IDEAL DECOMPOSITIONS AND COMPUTATION OF TENSOR NORMAL FORMS

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ABSTRACT. Symmetry properties of r-times covariant tensors T can be described by certain linear subspaces W of the group ring $\mathbb{K}[\mathcal{S}_r]$ of a symmetric group \mathcal{S}_r . If for a class of tensors T such a W is known, the elements of the orthogonal subspace W^{\perp} of W within the dual space $\mathbb{K}[\mathcal{S}_r]^*$ of $\mathbb{K}[\mathcal{S}_r]$ yield linear identities needed for a treatment of the term combination problem for the coordinates of the T. We give the structure of these W for every situation which appears in symbolic tensor calculations by computer. Characterizing idempotents of such W can be determined by means of an ideal decomposition algorithm which works in every semisimple ring up to an isomorphism. Furthermore, we use tools such as the Littlewood-Richardson rule, plethysms and discrete Fourier transforms for \mathcal{S}_r to increase the efficience of calculations. All described methods were implemented in a Mathematica package called PERMS.

1. The Term Combination Problem for Tensors

The use of computer algebra systems for symbolic calculations with tensor expressions is very important in differential geometry, tensor analysis and general relativity theory. The investigations of this paper¹ are motivated by the following *term combination problem* or *normal form problem* which occurs within such calculations.

Let us consider real or complex linear combinations

(1.1)
$$\tau = \sum_{i=1}^{n} \alpha_i T_{(i)} , \quad \alpha_i \in \mathbb{R}, \mathbb{C}$$

of expressions $T_{(i)}$ which are formed from the coordinates of certain tensors A, B, C, \ldots by multiplication and, possibly, contractions of some pairs of indices. An example of such an expression is

(1.2)
$$A_{iabc} A^a_{ikd} B^{bd}_{e} C^{ec} \quad .$$

In (1.2) we use Einstein's summation convention. Further we assume that each of the numbers of A, B, C, \ldots is constant if we run through the set of the $T_{(i)}$. Now we aim to carry out symbolic calculations with expressions of the type (1.1), (1.2)

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¹This paper is a summary of our Habilitationsschrift [10], where all proofs can be found. A part of the proofs were published in earlier papers [8, 9, 12], too. An abridged version of this summary is the paper [13].

according to the rules of the Ricci calculus. We assume that there is a metric tensor g which allows us to raise or lower indices, for instance

$$T_i^{\ jkb} = g^{bc} T_i^{\ jk} , \quad T_{ia}^{\ j} = g_{ac} T_i^{\ cj} ,$$

If now the tensors A, B, C, \ldots have symmetries and/or fulfil linear identities², then there exist relations between the $T_{(i)}$ in (1.1). (We restrict us to linear relations.) Thus the problem arises to detect such relations in sums (1.1), generated by symbolic calculations, and to reduce (1.1) to linear combinations of linearly independent $T_{(i)}$ (normal forms).

It is well-known that the representation theory of symmetric groups S_r yields powerful tools to treat this *term combination problem*. The connection between tensors and the representation theory of S_r has been considered already in books of J.A. Schouten [26] (1924), H. Weyl [27] (1939) and H. Boerner [2] (1955). In the 1940s Littlewood has developed and used tools such as the Littlewood-Richardson rule and plethysms for the investigation of tensors (see S.A. Fulling et al. [14] (references) and D.E. Littlewood [20] (appendix)).

Applying the same methods, Fulling, King, Wybourne and Cummins [14] have calculated large lists of normal form terms of polynomials of the Riemann curvature tensor and its covariant derivatives (by means of the program package Schur [28]). In their paper [14] they formulated the following steps to solve the above term combination problem for tensors:

- (a) Generate the space spanned by the set of homogeneous monomials of a definite 'order' or 'degree of homogeneity' formed from the coordinates of tensors of relevance by multiplication and index-pair contraction.
- (b) Construct a basis of this space (*normal forms*).
- (c) Present an algorithm for expressing an arbitrary element of the space in terms of the basis.

Our present paper yields a method to solve (b) and (c) for arbitrary tensors.

2. Tensors and the Group Ring of a Symmetric Group

We make use of the following connection between r-times covariant tensors $T \in \mathcal{T}_r V$ over a finite-dimensional K-vector space V and elements of the group ring $\mathbb{K}[\mathcal{S}_r]$ of a symmetric group \mathcal{S}_r over a field $\mathbb{K} = \mathbb{R}$ or $\mathbb{K} = \mathbb{C}$.

Definition 2.1. Any tensor $T \in \mathcal{T}_r V$ and any *r*-tuple $b := (v_1, \ldots, v_r) \in V^r$ of *r* vectors from *V* induce a function $T_b : \mathcal{S}_r \to \mathbb{K}$ according to the rule

(2.1)
$$T_b(p) := T(v_{p(1)}, \dots, v_{p(r)}) , p \in S_r.$$

We identify this function with the group ring element $T_b := \sum_{p \in S_r} T_b(p) \, p \in \mathbb{K}[S_r].$

²For instance, the Riemann tensor has the symmetry $R_{ijkl} = -R_{jikl} = -R_{ijlk} = R_{klij}$ and fulfils $R_{ijkl} + R_{iklj} + R_{iljk} = 0$. See P. Günther [15] for a large collection of identities for the curvature tensor, Weyl tensor, etc. and covariant derivatives of these tensors.

We allow the linear dependence of the v_i and repetitions of vectors in the above *r*-tuple *b*. Obviously, two tensors $S, T \in \mathcal{T}_r V$ fulfil S = T iff $S_b = T_b$ for all $b \in V^r$.

We try to describe symmetry properties of tensors with the help of the T_b by the following *principle*: If a "class" of tensors with a certain symmetry property is given, then we search such a linear subspace $W \subseteq \mathbb{K}[\mathcal{S}_r]$ that contains all T_b of the tensors from the "class" being considered. The linear identities which characterize W yield then identities for the coordinates of the T which can be used in the treatment of the term combination problem.

Every $a = \sum_{p \in S_r} a(p) p \in \mathbb{K}[S_r]$ acts as so-called symmetry operator $a: T \mapsto aT$ on tensors $T \in \mathcal{T}_r V$ if we define

(2.2)
$$(aT)(v_1, \ldots, v_r) := \sum_{p \in S_r} a(p) T(v_{p(1)}, \ldots, v_{p(r)}) , \quad v_i \in V.$$

We denote by '*' the mapping $*: a = \sum_{p \in S_r} a(p) p \mapsto a^* := \sum_{p \in S_r} a(p) p^{-1}$. Furthermore, if $p \in S_r$ and $b = (v_1, \ldots, v_r) \in V^r$, then we denote by pb the *r*-tuple $pb := (v_{p(1)}, \ldots, v_{p(r)})$. Many of our calculations are based on

Lemma 2.2. ³ If $a = \sum_{q \in S_r} a(q) q \in \mathbb{K}[S_r], T \in \mathcal{T}_r V, p, q \in S_r and b =$ $(v_1,\ldots,v_r) \in V^r$, then we have

 $T_b(p \circ q) = (qT)_b(p) = T_{pb}(q)$ $q(pb) = (p \circ q)b$ (2.3)

$$(2.4) q(pb) = (p \circ q)b$$

$$(2.5) (aT)_b = T_b \cdot a^*$$

$$(2.6) T_b = p \cdot T_{pb}.$$

The following symmetry concepts are used for tensors. (See B. Fiedler [10, Sec.III.2] and B. Fiedler [12]. See also R. Merris [22, pp.151,153,157], H. Boerner [2, p.127], G. Eisenreich [7, p.601].)

