

# MATRIX TABLEAU-PAIRS WITH KEY AND SHUFFLING CONDITIONS

OLGA AZENHAS<sup>1</sup> AND RICARDO MAMEDE<sup>2</sup>

CMUC

Department of Mathematics, University of Coimbra  
3001-454 Coimbra, Portugal

ABSTRACT. It has been shown that the sequence of Smith invariants defined by certain sequences of products of matrices, with entries in a local principal ideal domain, are combinatorially described by tableau-pairs  $(T, K)$  where  $T$  is a tableau of skew-shape which rectifies to the key-tableau  $K$ . It is a fact that the set of all shuffles of the columns of a key-tableau is a subset of its Knuth class. Here, under the condition that the word of  $T$  is a shuffle of the columns of the key-tableau  $K$ , we show the converse, that is, every tableau-pair under the aforesaid restrictions has a matrix construction. In the case of a four-letter alphabet, since we are able to give an explicit description of the Knuth class of a key-tableau as a union of the shuffles of certain subsets of words containing the key-tableau columns, our construction is general. This may be seen as an indication of a general procedure if a subset of shuffling generators of a generic key-tableau Knuth class is provided. At the moment however, this seems to be a very difficult problem.

## 1. INTRODUCTION

Given an  $n$  by  $n$  nonsingular matrix  $A$ , with entries in a local principal ideal domain with prime  $p$ , by Gaussian elimination one can reduce  $A$  to a diagonal matrix  $\Delta_a$  with diagonal entries  $p^{a_1}, \dots, p^{a_n}$ , for unique nonnegative integers  $a_1 \geq \dots \geq a_n$ , called the *Smith normal form* of  $A$ . The diagonal entries  $p^{a_1}, \dots, p^{a_n}$  are called the Smith invariants or invariant factors of  $A$ , and  $a = (a_1, \dots, a_n)$  the *invariant partition* of  $A$ . It is known that  $a$ ,  $b$ , and  $c$  are invariant partitions of nonsingular matrices  $A$ ,  $B$ , and  $C$  such that  $AB = C$  if and only if there exists a Littlewood–Richardson tableau  $T$  of type  $(a, b, c)$ , that is, a tableau of shape  $c/a$  which rectifies to the key-tableau of weight  $b$  (also known as Yamanouchi tableau of weight  $b$ ). (See [1, 2, 3, 4, 5, 11, 12, 14, 17, 28].) The relationship between Smith invariants and the product of Schur functions was noticed earlier by several authors [14, 17, 28]. For an overview and other relations, see the survey by W. Fulton [12] as well as [11, 13].

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Given a natural number  $n \geq 1$ ,  $[n]$  denotes the set  $\{1, \dots, n\}$ . Let  $m = (m_1, \dots, m_t)$  be a finite sequence of nonnegative integers. The symmetric group  $\mathfrak{S}_t$  acts on sequences of  $t$  nonnegative integers *via* the left action  $s_i m = (m_1, \dots, m_{i+1}, m_i, \dots, m_t)$  with  $s_i$  the simple transpositions of  $\mathfrak{S}_t$ ,  $1 \leq i \leq t - 1$ . Let  $\alpha(m)$  be the unique partition in the orbit  $\mathfrak{S}_t m$ , and  $\alpha(m)'$  the conjugate partition.  $K(m)$  denotes the key-tableau [21] of weight  $m$ , that is, the tableau of weight  $m$  and *shape*  $\alpha(m)$ , and  $D_{[m_k]}$  denotes the  $n$  by  $n$  diagonal matrix having the  $i$ th diagonal entry equal to  $p$  whenever  $i \in [m_k]$  and 1 otherwise. The conjugate invariant partition of  $D_{[m_k]}$  is the one-row partition  $(m_k)$ . We identify  $K(m)$  with the sequence of diagonal matrices  $(D_{[m_1]}, \dots, D_{[m_t]})$  in the following sense: the sequence of conjugate invariant partitions associated with the sequence of products of matrices

