r is equivalent to

$$(26) \quad \sum_{r=0}^{\infty} a_r x^r = (x^2 + x + 1)^{l-1} (2x + 1)(x + 2)(x - 1).$$

Let us denote the righthand side of this equation by $A(x)$. Now we multiply both sides of (25) by x^m , and we sum over all $m = 0, 1, \dots$. We obtain

$$\sum_{m=0}^{\infty} \sum_{r=0}^m \binom{2l+1-r}{m-r} (-1)^r a_r = A(x),$$

and after interchanging summations on the lefthand side and summing the (now) inner sum over m by means of the binomial theorem,

$$(1+x)^{2l+1} A\left(-\frac{x}{1+x}\right) = A(x).$$

It is trivial to verify this equation. Thus also the equivalent equation (25) must be true. Due to the preceding considerations, this completes the proof of Claim 1.

Next we turn to Claim 2. We show that there is a unique polynomial $s_1(k)$ of degree $2l$ that satisfies the recurrence (20). (We leave it as an exercise that the “degree calculus” of the Gosper algorithm shows that if there is a solution to the recurrence (20) then it has to be a polynomial of degree at most $2l$.) So, let $s_1(k) = \sum_{m=0}^{2l} c(m)(k-i-l)_m$. We substitute this into (20), then expand everything with respect to the basis $(k-i-l)_m$, $m = 0, 1, \dots$, for the space of polynomials in k , and finally compare coefficients of $(k-i-l)_m$ on both sides of (20). This leads to the following system of equations for the coefficients $c(m)$:

$$(27) \quad \begin{aligned} a_{2l+1-m}(2i)_{2l+1-m} &= (y+i-l-j+m-1)c(m-1) \\ &\quad - (2i+2l-m)(y+i+l-j-m)c(m), \\ &\quad m = 0, 1, \dots, 2l+1, \end{aligned}$$

where, by convention, we put $c(-1) = c(2l + 1) = 0$. For convenience, we set

$$c(m) = (2i)_{2l-m} (y + i - l - j)_{2l-m} (y + i - l - j)_m \tilde{c}(m).$$

By substituting this in (27), we obtain the simpler system of equations

$$(28) \quad \frac{a_m}{(y + i - l - j)_{2l-m+1} (y + i - l - j)_m} = \tilde{c}(m - 1) - \tilde{c}(m),$$

$$m = 0, 1, \dots, 2l + 1.$$

This is a system of $2l + 2$ equations for $2l + 1$ variables. (Recall the convention $c(-1) = c(2l + 1) = 0$, which of course implies $\tilde{c}(-1) = \tilde{c}(2l + 1) = 0$.) So, it is overdetermined. It is easy to see that this inhomogeneous system of linear equations has a (unique) solution if and only if the sum of the lefthand sides of (28) over all m equals 0, i.e., if and only if

$$(29) \quad \sum_{m=0}^{2l+1} \frac{a_m}{(y + i - l - j)_{2l-m+1} (y + i - l - j)_m} = 0.$$

This would follow immediately from the antisymmetry property $a_m = -a_{2l+1-m}$, because then the m th and $(2l + 1 - m)$ th summand in the sum in (29) would cancel each other. Indeed, the substitution $x \rightarrow 1/x$ in (26) yields $a_m = -a_{2l+1-m}$. Therefore, the system of equations (27) has indeed a unique solution, which implies that there is a unique polynomial $s_1(k)$ satisfying the recurrence (20), which is exactly the assertion of Claim 2.

The proof that $\prod_{i=1}^n (x + 2y + 3i + 1)_{n-i}$ divides $D_n(x, y)$ is now complete.

Step 4. $D_n(x, y)$ is a polynomial in x of maximal degree $n(3n + 1)/2$, and the same is true for the maximal degree of $D_n(x, y)$ in y . This is because each term in the defining expansion of the determinant $D_n(x, y)$ has degree $n(3n + 1)/2$ in x , and the same in y . Since the righthand side of (4), which by Steps 1–3 divides $D_n(x, y)$ as a polynomial in x and y , also has degree $n(3n + 1)/2$ in x , respectively y , $D_n(x, y)$ and the righthand side of (4) differ only by a multiplicative constant.

Step 5. The evaluation of the multiplicative constant. By the preceding steps we know that (5) holds. In particular, if we set $y = 0$, we have

$$(30) \quad \det_{1 \leq i, j \leq n} ((x+2j-i+1)_{n+i} (i-j+1)_{n-i}) = C(n) \prod_{j=1}^n (x+j+1)_{n+j}.$$

(The reader should be aware that the second term in the determinant $D_n(x, y)$, as given by (4), vanishes for $y = 0$ because of the presence of the factor $(y - i - j + 1)_{n+i}$.) The determinant on the lefthand side of (30) is a lower triangular matrix, hence it equals the product of its diagonal entries, which is $\prod_{j=1}^n (x+j+1)_{n+j} (n-j)!$. Therefore $C(n)$ is equal to $\prod_{j=1}^n (n-j)! = \prod_{j=1}^n (j-1)!$.

This finishes the proof of (4) and thus of the Theorem. \square

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DEPARTMENT OF MATHEMATICS, GEORGIA INSTITUTE OF TECHNOLOGY, ATLANTA, GA 30332-0160

Email address: ciucu@math.gatech.edu

INSTITUT FÜR MATHEMATIK DER UNIVERSITÄT WIEN, STRUDLHOFGASSE 4, A-1090 WIEN, AUSTRIA

Email address: kratt@ap.univie.ac.at