- (a) Let $\mathfrak{r} \subseteq \mathbb{K}[\mathcal{S}_r]$ be a right ideal of $\mathbb{K}[\mathcal{S}_r]$ for which an $a \in \mathfrak{r}$ Definition 2.3. and a $T \in \mathcal{T}_r V$ exist such that $aT \neq 0$. Then the tensor set $\mathcal{T}_r := \{aT \mid a \in \mathcal{T}\}$ $\mathfrak{r}, T \in \mathcal{T}_r V$ is called the *symmetry class* of tensors defined by \mathfrak{r} .
 - (b) Let $a_1, \ldots, a_n \in \mathbb{K}[S_r]$ be a finite set of group ring elements. We say that a tensor $T \in \mathcal{T}_r V$ possesses a symmetry defined by a_1, \ldots, a_n if T satisfies the linear equation system $a_i T = 0, (i = 1, ..., n).$

If e is a generating idempotent of a right ideal $\mathfrak{r} = e \cdot \mathbb{K}[S_r]$ that defines a symmetry class $\mathcal{T}_{\mathfrak{r}}$ then $\mathcal{T}_{\mathfrak{r}}$ fulfils $\mathcal{T}_{\mathfrak{r}} = \{eT \mid T \in \mathcal{T}_r V\}$ and a tensor $T \in \mathcal{T}_r V$ belongs to $\mathcal{T}_{\mathfrak{r}}$ iff eT = T. (See H. Boerner [2, p.127] or B. Fiedler [10, Sec.III.2.1].)

Now it can be shown that all T_b of tensors T which have a symmetry (a) or (b) lie in a certain left ideal of $\mathbb{K}[\mathcal{S}_r]$.

³See B. Fiedler [10, Sec.III.1] and B. Fiedler [12].

Proposition 2.4. ⁴ Let $e \in \mathbb{K}[S_r]$ be an idempotent. Then a $T \in \mathcal{T}_r V$ fulfils the condition eT = T iff $T_b \in \mathfrak{l} := \mathbb{K}[S_r] \cdot e^*$ for all $b \in V^r$, i.e. all T_b of T lie in the left ideal \mathfrak{l} generated by e^* .

Proposition 2.5. ⁵ Let $a_1, \ldots, a_m \in \mathbb{K}[S_r]$ be given group ring elements. A $T \in \mathcal{T}_r V$ satisfies a system of linear identities $a_i T = 0$, $(i = 1, \ldots, m)$, iff $T_b \in \mathfrak{l} := \{x \in \mathbb{K}[S_r] \mid x \cdot a_i^* = 0, i = 1, \ldots, m\}$ for all $b \in V^r$, i.e. all T_b of T lie in the left annihilator ideal \mathfrak{l} of the set $\{a_1^*, \ldots, a_m^*\}$.

The proofs follow easily from (2.5). A further result is

Proposition 2.6. ⁶ If dim $V \ge r$, then every left ideal $\mathfrak{l} \subseteq \mathbb{K}[S_r]$ fulfils $\mathfrak{l} = \mathcal{L}_{\mathbb{K}}\{T_b \mid T \in \mathcal{T}_{\mathbb{I}^*}, b \in V^r\}$. (Here $\mathcal{L}_{\mathbb{K}}$ denotes the forming of the linear closure.)

If dim V < r, then the T_b of the tensors from $\mathcal{T}_{\mathfrak{l}^*}$ will span only a linear subspace of \mathfrak{l} in general.

In the case of tensors T with index contractions the role of the T_b is played by certain sums $\sum_{b \in \mathfrak{B}_{b_0}} \gamma_b T_b$, which we now define.

Definition 2.7. Let $g \in \mathcal{T}_2 V$ be a fundamental tensor with arbitrary signature on V and $\mathcal{B} = \{n_1, \ldots, n_d\}$ be an orthonormal basis of V with respect to g. Further let r, l be integers with $2 \leq 2l < r$ and $b_0 = (v_{2l+1}, \ldots, v_r) \in \mathcal{B}^{r-2l}$ be a fixed (r-2l)-tuple of vectors from \mathcal{B} . Then we denote by \mathfrak{B}_{b_0} the set ⁷ of r-tuples of basis vectors

$$\mathfrak{B}_{b_0} := \left\{ (w_1, w_1, w_2, w_2, \dots, w_l, w_l, v_{2l+1}, \dots, v_r) \in \mathcal{B}^r \mid (w_1, \dots, w_l) \in \mathcal{B}^l \right\}.$$

Moreover, we set $\gamma_b := \prod_{i=1}^l g(w_i, w_i) \in \{1, -1\}$ for every $b \in \mathfrak{B}_{b_0}$.

Proposition 2.8.⁸ Let $T \in \mathcal{T}_r V$ be a tensor and $g \in \mathcal{T}_2 V$ be a fundamental tensor. We determine all tensor coordinates with respect to an orthonormal basis $\mathcal{B} = \{n_1, \ldots, n_d\}$ of V. Let $b_0 = (n_{i_{2l+1}}, \ldots, n_{i_r}) \in \mathcal{B}^{r-2l}$ be a fixed (r-2l)-tuple of basis vectors. Then

(2.7)
$$\sum_{p \in S_r} (pT)_{j_1 \ j_2 \ \dots \ j_l \ i_{2l+1} \dots i_r}^{j_1 \ j_2 \ \dots \ j_l} p = \sum_{b \in \mathfrak{B}_{b_0}} \gamma_b T_b .$$

Due to Prop. 2.8 the $\sum_{b \in \mathfrak{B}_{b_0}} \gamma_b T_b$ are objects which contain information about the tensor coordinates of T with l index-pair contractions. In the case of tensors Twith index contractions we search for linear subspaces $W \subseteq \mathbb{K}[\mathcal{S}_r]$ which contain all $\sum_{b \in \mathfrak{B}_{b_0}} \gamma_b T_b$ for a fixed b_0 .

⁴See B. Fiedler [12] or B. Fiedler [10, Prop. III.2.5, III.3.1, III.3.4].

⁷We set $b_0 := \emptyset$ and $\mathfrak{B}_{\emptyset} := \{(w_1, w_1, w_2, w_2, \dots, w_l, w_l) \in \mathcal{B}^r \mid (w_1, \dots, w_l) \in \mathcal{B}^l\}$ in the case r = 2l > 0.

⁸See B. Fiedler [12] or B. Fiedler [10, Prop. III.3.31].

⁵See B. Fiedler [12] or B. Fiedler [10, Prop. III.3.3, III.3.4].

⁶See B. Fiedler [12] or B. Fiedler [10, Prop. III.2.6].