$$D_{[m_1]}, D_{[m_1]}D_{[m_2]}, \dots, D_{[m_1]}D_{[m_2]} \cdots D_{[m_t]}$$

is the nested sequence of partitions  $(m_1) \subseteq \alpha(m_1, m_2) \subseteq \dots \subseteq \alpha(m)$  which defines the key-tableau  $K(m)$  by telling us that for each  $1 \leq i \leq t$  the first  $m_i$  columns contain the letter  $i$ . (We adopt the English convention for Young diagrams.) For instance,

$$K(1032) = \begin{array}{ccc} & 1 & 3 & 3 \\ & 3 & 4 & \\ & 4 & & \end{array} \quad \text{is identified with } (D_{[1]}, D_\emptyset, D_{[3]}, D_{[2]}).$$

Let  $U$  be an  $n$  by  $n$  unimodular matrix, that is, the determinant of  $U$  is not divisible by  $p$ . Put  $\Delta_a UK(m)$  for the sequence  $\Delta_a, \Delta_a U D_{[m_1]}, \Delta_a U D_{[m_1]} D_{[m_2]}, \dots, \Delta_a U D_{[m_1]} D_{[m_2]} \cdots D_{[m_t]}$ . The sequence of conjugate invariant partitions  $a^{0'} = a' \subseteq a^{1'} \subseteq \dots \subseteq a^{t'} = c'$ , associated with that sequence of matrices, is such that the skew-shape  $a^{k+1'}/a^{k'}$  has  $m_{k+1}$  boxes with at most one box in each column,  $k = 0, 1, \dots, t - 1$ . Thus  $\Delta_a UK(m)$  is identified with the tableau-pair  $(T, K(m))$  where  $T$  is the tableau of weight  $m$  defined by that nested sequence of the *conjugate* partitions. It is shown in [6] that  $T$  rectifies to  $K(m)$ . The reverse question arises naturally: *Given a tableau-pair  $(T, K)$  such that  $K$  is a key-tableau and  $T$  is a semi-standard Young tableau that rectifies to  $K$ , does there exist a matrix construction of the form  $\Delta_a UK(m)$  for that pair?* This question has been already answered positively for some instances of the weight in the orbit of  $\mathfrak{S}_t \alpha(m)$  [1, 3, 4, 5, 7]. These are precisely instances where the Knuth class of a key-tableau is equal to the set of shuffles of its columns. In [6], a necessary and sufficient condition has been given for the Knuth class of a key-tableau to be equal to the set of the shuffles of its columns. In particular, the shuffles of the columns of a key-tableau are always contained in its Knuth class. Here, in Section 4, an algorithm is provided for the matrix construction of such a tableau-pair  $(T, K)$  whenever the word of  $T$  is a shuffle of the columns of  $K$ . It is also shown how to extend our matrix construction to any pair  $(T, K)$ , over a four-letter alphabet, whenever  $T$  rectifies to the key-tableau  $K$ . In the case of a four-letter alphabet it is shown in [6] that the Knuth class of a key-tableau is the union of the shuffles of some subsets of words where the set of columns is included.

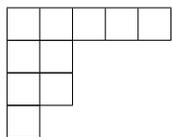
The matrix construction of a tableau-pair  $(T, K)$ , where the word of  $T$  is a shuffle of the columns of the key-tableau  $K$ , relies on the following idea: a semistandard Young

tableau, with a word of length  $\ell$ , can be encoded by a biword  $\Lambda$  without repeated biletters  $\begin{pmatrix} \pi_j \\ x_j \end{pmatrix}$  where  $x_j$  is a letter in column  $\pi_j$  of  $T$ , for  $j = 1, \dots, \ell$ . Among the semistandard Young tableaux with biword  $\Lambda$  there is one in compact form (a tableau where the number of rows of size two of any two consecutive columns is maximal when shifting down the rightmost with respect to the other). Let  $c/a$  be its skew-shape. A semistandard Young tableau with biword  $\Lambda$  has therefore skew-shape  $(c' + e)' / (a' + e)'$ , for some partition  $e$ . As the word of  $T$  is a shuffle of the columns of  $K$ , we may choose a biword  $\Pi$  factorized into biwords whose bottom words are the columns of the key-tableau  $K$ . These factors are identified with east–northeastward paths in the lattice  $\mathbb{N}^2$ , seen as an  $\mathbb{N}$ -matrix, whose vertices  $(\pi_j, x_j)$  are the biletters  $\begin{pmatrix} \pi_j \\ x_j \end{pmatrix}$ . These lattice paths do not intersect and satisfy some properties which allow us to define an injective map to be used in the Main Algorithm in Section 4. The Main Algorithm depends only on the biword  $\Pi$  identified with its lattice path representation. It provides a matrix construction, for any  $T$  with that biword, based on elementary matrices  $T_{ij}(x) = I + xE_{ij}$  whose indices  $\{i j\}$  are generated by the vertices of those lattice paths. If we want a simultaneous matrix construction of  $T$  and  $K$ , we have to select a specific biword that one shows to exist always. This is achieved in Section 5 by relating certain type of vertices in the lattice paths of the biword with the existence of a commutation property when passing an elementary matrix  $T_{ij}(x)$  past another  $T_{ab}(y)$ .

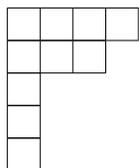
The paper is organized in six sections. In the next section, which in its turn is divided into three subsections, we provide the necessary definitions on tableau and word combinatorics. Special attention is given, in Subsection 2.2, to key-tableaux and, in particular, to those over a four-letter alphabet where the description of the Knuth class is translated into the shuffling operation. In Section 3, divided into two subsections, a semistandard Young tableau is encoded by a biword without repeated letters. When the word of the tableau is a shuffle of the columns of a key-tableau, we choose a biword that factorizes into biwords which are identified with east–northeastward lattice paths in  $\mathbb{N}^2$ . These nonintersecting lattice paths enjoy certain properties that allow us to define in Algorithm 3.1 an injective map that will be an important tool in the next section. Section 4 provides the Main Algorithm which generates a matrix construction of  $T$  whenever the word of  $T$  is a shuffle of the columns of a key-tableau  $K$ . Stated informally, the main theorem (see Theorem 4.4) is: there exists a biword that generates simultaneously a matrix construction of  $T$  and  $K$  whenever the word of  $T$  is a shuffle of the columns of the key-tableau  $K$ . In the succeeding section, a biword obeying certain conditions is exhibited and the proof of the main theorem is given. Finally, in Section 6, it is shown that, over a four-letter alphabet, every tableau-pair  $(T, K)$ , where  $T$  rectifies to the key-tableau  $K$ , has a matrix construction. In the last section, we discuss the generalization of our construction to a  $t$ -letter alphabet, for  $t \geq 5$ . A pertinent problem is the enumeration of the elements of a key-tableau Knuth class which is related with the enumeration of distinct permutation shuffles in the work of C. S. Barnes [8].

## 2. YOUNG TABLEAUX, KEYS, AND SHUFFLES

**2.1. Young tableaux, words, and Knuth equivalence.** A *partition*  $a = (a_1, \dots, a_l)$  is a weakly decreasing sequence of positive integers. We call  $l$  the *length* of  $a$  which is denoted by  $\ell(a)$ . We say that  $a$  is a partition of  $|a| := a_1 + \dots + a_l$ . It is convenient to set  $a_k = 0$  for  $k > \ell(a)$ , and we identify  $a$  with  $(a_1, \dots, a_{\ell(a)}, 0, \dots, 0)$ , where the tail of zeros is of arbitrary length. The unique partition of 0 is denoted by  $(0)$ . A (*weak*) *composition*  $m = (m_1, \dots, m_t)$  is a finite sequence of nonnegative integers. We also say that  $m$  is a weak composition of  $|m| := m_1 + \dots + m_t$ . The symmetric group  $\mathfrak{S}_t$  acts on (weak) compositions with  $t$  entries *via* the left action  $\sigma m := (m_{\sigma^{-1}(1)}, \dots, m_{\sigma^{-1}(t)})$ , with  $\sigma \in \mathfrak{S}_t$ . The unique partition in the orbit  $\mathfrak{S}_t m$  is denoted by  $\alpha(m)$ . We think of a partition in terms of its *Young diagram* consisting of a left-justified array of boxes that, according to the English convention, has  $a_i$  boxes in the  $i$ th row counting from top to bottom. For instance, the Young diagram of  $a = (5, 2, 2, 1)$  is