The left ideals l from the Propositions 2.4, 2.5 and 2.6 are the simplest examples of linear subspaces W describing tensor symmetries. Before we give such subspaces W in more complicated cases, we will explain how they can be used in the treatment of the term combination problem for tensors.⁹

3. The Treatment of the Term Combination Problem

The term combination problem from Sec. 1 can be reformulated as the problem to find all linear identities between the summands of given tensor expressions

(3.1)
$$\tau_{i_1\dots i_r} := \sum_{p\in\mathcal{P}} c_p T_{i_{p(1)}\dots i_{p(r)}} , \quad c_p\in\mathbb{K}, \,\mathcal{P}\subseteq\mathcal{S}_r \text{ or}$$

(3.2)
$$\tau_{i_{2l+1}\dots i_r} := \sum_{p \in \mathcal{P}} c_p (pT)_{j_1 \ j_2 \ \dots \ j_l \ i_{2l+1}\dots i_r}^{j_1 \ j_2 \ \dots \ j_l} , \quad c_p \in \mathbb{K}, \ \mathcal{P} \subseteq \mathcal{S}_r$$

where $T \in \mathcal{T}_{\mathfrak{l}^*}$ is a tensor from a given symmetry class defined by a left ideal \mathfrak{l} (or a right ideal $\mathfrak{r} = \mathfrak{l}^*$). We assume that (3.1) and (3.2) are results of symbolic computer calculations. \mathcal{P} is a subset of permutations, which is determined by the concrete form of the given expressions (3.1) or (3.2).

Let $W \subseteq \mathfrak{l}$ be a linear subspace which contains all T_b or $\sum_{b \in \mathfrak{B}_{b_0}} \gamma_b T_b$ of T. If we consider the *orthogonal subspace* $W^{\perp} := \{x \in \mathbb{K}[\mathcal{S}_r]^* \mid \forall w \in W : \langle x, w \rangle = 0\}$ of W, then every $x \in W^{\perp}$ yields a linear identity for the coordinates of T since

$$(3.3) \quad 0 = \langle x, T_b \rangle = \sum_{p \in \mathcal{S}_r} x_p T_b(p) = \sum_{p \in \mathcal{S}_r} x_p T_{i_{p(1)} \dots i_{p(r)}} \quad \text{or}$$

$$0 = \langle x, \sum_{b \in \mathfrak{B}_{b_0}} \gamma_b T_b \rangle = \sum_{\substack{b \in \mathfrak{B}_{b_0} \\ p \in \mathcal{S}_r}} \gamma_b T_b(p) x_p = \sum_{p \in \mathcal{S}_r} x_p (pT)_{j_1 \dots j_l}^{j_1 \dots j_l} \sum_{\substack{i_{2l+1} \dots i_r}} y_{i_{2l+1} \dots i_r}$$

where $x_p := \langle x, p \rangle$, $p \in S_r$. (The last steps are correct if all *b* occurring in (3.3) are *r*-tuples of basis vectors of *V*.) Every identity (3.3) can be used to eliminate certain summands in (3.1), (3.2). If *W* is spanned by all T_b or $\sum_{b \in \mathfrak{B}_{b_0}} \gamma_b T_b$ of the tensors considered, then W^{\perp} contains all linear identities which are possible between summands of expressions (3.1), (3.2) (compare Prop. 2.6).

If a basis $\{h_1, \ldots, h_k\}$ of W is known, then the coefficients x_p of the $x \in W^{\perp}$ can be obtained from the linear equation system

(3.4)
$$\langle x, h_i \rangle = \sum_{p \in \mathcal{S}_r} h_i(p) x_p = 0 \qquad (i = 1, \dots, k).$$

Thus an important goal is to find such a basis $\{h_1, \ldots, h_k\}$ of W. (An efficient algorithm for that purpose is given in Prop. 5.1.)

Note that (3.4) is a very large system with a $(k \times r!)$ -coefficient matrix, $k = \dim W$. However, since we only need identities to reduce sums (3.1), (3.2), we can restrict us to solutions of (3.4) which fulfil $x_p = 0$ for $p \in S_r \setminus \mathcal{P}$. This reduces the number of columns to $|\mathcal{P}|$. Furthermore, every of our spaces W is a linear

⁹See B. Fiedler [12] or B. Fiedler [10, Sec. III.1, III.4.1].

subspace of a left ideal $\mathfrak{l} = \mathbb{K}[\mathcal{S}_r] \cdot e$. A decomposition $e = e_1 + \ldots + e_m$ of the generating idempotent into pairwise orthogonal, primitive idempotents e_i induces a decomposition $W = W_1 \oplus \ldots \oplus W_m$ with $W_i \subseteq \mathbb{K}[\mathcal{S}_r] \cdot e_i$ and a decomposition of the tensors $T \in \mathcal{T}_{\mathfrak{l}^*}$: $T = e_1^*T + \ldots + e_m^*T$. Then we can transform (3.1), (3.2) into expressions formed from the e_i^*T , for instance

(3.5) (3.2)
$$\Rightarrow \tau_{i_{2l+1}\dots i_r} := \sum_{i=1}^m \sum_{p \in \mathcal{P}} c_p \left(p(e_i^*T) \right)_{j_1 \ j_2 \ \dots \ j_l \ i_{2l+1}\dots i_r}^{j_1 \ j_2 \ \dots \ j_l} ,$$

and use the smaller equation systems of the smaller spaces W_i to determine linear identities for the coordinates of the e_i^*T . Note, however, that a reduction of expressions such as (3.5) by means of identities of the W_i leads to a linear combination of coordinates of the e_i^*T which cannot be reckoned back into a linear combination of the coordinates of T in general.

4. The Algorithm for Ideal Decompositions

Two problems became visible up to now: We need methods

- (a) to determine generating idempotents e for left/right ideals of $\mathbb{K}[\mathcal{S}_r]$ for which such idempotents are unknown (such as in Prop. 2.5).
- (b) to decompose a given idempotent $e \in \mathbb{K}[S_r]$ into a sum $e = e_1 + \ldots + e_m$ of pairwise orthogonal primitive idempotents e_i .

We developed an algorithm which solves these problems. First versions of this algorithm were designed for $\mathbb{K}[S_r]$ (see B. Fiedler [8, 9]). But it turned out that this algorithm works even in an arbitrary semisimple ring \mathfrak{R} which fulfils:

(A) We know explicitly a decomposition

(4.1)
$$\mathfrak{R} = \bigoplus_{i=1}^{m} \mathfrak{R} \cdot y_i \quad \text{or} \quad \mathfrak{R} = \bigoplus_{i=1}^{m} y_i \cdot \mathfrak{R}$$

of the full ring \mathfrak{R} into minimal left or right ideals generated by known primitive idempotents y_i . Pairwise orthogonality of the y_i is not required.

(B) We know in \mathfrak{R} a method to construct explicitly a solution $x \in \mathfrak{R}$ for every equation

(4.2)
$$e \cdot a \cdot x \cdot e = e$$
 or $e \cdot x \cdot a \cdot e = e$,

where $e \in \mathfrak{R}$ is a primitive idempotent and $a \in \mathfrak{R}$ is a ring element with $e \cdot a \neq 0$ or $a \cdot e \neq 0$, respectively.

We describe now a version (L) of the algorithm for left ideals of a semisimple ring \mathfrak{R} . Obviously an analogous version (R) for right ideals can be formulated, too. (See B. Fiedler [10, Sec.I.2] and B. Fiedler [8, 9].)