The *conjugate partition* of  $a$  is the partition  $a'$  whose Young diagram is obtained from  $a$  by interchanging rows with columns. Another way to look at the conjugate of  $a$ , which will be used in the sequel, is to say that if  $a = (a_1, \dots, a_t)$  and  $\ell_k$  is the number of columns of length  $k$  in the Young diagram of  $a$ ,  $1 \leq k \leq t$ , then  $a' = (t^{\ell_t}, \dots, 2^{\ell_2}, 1^{\ell_1})$ , where the notation  $i^{\ell_i}$  means that the positive integer  $i$  appears  $\ell_i$  times in the sequence  $a'$ . Conversely, we have  $a_k = a_{k+1} + \ell_k$ , for  $k = 1, \dots, t$ , with  $a_{t+1} := 0$ . positive integer  $i$  appears  $\ell_i$  times in the sequence  $a'$ . For example, the conjugate partition of  $a = (5, 2, 2, 1)$  is the partition  $(4, 3, 1, 1, 1) = (4, 3, 2^0, 1^3)$ , and its Young diagram is



A partition  $a$  is contained in a partition  $c$ , written  $a \subseteq c$ , if the Young diagram of  $a$  is contained in the Young diagram of  $c$ . In this case, we define the skew-shape  $c/a$  to be the set of boxes in the Young diagram of  $c$  that remains after we remove those boxes corresponding to  $a$ . When  $a = (0)$ , we obtain the Young diagram of  $c$ .

Let  $\mathbb{N}$  be the set of positive integers with the usual order “ $\leq$ ”. If  $t \in \mathbb{N}$ ,  $[t]$  denotes the set  $\{1, \dots, t\}$ , and  $[t]^*$  the free monoid in the alphabet  $[t]$ . That is, the collection of all finite words over the alphabet  $[t]$ , with the concatenation operation. A word over the alphabet  $[t]$  is defined to be a finite string of elements (known as letters) of the set  $[t]$ .

Given the word  $w = x_1 \dots x_k$  over the alphabet  $[t]$ , the *weight* of  $w$ , in the alphabet  $[t]$ , is the vector  $(m_1, \dots, m_t)$ , where  $m_x$  denotes the multiplicity of the letter  $x \in [t]$  in the word  $w$ . Here  $k$  is the *length* of  $w$ , denoted by  $\ell(w)$ . We have  $\ell(w) = m_1 + \dots + m_t$  and, in particular, the length and the weight of the empty word is zero. A *subword* of a word  $w$  is a word obtained by deleting the letters at some (not necessarily

adjacent) positions in  $w$ . As usual, we write  $w^k$ ,  $k \geq 0$ , to mean the concatenation of  $w$  with itself  $k$  times. When one writes  $x_1x_2 \dots x_k = 1^{m_1}2^{m_2} \dots t^{m_t}$  this means that  $w$  is a weakly increasing word, called a *row word*, and  $x_1 = \dots = x_{m_1} = 1$ ,  $x_{m_1+1} = \dots = x_{m_1+m_2} = 2$ ,  $\dots$ ,  $x_{m_1+\dots+m_{t-1}+1} = \dots = x_k = t$ . Similarly we have a weakly decreasing word when we write  $x_1x_2 \dots x_k = t^{m_t} \dots 2^{m_2}1^{m_1}$ . If the letters in  $w$  are in strictly decreasing order, that is,  $x_i > x_{i+1}$  for all  $i$ ,  $w$  is called a *column word*. A column word is often identified with its *support*, that is, with the set consisting of its letters. For example, the word 431 is identified with the set  $\{1, 3, 4\}$ .