A frequent step of the algorithm is the construction of a generating idempotent for a left ideal $\mathfrak{l} = \mathfrak{R} \cdot e \cdot a$, where $e \in \mathfrak{R}$ is a primitive idempotent and $a \in \mathfrak{R}$ is a ring element with $e \cdot a \neq 0$. This is possible by **Proposition 4.1.**¹⁰ For the above $e, a \in \mathfrak{R}$ there exists such an $x \in \mathfrak{R}$ that

 $(4.3) e \cdot a \cdot x \cdot e = e.$

Moreover, the ring element $e' := x \cdot e \cdot a$ formed with this x is an idempotent which generates the minimal left ideal $\Re \cdot e \cdot a$.

A second construction orthogonalizes idempotents. Let $\mathfrak{l} = \mathfrak{R} \cdot e$ and $\tilde{\mathfrak{l}} = \mathfrak{R} \cdot \tilde{e}$ be two left ideals with known generating idempotents e and \tilde{e} . Assume that e is primitive and $e \cdot \tilde{e} \neq e$, i.e. $\mathfrak{l} \not\subseteq \tilde{\mathfrak{l}}$. Then the sum $\mathfrak{l} + \tilde{\mathfrak{l}} = \mathfrak{l} \oplus \tilde{\mathfrak{l}}$ is direct since the minimality of \mathfrak{l} yields $\mathfrak{l} \cap \tilde{\mathfrak{l}} = \{0\}$. Now we search for new generating idempotents f and \tilde{f} of \mathfrak{l} and $\tilde{\mathfrak{l}}$ which fulfil $f \cdot \tilde{f} = \tilde{f} \cdot f = 0$.

Theorem 4.2.¹¹ The above orthogonalization problem can be solved in 2 steps:

(i) We can find such a ring element $x \in \mathfrak{R}$ that

$$(4.4) e \cdot (1 - \tilde{e}) \cdot x \cdot e = e .$$

If we use this x to form $f := (1 - \tilde{e}) \cdot x \cdot e$, then f is a generating idempotent of \mathfrak{l} which satisfies $\tilde{e} \cdot f = 0$.

(ii) For a result f of step (i) there exists an $\tilde{x} \in \mathfrak{R}$ such that

(4.5)
$$f \cdot (1 - \tilde{e}) \cdot \tilde{x} \cdot f = f.$$

If we make use of \tilde{x} to form $\tilde{f} := \tilde{e} - (1 - \tilde{e}) \cdot \tilde{x} \cdot f \cdot \tilde{e}$, then \tilde{f} is a generating idempotent of $\tilde{\mathfrak{l}}$ which fulfils $f \cdot \tilde{f} = \tilde{f} \cdot f = 0$.

Now we can describe our ALGORITHM (L), which allows us to decompose every left ideal $\mathfrak{l} = \mathfrak{R} \cdot a$ with known generating element $a \neq 0$ into a direct sum $\mathfrak{l} = \bigoplus_{i=1}^{m} \mathfrak{l}_i$ of minimal left ideals \mathfrak{l}_i explicitly (see B. Fiedler [10, Sec.I.2] and B. Fiedler [8, 9]).

A multiplication of (4.1) by a yields a sum

(4.6)
$$\mathfrak{l} = \mathfrak{R} \cdot a = \sum_{\substack{i=1\\y_i \cdot a \neq 0}}^{m} \mathfrak{R} \cdot y_i \cdot a$$

of minimal left ideals for l which however is not direct in general. Now we can carry out the following steps:

- (1) The first summand in (4.6) is a minimal left ideal. We denote it by l_1 and we determine a generating idempotent e_1 of l_1 by means of Prop. 4.1.
- (2) We search for the first minimal left ideal $\Re \cdot y_i \cdot a$ in (4.6) which is not contained in \mathfrak{l}_1 that means for which

$$(4.7) y_i \cdot a \cdot e_1 \neq y_i \cdot a.$$

We denote it by \mathfrak{l}_2 and we construct a generating idempotent e_2 of \mathfrak{l}_2 by means of Prop. 4.1. Since \mathfrak{l}_2 is minimal and $\mathfrak{l}_2 \not\subseteq \mathfrak{l}_1$, we obtain $\mathfrak{l}_1 \cap \mathfrak{l}_2 = \{0\}$.

 $^{^{10}\}mathrm{See}$ B. Fiedler [10, Prop. I.2.1]. Compare B. Fiedler [8, 9].

¹¹B. Fiedler [10, Thm. I.2.4] and B. Fiedler [8, 9]. The proof uses the fact that the set $\{e - x \cdot e + e \cdot x \cdot e \mid x \in \Re\}$ is the set of all generating idempotents of $\mathfrak{l} = \Re \cdot e$ (compare D.S. Passman [23, p.137]).

Thus \mathfrak{l}_1 and \mathfrak{l}_2 form a direct sum $\tilde{\mathfrak{l}}_2 := \mathfrak{l}_1 \oplus \mathfrak{l}_2$. Because e_2 is primitive and $e_2 \cdot e_1 \neq e_2$, we can determine new, orthogonal, generating idempotents \hat{f}_1, f_2 of $\mathfrak{l}_1, \mathfrak{l}_2$ by means of Theorem 4.2. Then $\tilde{f}_2 := \hat{f}_1 + f_2$ is a generating idempotent of $\tilde{\mathfrak{l}}_2$.

(3) Now we search for the next minimal left ideal $\Re \cdot y_i \cdot a$ in (4.6) which is not contained in \tilde{l}_2 that means for which

$$(4.8) y_i \cdot a \cdot f_2 \neq y_i \cdot a \,.$$

We denote it by \mathfrak{l}_3 . We construct a primitive generating idempotent e_3 of \mathfrak{l}_3 and pass over to new orthogonal idempotents \hat{f}_2, f_3 instead of \tilde{f}_2, e_3 . This leads us to the left ideal $\tilde{\mathfrak{l}}_3 := \tilde{\mathfrak{l}}_2 \oplus \mathfrak{l}_3$ which has the generating idempotent $\tilde{f}_3 := \hat{f}_2 + f_3$.

(4) We continue this procedure until we have processed all left ideals in (4.6). The result is a left ideal $\tilde{\mathfrak{l}}_n := \mathfrak{l}_1 \oplus \ldots \oplus \mathfrak{l}_n$ and a generating idempotent \tilde{f}_n of $\tilde{\mathfrak{l}}_n$.

Obviously, $\tilde{\mathfrak{l}}_n \subseteq \mathfrak{l}$ since every left ideal \mathfrak{l}_i is a summand in (4.6). Furthermore, every summand $\mathfrak{R} \cdot y_i \cdot a$ from (4.6) which had been considered before we had reached the ideal $\tilde{\mathfrak{l}}_n$ is contained in $\tilde{\mathfrak{l}}_{n-1} = \mathfrak{l}_1 \oplus \ldots \oplus \mathfrak{l}_{n-1} \subseteq \tilde{\mathfrak{l}}_n$. All other summands $\mathfrak{R} \cdot y_i \cdot a$ of (4.6) lie in $\tilde{\mathfrak{l}}_n$. This leads to $\mathfrak{l} \subseteq \tilde{\mathfrak{l}}_n$ and $\mathfrak{l} = \tilde{\mathfrak{l}}_n$.

According to Theorem 4.2 (ii), every generating idempotent f_k of \mathfrak{l}_k can be written as $\hat{f}_k = (1 - z_k) \cdot \tilde{f}_k$ with a $z_k \in \mathfrak{R}$ which we have already determined to carry out the orthogonalization $(e_{k+1}, \tilde{f}_k) \mapsto (f_{k+1}, \hat{f}_k)$. Thus we can write

$$\tilde{f}_n = \hat{f}_{n-1} + f_n = (1 - z_{n-1}) \cdot \tilde{f}_{n-1} + f_n = (1 - z_{n-1}) \cdot (\hat{f}_{n-2} + f_{n-1}) + f_n$$

$$= (1 - z_{n-1}) \cdot (1 - z_{n-2}) \cdot \tilde{f}_{n-2} + (1 - z_{n-1}) \cdot f_{n-1} + f_n$$

$$\vdots$$

$$(4.9) = \sum_{k=1}^{n-1} (1 - z_{n-1}) \cdot (1 - z_{n-2}) \cdot \dots \cdot (1 - z_k) \cdot f_k + f_n.$$

Formula (4.9) presents a decomposition $\tilde{f}_n = \sum_{k=1}^n h_k$ of \tilde{f}_n into summands which fulfil $h_k := (1 - z_{n-1}) \cdot (1 - z_{n-2}) \cdot \ldots \cdot (1 - z_k) \cdot f_k \in \mathfrak{l}_k$ and $h_n := f_n \in \mathfrak{l}_n$. Thus, $\tilde{f}_n = \sum_{k=1}^n h_k$ is the decomposition of \tilde{f}_n which corresponds to $\mathfrak{l} = \bigoplus_{k=1}^n \mathfrak{l}_k$ and the h_k are pairwise orthogonal generating idempotents of the \mathfrak{l}_k .

Obviously, the algorithm (L) solves the above problem (b). Furthermore, the algorithm (L) can be extended to left ideals which are non-direct sums $\mathbf{l} = \sum_{i=1}^{h} \mathfrak{R} \cdot a_i$ by applying its steps to the summands of $\mathbf{l} = \sum_{i=1}^{h} \sum_{j=1}^{m} \mathfrak{R} \cdot y_j \cdot a_i$. Likewise, we can construct generating idempotents and decompositions for right ideals $\mathbf{r} = \sum_{i=1}^{h} a_i \cdot \mathfrak{R}$ by the algorithm version (R). If our left/right ideals are intersections $\mathbf{l} = \bigcap_{i=1}^{h} \mathfrak{R} \cdot e_i$ or $\mathbf{r} = \bigcap_{i=1}^{h} e_i \cdot \mathfrak{R}$ of left/right ideals $(e_i \text{ idempotents})$, then their right/left annihilator ideals are $\mathcal{A}_r(\mathbf{l}) = \sum_{i=1}^{h} (1-e_i) \cdot \mathfrak{R}$ or $\mathcal{A}_l(\mathbf{r}) = \sum_{i=1}^{h} \mathfrak{R} \cdot (1-e_i)$, respectively. In this case we can construct a generating idempotent e of $\mathcal{A}_r(\mathbf{l})$ or $\mathcal{A}_l(\mathbf{r})$ by (R) or (L), respectively, and form e' := 1 - e to obtain a generating

idempotent e' of \mathfrak{l} or \mathfrak{r} . Thus our algorithms solve problem (a) for non-direct sums or intersections of left/right ideals. See B.Fiedler [10, Sec.I.4] or B. Fiedler [9] for further details.

5. Completions of the Decomposition Algorithms

The basic assumptions (A) and (B). Actual decomposition constructions can only be carried out by the algorithms (L) or (R) if our semisimple ring \Re fulfils the above assumptions (A) and (B). This is the case for

- (1) the group ring $\mathfrak{R} = \mathbb{K}[\mathcal{S}_r]$ of a symmetric group,
- (2) a ring $\mathfrak{R} = \bigotimes_{i=1}^{m} \mathbb{S}_{i}^{n_{i} \times n_{i}}$, that is an outer direct product of full $(n_{i} \times n_{i})$ -matrix rings over skew fields \mathbb{S}_{i} ,
- (3) all semisimple rings \mathfrak{R} for which an isomorphism $D: \mathfrak{R} \to \mathfrak{R}' = \bigotimes_{i=1}^m \mathbb{S}_i^{n_i \times n_i}$ onto a ring \mathfrak{R}' of the second type is explicitly known.

(See B. Fiedler [10, Sec.I.3].) According to Wedderburn's Theorem every semisimple ring is isomorphic to a ring of type 2. Thus statement 2 means that the decomposition algorithms (L) and (R) work in every semisimple ring up to an isomorphism.

Let us consider $\mathfrak{R} = \bigotimes_{i=1}^{m} \mathbb{S}_{i}^{n_{i} \times n_{i}}$. We denote by $C_{kl} \in \mathbb{S}_{i}^{n_{i} \times n_{i}}$ a matrix in which exactly the element located in the k-th row and the j-th column is equal to $1 \in \mathbb{S}_{i}$ whereas all other elements vanish. Then a decomposition of \mathfrak{R} into minimal left/right ideals is given by the decompositions $\mathbb{S}_{i}^{n_{i} \times n_{i}} = \bigoplus_{j=1}^{n_{i}} \mathbb{S}_{i}^{n_{i} \times n_{i}} \cdot C_{jj}$ and $\mathbb{S}_{i}^{n_{i} \times n_{i}} = \bigoplus_{j=1}^{n_{i}} C_{jj} \cdot \mathbb{S}_{i}^{n_{i} \times n_{i}}$ of the matrix rings into minimal left/right ideals. Furthermore there exists a very fast procedure to solve (4.2) in \mathfrak{R} .

Since a primitive idempotent $e \in \mathfrak{R}$ has only one non-vanishing block matrix (i.e. $e = (0, \ldots, E, \ldots, 0)$, where $E \in \mathbb{S}_i^{n_i \times n_i}$ is also a primitive idempotent), an equation such as $e \cdot a \cdot x \cdot e = e$ leads to a single matrix equation $E \cdot A \cdot X \cdot E = E$. Moreover, E can be written as $E = f^t \cdot h$ with row vectors $f, h \in \mathbb{S}_i^{n_i}$, where h is a non-vanishing row of E. If we set $m = h \cdot A$ and determine non-vanishing elements m_{j_0}, f_{k_0} of m, f, then $X := (m_{j_0})^{-1} (f_{k_0})^{-1} C_{j_0 k_0}$ is a solution of $E \cdot A \cdot X \cdot E = E$, which yields a solution $x = (0, \ldots, X, \ldots, 0)$ of $e \cdot a \cdot x \cdot e = e$ (B. Fiedler [10, Sec. I.1.2, I.3.2]). Obviously, this procedure will run very fast on a computer.

For $\mathbb{K}[\mathcal{S}_r]$ the well-known decomposition of $\mathbb{K}[\mathcal{S}_r]$ into minimal left/right ideals by means of *Young symmetrizers* guarantees (A). See B. Fiedler [8, 9] for (B).

If for a semisimple ring \mathfrak{R} an above isomorphism $D : \mathfrak{R} \to \mathfrak{R}'$ is known (and practicable on a computer), then every ideal decomposition problem for \mathfrak{R} can be transferred to \mathfrak{R}' and treated there by the algorithms (L) and (R).

If $\mathfrak{R} = \mathbb{C}[G]$ is the group ring of a finite group G, then we have $\mathbb{S}_i = \mathbb{C}$ for all i and the isomorphism D is called a *discrete Fourier transform* for G. Explicit algorithms for such Fourier transforms are known at least for *abelian groups*, *solvable groups*, *supersolvable groups* and *symmetric groups* (see M. Clausen und U. Baum [5]).