A (semistandard) Young tableau  $T$  of shape  $c/a$  (for short skew-tableau) is a filling of the skew-shape  $c/a$  with positive integers, weakly increasing across each row and strictly increasing down each column [10, 19]. If  $a = (0)$  then  $T$  is said to be of partition shape. The (reading) word of  $T$  is the word obtained by reading its columns from bottom to top, starting on the left and moving to the right. The weight of  $T$  is the weight of its word. A Young tableau of partition shape is identified with its (reading) word. The empty tableau is the Young tableau of shape  $(0)$  identified with the empty word.

**Example 2.1.** A (semistandard) Young tableau  $T$  of skew-shape  $(4, 4, 2, 1)/(3, 2)$  with weight  $(1, 2, 2, 1)$  and word 433221,

$$(2.1) \quad \begin{array}{cccc} & & & 1 \\ & & 2 & 2 \\ 3 & 3 & & \\ 4 & & & \end{array} .$$

A Young tableau  $T$  of shape  $c/a$  and weight  $(m_1, \dots, m_t)$  may also be seen as a nested sequence of partitions  $T = (a^0, a^1, \dots, a^t)$ , where  $a = a^0 \subseteq a^1 \subseteq \dots \subseteq a^t = c$ , such that for  $k = 1, \dots, t$ , the skew-diagram  $a^k/a^{k-1}$  has  $m_k$  boxes with at most one box in each column, and it is filled with the letter  $k$  [22]. In Example 2.1, we may write  $T = (a^0, a^1, a^2, a^3, a^4)$ , where  $a^0 = (3, 2) \subseteq a^1 = (4, 2) \subseteq a^2 = (4, 4) \subseteq a^3 = (4, 4, 2) \subseteq a^4 = (4, 4, 2, 1)$ .

The Knuth or plactic congruence  $\equiv [20, 10, 18, 19]$  on the words over the alphabet  $[t]$  is the congruence in  $[t]^*$  defined by the transitive closure of the relations

$$uxzyv \equiv uzxyv, \quad x \leq y < z,$$

$$uyz xv \equiv uyxzv, \quad x < y \leq z,$$

where  $x, y$  and  $z$  are letters in  $[t]$ , and  $u$  and  $v$  are words in  $[t]^*$ . Every Knuth class has one and only one Young tableau of partition shape. Henceforth, each Knuth or plactic class consists of all words Knuth equivalent to the unique Young tableau in the Knuth class. Using *jeu de taquin slides* every semistandard tableau can be *rectified* to the unique tableau in the Knuth class of its word [10, 26, 27]. In Example 2.1,  $T$  is rectified to 4321 32  $\equiv$  433221. Two Young tableaux are said to be Knuth equivalent if they can be obtained from another by *jeu de taquin slides*. Equivalently, they have the same rectification.

The *overlap* of a pair  $(u, v)$  of column words  $u$  and  $v$  is the maximum number of rows of length two that one can obtain by putting the column  $u$  to the left of the column  $v$  so that the two columns together form a skew-tableau. For instance, the

overlap of (43, 21) is zero while the overlap of (43, 32) is one. A skew-tableau is said to be in *compact form* if the number of rows of size two of any two consecutive columns is the overlap of the words comprising those columns. For instance the skew-tableau in Example 2.1 is in compact form. Any skew-tableau may be put in compact form using *jeu de taquin* slides in consecutive rows.

**2.2. Key-tableaux and shuffles.** A *key-tableau*  $K$  is a tableau whose columns are pairwise comparable in the inclusion order [21]. Given a composition  $m$ , the key-tableau of weight  $m$  is the tableau  $K(m)$  of shape  $\alpha(m)$  whose first  $m_j$  columns contain the letter  $j$  for all  $j$ . This defines an obvious bijection between key-tableaux and compositions [24]. A key-tableau is also a tableau whose weight is a permutation of its shape. The key-tableau  $K(m)$  [6, 21] may be written as

$$(2.2) \quad K(m) = v_t^{\ell_t} v_{t-1}^{\ell_{t-1}} \dots v_1^{\ell_1},$$

where  $\{1, \dots, t\} = v_t \supsetneq v_{t-1} \supsetneq \dots \supsetneq v_1 \neq \emptyset$  are column words such that  $\alpha(m) = (\ell_1 + \dots + \ell_t, \dots, \ell_{t-1} + \ell_t, \ell_t)$ , and its conjugate is  $\alpha'(m) = (t^{\ell_t}, \dots, 1^{\ell_1})$ . Let  $\sigma \in \mathfrak{S}_t$ , written as a word  $a_1 a_2 \dots a_t$  in  $[t]^*$ , be such that  $\{a_1, \dots, a_i\}$  is the support of  $v_i$ , for  $i = 1, \dots, t$ . One has  $\sigma\alpha(m) = m = (m_1, \dots, m_t)$ , with  $m_j = \sum_{k=\sigma^{-1}(j)}^t \ell_k$  for all  $j$ . Therefore the key  $K(m)$  may be encoded by a permutation  $\sigma$  and a nonnegative integral vector  $(\ell_t, \dots, \ell_1)$ , that is,  $K(\sigma, (\ell_t, \dots, \ell_1)) := v_t^{\ell_t} v_{t-1}^{\ell_{t-1}} \dots v_1^{\ell_1}$ . (This presentation of a key-tableau will be used in Section 6.)

**Example 2.2.** The key-tableau

$$K(1, 2, 2, 1) = \begin{array}{|c|c|} \hline 1 & 2 \\ \hline 2 & 3 \\ \hline 3 & \\ \hline 4 & \\ \hline \end{array} = 4321\ 32$$

has shape  $(2, 2, 1, 1)$  and weight  $m = (1, 2, 2, 1)$ .

The Knuth class of the word 433221 of the skew-tableau  $T$  in Example 2.1 is the key-tableau  $K(1, 2, 2, 1) = 4321\ 32$  (the empty space means the ending of a column and the starting of a new one) which can be encoded either by permutation 2314 or 3214, and the nonnegative integral vector  $(1, 0, 1, 0)$ .

Let  $w = x_1 \dots x_k \in [t]^*$ . Given  $I = \{i_1 < \dots < i_r\} \subseteq [k]$ ,  $w|I$  denotes the subword  $x_{i_1} \dots x_{i_r}$ . In particular,  $w|I$  is the empty word if  $I = \emptyset$ . Two subwords  $w|I$  and  $w|J$  of  $w$  are said to be disjoint if the sets  $I$  and  $J$  are disjoint.

The word  $w = x_1 \dots x_k$  is said to be a *shuffle* of the words  $u$  and  $v$  if there exists a partition  $\{I, J\}$  of the set  $[k]$  such that  $u = w|I$  and  $v = w|J$ . Let  $u_1, \dots, u_q \in [t]^*$  be  $q$  words of lengths  $k_1, \dots, k_q$ , respectively. If  $q = 0$  the shuffle is the empty word. If  $q \geq 1$ , let  $\{I_1, \dots, I_q\}$  be a partition of the set  $[k]$ . Then the word  $w|\{I_1, \dots, I_q\}$  defined by  $w|I_j = u_j$ , for  $j = 1, \dots, q$ , is a *shuffle* of  $u_1, \dots, u_q$  [15, 25]. The words  $u_1, \dots, u_q$  are said to be the *shuffle components* of  $w|\{I_1, \dots, I_q\}$ . Indeed  $w|\{I_1, \dots, I_q\} = w|\{F_1, \dots, F_q\}$  whenever  $\{I_1, \dots, I_q\} = \{F_1, \dots, F_q\}$ . However, a word may have different shuffle decompositions of the same words, that is,  $w|\{I_1, \dots, I_q\} = w|\{J_1, \dots, J_q\}$  with  $\{I_1, \dots, I_q\} \neq \{J_1, \dots, J_q\}$  set partitions of  $[k]$ . For example, if  $w = 433221 \in [4]^*$ , we have  $w|\{\{1, 2, 5, 6\}, \{3, 4\}\} =$