Discrete Fourier transforms. In group rings $\mathfrak{R} = \mathbb{C}[G]$ of large finite groups G (such as $G = \mathcal{S}_r, r \geq 8$), even a single product $a \cdot b$, $a, b \in \mathbb{C}[G]$ can lead to high costs in time and computer memory (see B. Fiedler [10, Sec. I.1.3, I.5.1]). Here

the use of a discrete Fourier transform $D: \mathfrak{R} = \mathbb{C}[G] \to \mathfrak{R}' = \bigotimes_{i=1}^k \mathbb{C}^{n_i \times n_i}$ and the transfer of ideal decomposition problems to \mathfrak{R}' is the most important tool to surmount difficulties. Calculations in \mathfrak{R}' have the following advantages:

- (1) Decompositions (4.1) and solutions of (4.2) can be constructed very fast in \mathfrak{R}' .
- (2) Every product formed during a run of (L) or (R) contains a factor which is a primitive idempotent $e \in \mathfrak{R}'$. Since every such e has only 1 non-vanishing block matrix $E \in \mathbb{C}^{n_i \times n_i}$, i.e. $e = (0, \ldots, 0, E, 0, \ldots, 0)$, the costs for every step of (L) or (R) reduce to the costs of calculations in a ring $\mathbb{C}^{n_i \times n_i}$.
- (3) The algorithms (L) and (R) can be carried out completely within \mathfrak{R}' . Only input and output data have to be mapped between \mathfrak{R} and \mathfrak{R}' by means of D and D^{-1} . Thus, "less fast" Fourier transforms can be useful, too. (See B. Fiedler [10, Sec.I.5.1].)
- (4) In \mathfrak{R}' there is a fast construction of bases of linear subspaces W, which we need to form linear equation systems (3.4).

To describe this construction, we denote by $C_{i,a} \in \mathbb{K}^{n \times n}$ that matrix in which the *i*-th row is equal to a given $a \in \mathbb{K}^n$ whereas all other rows are filled with 0.

Proposition 5.1. Let $\mathfrak{l} = \mathbb{K}^{n \times n} \cdot A$ be a minimal left ideal of $\mathbb{K}^{n \times n}$ with known generating element $0 \neq A \in \mathbb{K}^{n \times n}$ and $B = [b_{ij}]_{n,n} \neq 0$ be a matrix from $\mathbb{K}^{n \times n}$. Determine a row $a \neq 0$ of A and a parametric form Λ of the solution of the linear equation system

(5.1)
$$\sum_{j=1}^{n} b_{ij} \lambda_j = 0 , \quad i = 1, \dots, n , \quad \lambda_i \in \mathbb{K} \text{ (unknowns)}.$$

Then $\mathcal{B} := \{ B \cdot C_{i,a} \mid i \text{ index for which } \lambda_i \text{ is not a parameter in } \Lambda \}$ is a basis of the \mathbb{K} -vector space $B \cdot \mathfrak{l} = B \cdot \mathbb{K}^{n \times n} \cdot A$.

See B. Fiedler [10, Sec.I.1.2] for the proof and other fast basis constructions. Spaces with a structure $B \cdot \mathfrak{l}$ are typical examples of spaces W (see Sec. 6, 7). Further, we see that \mathfrak{l} has the basis $\mathcal{B} = \{C_{i,a} \mid i = 1, \ldots, n\}$ if we use the identity matrix $B = \mathrm{Id} \in \mathbb{K}^{n \times n}$ for B.

For our tensor investigations we need $\Re = \mathbb{K}[S_r]$. M. Clausen und U. Baum [5, 6] developed a very fast Fourier transform for $\mathbb{K}[S_r]$, which bases on Young's seminormal representation of S_r (see also H. Boerner [3] and A. Kerber [17, Vol.I, p.75,76]). However, since the interpreter Mathematica does not allow the full speed and the optimal storage handling of this ingenious algorithm, we use Young's natural representation of S_r as discrete Fourier transform in our Mathematica package PERMS [11]. (See H. Boerner [2, pp.102–108], B. Fiedler [10, Sec.I.5.2].) This implementation works good at least for $S_r, r \leq 8$.

Multiplicities. Obviously, the efficiency of the algorithms (L) and (R) can be improved if we know before a run of (L) or (R) the multiplicities of equivalent minimal left/right ideals l_i or \mathbf{r}_i within decompositions $\mathbf{l} = \bigoplus_{i=1}^m l_i$ or $\mathbf{r} = \bigoplus_{i=1}^m \mathbf{r}_i$

searched for. If the algorithms have constructed such a direct sum of minimal left/right ideals of a fixed equivalence class that the number of summands equals the known multiplicity for this class, then the investigation of the remaining ideals of the class can be cancelled. This reduces the calculation time.

In the case of $\Re = \mathbb{K}[S_r]$ such multiplicies can be calculated by means of the *irreducible characters* of S_r , *Frobenius reciprocity*, the *Littlewood-Richardson rule* and *plethysms*. The determination of the irreducible characters of S_r is possible by the *Murnaghan-Nakayama formula*. We implemented all these tools in our Mathematica package PERMS [11] (see B. Fiedler [10, Sec. II.3–II.6] for descriptions of implementations). For plethysms we use a very efficient method of F. Sänger [24, pp. 29–33]. (See B. Fiedler [10, Sec. II.6.3].)

6. Characterizing Left Ideals of Tensor Products

We continue to list linear subspaces $W \subseteq \mathbb{K}[\mathcal{S}_r]$ describing tensor symmetries. In the case of tensor products two types of products can be considered: $T^{(1)} \otimes \ldots \otimes T^{(m)}$ with possibly different $T^{(i)}$ and $T \otimes \ldots \otimes T$.

Proposition 6.1.¹² Let $\mathfrak{l}_i \subseteq \mathbb{K}[S_{r_i}]$ (i = 1, ..., m) be left ideals and $T^{(i)} \in \mathcal{T}_{\mathfrak{l}_i^*} \subseteq \mathcal{T}_{r_i}V$ be r_i -times covariant tensors from the symmetry classes characterized by the \mathfrak{l}_i . Consider the product

(6.1)
$$T := T^{(1)} \otimes \ldots \otimes T^{(m)} \in \mathcal{T}_r V \quad , \quad r := r_1 + \ldots + r_m \, .$$

For every i we define an embedding

(6.2)
$$\iota_i : \mathcal{S}_{r_i} \to \mathcal{S}_r$$
 , $(\iota_i s)(k) := \begin{cases} \Delta_i + s(k - \Delta_i) & \text{if } r_{i-1} < k \le r_i \\ k & \text{else} \end{cases}$

where $\Delta_i := r_0 + \ldots + r_{i-1}$ and $r_0 := 0$. Then the T_b of the tensor (6.1) fulfil

(6.3)
$$\forall b \in V^r : T_b \in \mathfrak{l} := \mathbb{K}[\mathcal{S}_r] \cdot \mathcal{L}\{\tilde{\mathfrak{l}}_1 \cdot \ldots \cdot \tilde{\mathfrak{l}}_m\} = \mathbb{K}[\mathcal{S}_r] \cdot (\tilde{\mathfrak{l}}_1 \otimes \ldots \otimes \tilde{\mathfrak{l}}_m)$$

where $\tilde{\mathfrak{l}}_i := \iota_i(\mathfrak{l}_i)$ are the embeddings of the \mathfrak{l}_i into $\mathbb{K}[\mathcal{S}_r]$ induced by the ι_i . If $\dim V \geq r$, then the above left ideal \mathfrak{l} is generated by all $T_b \in \mathbb{K}[\mathcal{S}_r]$ which are formed from tensor products (6.1) of arbitrary tensors $T^{(i)} \in \mathcal{T}_{\mathfrak{l}_i^*}$.

Proposition 6.2.¹³ Let $\mathfrak{l}_0 \subseteq \mathbb{K}[\mathcal{S}_m]$ be a left ideal and $T \in \mathcal{T}_{\mathfrak{l}_0^*} \subseteq \mathcal{T}_m V$ be a tensor of order *m* from the symmetry class $\mathcal{T}_{\mathfrak{l}_0^*}$. Consider the product

(6.4)
$$\hat{T} := \underbrace{T \otimes \ldots \otimes T}_{n} \in \mathcal{T}_{mn} V.$$

Then all \hat{T}_b , $b \in V^{mn}$, lie in the left ideal

(6.5)
$$\mathfrak{l} := \mathbb{K}[\mathcal{S}_{mn}] \cdot \mathcal{L}\{\mathfrak{l}_1 \cdot \ldots \cdot \mathfrak{l}_n \cdot \mathfrak{l}'\} = \mathbb{K}[\mathcal{S}_{mn}] \cdot (\mathfrak{l}_1 \otimes \ldots \otimes \mathfrak{l}_n \otimes \mathfrak{l}')$$

 $^{^{12}}$ See B. Fiedler [10, Sec.III.3.2] and B. Fiedler [12].

¹³See B. Fiedler [10, Sec.III.3.2] and B. Fiedler [12].

where $\mathfrak{l}_i := \iota_i(\mathfrak{l}_0)$ are embeddings of \mathfrak{l}_0 into $\mathbb{K}[S_{mn}]$ which are formed by means of mappings (6.2) with $r_1 = \ldots = r_n = m$ and r = mn. Further \mathfrak{l}' denotes the 1-dimensional ideal $\mathfrak{l}' := \mathcal{L}\{\sum_{q \in Q} q\}$ of $\mathbb{K}[Q]$ where $Q \subset S_{mn}$ is the subgroup

(6.6)
$$Q := \left\{ q = \binom{k \cdot m - l}{s(k) \cdot m - l} \underset{\substack{0 \le l \le m - 1}}{\overset{1 \le k \le n}{0 \le l \le m - 1}} \in \mathcal{S}_{mn} \ \middle| \ s \in \mathcal{S}_n \right\} \cong \mathcal{S}_n$$

If dim $V \ge m \cdot n$, then the above left ideal \mathfrak{l} is generated by all $\hat{T}_b \in \mathbb{K}[\mathcal{S}_{mn}]$ which are formed from tensor products (6.4) of arbitrary tensors $T \in \mathcal{T}_{\mathfrak{l}_0^*}$.

Let $\breve{\omega}_G : G \to GL(\mathbb{K}[G])$ denote the regular representation of a finite group G defined by $\breve{\omega}_g(f) := g \cdot f, g \in G, f \in \mathbb{K}[G]$. If we use the above left ideals \mathfrak{l}_i , \mathfrak{l}_0 , \mathfrak{l} to define subrepresentations $\alpha_i := \breve{\omega}_{\mathcal{S}_{r_i}}|_{\mathfrak{l}_i}, \alpha := \breve{\omega}_{\mathcal{S}_m}|_{\mathfrak{l}_0}, \beta := \breve{\omega}_{\mathcal{S}_r}|_{\mathfrak{l}}$, then the representation β is equivalent to a Littlewood-Richardson product or a plethysm¹⁴, respectively (see B. Fiedler [10, Sec.III.3.2]):

- (6.7) $\mathfrak{l} \text{ according to (6.3)} \implies \beta \sim \alpha_1 \# \dots \# \alpha_m \uparrow S_r$
- (6.8) $\mathfrak{l} \text{ according to } (6.5) \implies \beta \sim \alpha \odot [n].$

These results correspond to statements of S.A. Fulling et al. [14]. (6.7) and (6.8) yield valuable information about multiplicities if one wishes to apply the algorithm (L) to \mathfrak{l} .

7. Subspaces Characterizing Tensors with Index Contractions

First we give a universal linear subspace which contains the group ring elements $\sum_{b \in \mathfrak{B}_{ba}} \gamma_b T_b$ of a tensor T with l index contractions for every value of dim V.

Theorem 7.1. Let $V, \mathcal{B}, r, l, g, b_0$ have the meaning given in Def. 2.7 and Prop. 2.8. Consider the partition $\lambda_0 := (2^l, 1^{r-2l}) \vdash r$ and the lexicographically smallest standard tableau t of λ_0 . Form the group ¹⁵ $G := \mathcal{H}_t \cdot Q$ where \mathcal{H}_t is the group of all horizontal permutations of t and $Q \subset \mathcal{V}_t$ is the subgroup of all such vertical permutations of t which only permute full rows of t with length 2. Then every tensor $T \in \mathcal{T}_{\mathfrak{l}^*} \subseteq \mathcal{T}_r V$ ($\mathfrak{l} = \mathbb{K}[\mathcal{S}_r] \cdot e$, e idempotent) fulfils

(7.1)
$$\sum_{b \in \mathfrak{B}_{b_0}} \gamma_b T_b \in \mathbf{1}_G \cdot \mathbb{K}[\mathcal{S}_r] \cdot e \quad , \quad \mathbf{1}_G := \sum_{g \in G} g \, .$$

¹⁴See the references [17, 18, 16, 24, 20, 21, 14] for the Littlewood-Richardson rule and plethysms.

¹⁵*G* is a semidirect product $\mathcal{H}_t \rtimes Q$ and isomorphic to the wreath product $\mathcal{S}_2 \wr \mathcal{S}_l$.

Furthermore, if dim $V \ge r - l$, then there is such a $b_0 \in \mathcal{B}^{r-2l}$ that ¹⁶

(7.2)
$$1_G \cdot \mathbb{K}[\mathcal{S}_r] \cdot e = \mathcal{L}_{\mathbb{K}} \Big\{ \sum_{b \in \mathfrak{B}_{b_0}} \gamma_b T_b \ \Big| \ T \in \mathcal{T}_{\mathfrak{l}^*} \Big\}.$$

The proof can be found in Sec. III.3.4 of our Habilitationsschrift [10]. If dim V < r-l, then the $\sum_{b \in \mathfrak{B}_{b_0}} \gamma_b T_b$ will span only a linear subspace of $1_G \cdot \mathbb{K}[\mathcal{S}_r] \cdot e$ in general. To describe this subspace, we define:

Definition 7.2. If $\lambda = (\lambda_1, \ldots, \lambda_k) \vdash r$ is a partition with length $|\lambda| = k$ and $(v_1, \ldots, v_k) \in V^k$ is a k-tuple of vectors, then we denote by $\langle \lambda; v_1, \ldots, v_k \rangle$ or short $\langle \lambda; v_i \rangle$ that r-tuple from V^r which has the structure

(7.3)
$$\langle \lambda; v_1, \dots, v_k \rangle := (\underbrace{v_1, \dots, v_1}_{\lambda_1}, \underbrace{v_2, \dots, v_2}_{\lambda_2}, \dots, \underbrace{v_k, \dots, v_k}_{\lambda_k}) \in V^r$$

For every $b = (v_1, \ldots, v_r) \in V^r$, there exists a unique partition $\lambda \vdash r$ and a permutation $q \in S_r$ such that b can be written as $b = q\langle \lambda; w_1, \ldots, w_{|\lambda|} \rangle$ where $w_1, \ldots, w_{|\lambda|}$ are the pairwise different, suitably renumbered vectors from b. We call $\langle \lambda; w_1, \ldots, w_{|\lambda|} \rangle$ a grouping of b and λ the grouping partition of b, which we also denote by $\lambda = b^{\vdash}$.