$w|\{\{1, 2, 4, 6\}, \{3, 5\}\}$ . The set of the shuffles of the  $q$  words  $u_1, \dots, u_q$ , is the set of all words obtained by shuffling together the  $q$  words  $u_1, \dots, u_q$ ,

$$\begin{aligned} & \sqcup\sqcup(u_1, \dots, u_q) \\ & = \{w|\{I_1, \dots, I_q\} : \{I_1, \dots, I_q\} \text{ a set partition of } [k], w|I_j = u_j, 1 \leq j \leq q\}. \end{aligned}$$

Let  $K(m) = v_t^{\ell_t} v_{t-1}^{\ell_{t-1}} \dots v_1^{\ell_1}$  be the key as in (2.2), and  $\sigma \in \mathfrak{S}_t$  such that  $\sigma\alpha(m) = m$ . By Greene's Theorem [16], the set  $\sqcup\sqcup(v_t^{\ell_t}, v_{t-1}^{\ell_{t-1}}, \dots, v_1^{\ell_1})$  is always a subset, in general proper, of the Knuth class of  $K(m)$ . In [6] it was proved that, if  $f_1 \supseteq f_2 \supseteq \dots \supseteq f_k$  are nonempty column words, then the Knuth class of  $K = f_1 \dots f_k$  is equal to the set of all shuffles of its columns if and only if each column is either an interval of  $f_1$ , or is obtained from an interval of  $f_1$  by removing a single letter. It is clear that this property is true for any key  $K(m) = v_t^{\ell_t} v_{t-1}^{\ell_{t-1}} \dots v_1^{\ell_1}$  over the alphabet  $\{1, 2\}$  or  $\{1, 2, 3\}$ . However, over the alphabet  $\{1, 2, 3, 4\}$ , considering  $S := \{1423, 1432, 4123, 4132\} \subseteq \mathfrak{S}_4$ , this property holds if and only if either  $\sigma \notin S$ , or  $\sigma \in S$  but  $\ell_2 = 0$  or  $\ell_4 = 0$ .

Denote by  $\widehat{v}_5^{n_5}$  the multiset consisting of  $n_5$  words  $\widehat{v}_5 := 431421 \equiv K(2, 1, 1, 2)$ . The Knuth class of a key over the alphabet  $\{1, 2, 3, 4\}$  was characterized in [6], in terms of the shuffling operation, as follows.

**Theorem 2.1.** [6] *Let  $\sigma \in \mathfrak{S}_4$ , and let  $(\ell_4, \dots, \ell_1)$  be a sequence of nonnegative integers. The Knuth class of  $K(\sigma, (\ell_4, \dots, \ell_1)) = v_4^{\ell_4} v_3^{\ell_3} v_2^{\ell_2} v_1^{\ell_1}$  is*

$$\begin{cases} \sqcup\sqcup(v_4^{\ell_4}, v_3^{\ell_3}, v_2^{\ell_2}, v_1^{\ell_1}), & \text{if } \sigma \in \mathfrak{S}_4 \setminus S, \text{ or } \sigma \in S \text{ and } (\ell_2 = 0 \text{ or } \ell_4 = 0), \\ \bigcup_{n_5=0}^{\min\{\ell_2, \ell_4\}} \sqcup\sqcup(\widehat{v}_5^{n_5}, v_4^{n_4}, \dots, v_1^{n_1}), & \text{if } \sigma \in S \text{ and } \ell_2 \neq 0, \ell_4 \neq 0, \end{cases}$$

where  $n_i = \ell_i$ ,  $i = 1, 3$ , and  $n_i = \ell_i - n_5$ ,  $i = 2, 4$ .

### 3. BIWORDS AND NONINTERSECTING E-NE LATTICE PATHS

**3.1. Biwords.** Our aim is, given a (skew) semistandard Young tableau in compact form, to encode it by a biword so that, identifying their biletters with points in the lattice  $\mathbb{N}^2$ , seen as an  $\mathbb{N}$ -matrix, one is able to define the elementary matrices of our matrix construction in Section 4. All the tableaux with that biword will be related with the tableau in compact form.