Definition 7.3. Let \mathcal{B} be an orthonormal basis with respect to a fundamental tensor $g \in \mathcal{T}_2 V$. We call $(n_{i_1}, \ldots, n_{i_{r'}}) \in \mathcal{B}^{r'}$ smaller than $(n_{j_1}, \ldots, n_{j_{r'}}) \in \mathcal{B}^{r'}$ if the first non-vanishing difference $j_k - i_k$ fulfils $j_k - i_k > 0$. If $\langle \lambda; w_1, \ldots, w_{|\lambda|} \rangle$ and $\langle \lambda; w'_1, \ldots, w'_{|\lambda|} \rangle$ are two groupings of a fixed *r*-tuple $b \in \mathcal{B}^r$ of basis vectors, then we call $\langle \lambda; w_1, \ldots, w_{|\lambda|} \rangle$ smaller than $\langle \lambda; w'_1, \ldots, w'_{|\lambda|} \rangle$ if the $|\lambda|$ -tuple $(n_{i_1}, \ldots, n_{i_{|\lambda|}}) := (w_1, \ldots, w_{|\lambda|})$ is smaller than the $|\lambda|$ -tuple $(n_{j_1}, \ldots, n_{j_{|\lambda|}}) := (w'_1, \ldots, w'_{|\lambda|})$.

For every *r*-tuple $b \in \mathcal{B}^r$ there exists a permutation $p_b \in \mathcal{S}_r$ such that *b* has a representation $b = p_b \langle \lambda; w_1, \ldots, w_{|\lambda|} \rangle$ where $\langle \lambda; w_1, \ldots, w_{|\lambda|} \rangle$ is the smallest grouping of *b* and $\lambda = b^{\vdash}$. We denote by **p** a single-valued mapping $\mathbf{p} : \mathcal{B}^r \to \mathcal{S}_r, b \mapsto \mathbf{p}(b) := p_b$ which assigns exactly one of such permutations p_b to *b*.

Let $b_0 \in \mathcal{B}^{r-2l}$ be an (r-2l)-tuple of vectors from the basis \mathcal{B} . We denote by Λ_{b_0} the set $\Lambda_{b_0} := \{\lambda \vdash r \mid \exists b \in \mathfrak{B}_{b_0} : \lambda = b^{\vdash}\}$. Furthermore, we assign to every partition $\lambda \in \Lambda_{b_0}$ the lexicographically smallest standard tableau t_{λ} of λ and the set $\mathcal{M}_{b_0,\lambda} := \{\mathfrak{p}(b)^{-1}b \in \mathcal{B}^r \mid b \in \mathfrak{B}_{b_0} \text{ with } b^{\vdash} = \lambda\}$ of such r-tuples which are the smallest groupings of the $b \in \mathfrak{B}_{b_0}$ with grouping partition λ .

Theorem 7.4. Let $V, \mathcal{B}, r, l, g, b_0, \mathcal{T}_{l^*}$ and e have the meaning given in Theorem 7.1 and $\mathfrak{p} : \mathcal{B}^r \to \mathcal{S}_r$ be a mapping of the type described in Definition 7.3. Then we

¹⁶Corollary: If r - 2l = 0, then the decomposition of $1_G \cdot \mathbb{K}[S_r]$ into minimal right ideals is characterized by a plethysm $[2] \odot [l] \sim \sum_{\mu \vdash l} [2\mu]$. Thus the number I of linearly independent invariants of T is bounded by the sum M of the multiplicities of minimal left ideals \mathfrak{l}_i belonging to partitions 2μ , $\mu \vdash l$, in $\mathbb{K}[S_r] \cdot e = \bigoplus \mathfrak{l}_i$. If dim $V \geq l$, then I = M. (See B. Fiedler [10, Sec.III.4.2]. Compare S.A. Fulling et al. [14].)

have

$$\mathcal{L}_{\mathbb{K}}\left\{\sum_{b\in\mathfrak{B}_{b_{0}}}\gamma_{b}T_{b} \mid T\in\mathcal{T}_{\mathfrak{l}^{*}}\right\} = \sum_{\lambda\in\Lambda_{b_{0}}}\sum_{\langle\lambda;w_{i}\rangle\in\mathcal{M}_{b_{0};\lambda}}a_{\langle\lambda;w_{i}\rangle}\cdot 1_{\mathcal{H}_{t_{\lambda}}}\cdot\mathbb{K}[\mathcal{S}_{r}]\cdot e$$

where

$$a_{\langle \lambda; w_i \rangle} := \sum_{\substack{b \in \mathfrak{B}_{b_0} \\ \mathfrak{p}(b)^{-1}b = \langle \lambda; w_i \rangle}} \gamma_b \, \mathfrak{p}(b)^{-1} \, da_{\langle \lambda; w_i \rangle}$$

The proof is given in Sec. III.3.4 of our Habilitationsschrift [10]. Furthermore, we clear in Sec. III.3.4 to what extent the subspaces from Theorem 7.4 are independent from the choice of b_0 .

8. Concluding Remarks

In Sec. III.3.3 of our Habilitationsschrift [10] we determined linear subspaces $W \subseteq \mathbb{K}[S_r]$ for tensors $T \in \mathcal{T}_r V$ for which the vector space V has a dimension $\dim V < r$. Some details about this case can be found in the paper [12], too.

In Sec. III.4.2 of the Habilitationsschrift [10] we applied Theorem 7.1 to *invariants* (i.e. r = 2l and $b_0 = \emptyset$). Footnote 16 gives one of the results.

We have tested our methods in computer calculations. Among other things, we treated the term combination problem for quadratic monomials in the coordinates of the Riemannian curvature tensor (B. Fiedler [10, Sec.III.5.1]). The correctness of this calculation was controlled by means of the results of S.A. Fulling et al. [14]. Furthermore, we used Theorem 7.4 to verify the standard identity¹⁷ $\sum_{p \in S_{2n}} \chi(p) A_{p(1)} \cdots A_{p(2n)} = 0$ for $(n \times n)$ -matrices A_i in the case n = 2 (B. Fiedler [10, Sec.III.5.2]). For this calculation the full algorithm (R) was required.

We implemented our methods in a Mathematica package PERMS. This package comprises tools for permutation groups and group rings, partitions and tableaux, special idempotents, characters, the Littlewood-Richardson rule, plethysms, discrete Fourier transforms and our algorithms (L) and (R). Furthermore, we used the packages SYMMETRICA [19] and GAP [25].

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¹⁷See S.A. Amitsur and J. Levitzki [1] and S. Bondari [4].

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