Given  $n, t \in \mathbb{N}$ , a biword over the alphabet  $[n] \times [t]$  is an ordered pair of words of the same length, written as

$$\begin{pmatrix} u \\ v \end{pmatrix} = \begin{pmatrix} \pi_1 & \cdots & \pi_k \\ x_1 & \cdots & x_k \end{pmatrix},$$

with  $u = \pi_1 \dots \pi_k \in [n]^*$  and  $v = x_1 \dots x_k \in [t]^*$ . Each biword  $\begin{pmatrix} u \\ v \end{pmatrix}$  can also be

seen as a word whose letters are the biletters  $\begin{pmatrix} \pi_1 \\ x_1 \end{pmatrix}, \dots, \begin{pmatrix} \pi_k \\ x_k \end{pmatrix}$ . We only consider

biwords with pairwise distinct biletters, that is,  $\begin{pmatrix} \pi_i \\ x_i \end{pmatrix} \neq \begin{pmatrix} \pi_j \\ x_j \end{pmatrix}$ , for every  $i \neq j$ . Two

biwords are said to be *equivalent* if they have the same set of biletters. Among the

biwords having the same set of biletters as  $\begin{pmatrix} u \\ v \end{pmatrix}$  we consider two special ones,  $\Sigma$  and  $\Sigma'$ . The weakly increasing rearrangement of  $\begin{pmatrix} u \\ v \end{pmatrix}$  using the anti-lexicographic order on biletters with priority in the top entry [19], that is,  $\begin{pmatrix} \pi \\ x \end{pmatrix} \leq \begin{pmatrix} \pi' \\ x' \end{pmatrix}$  if and only if  $\pi < \pi'$ , or  $\pi = \pi'$  and  $x \geq x'$ , gives the biword

$$(3.1) \quad \Sigma = \begin{pmatrix} 1^{f_1} & \cdots & n^{f_n} \\ w_1 & \cdots & w_n \end{pmatrix},$$

for some integers  $f_i \geq 0$ , such that  $w = w_1 \cdots w_n \in [t]^*$ , with  $w_i$  a column word of length  $f_i$ . The weakly decreasing rearrangement of  $\begin{pmatrix} u \\ v \end{pmatrix}$  using the lexicographic order in biletters with priority on the bottom entry [19], that is,  $\begin{pmatrix} \pi \\ x \end{pmatrix} \geq \begin{pmatrix} \pi' \\ x' \end{pmatrix}$  if and only if  $x > x'$ , or  $x = x'$  and  $\pi \geq \pi'$ , gives the biword

$$(3.2) \quad \Sigma' = \begin{pmatrix} J_t & \cdots & J_1 \\ t^{m_t} & \cdots & 1^{m_1} \end{pmatrix},$$

for some integers  $m_i \geq 0$ , such that  $J = J_t \cdots J_1 \in [n]^*$ , with  $J_i$  a column word of length  $m_i$ . We have then the word-pair  $(J, w) \in [n]^* \times [t]^*$  such that the weight of  $w$  is equal to  $(m_1, \dots, m_t) = (\#J_1, \dots, \#J_t)$  and the weight of  $J$  is  $(\ell(w_1), \dots, \ell(w_n))$ . (Recall that column words are identified with their supports.)

Young tableaux in compact form are in one to one correspondence with classes of biwords without repeated biletters. Young tableaux with the same biword have the same word pair  $(J, w)$ , hence their skew-shapes are  $(c' + e)' / (a' + e)'$ , where  $e$  is any partition of length at most  $n$ , and  $c/a$  is the skew-shape of the compact tableau with

$$\begin{array}{c} \boxed{1} \\ \boxed{2} \ \boxed{2} \\ \boxed{3} \ \boxed{3} \\ \boxed{4} \end{array} \quad \text{and} \quad \begin{array}{c} \boxed{1} \\ \boxed{2} \ \boxed{2} \\ \boxed{3} \\ \boxed{3} \\ \boxed{4} \end{array}$$

that biword. For instance,  $\begin{array}{c} \boxed{1} \\ \boxed{2} \ \boxed{2} \\ \boxed{3} \ \boxed{3} \\ \boxed{4} \end{array}$  and  $\begin{array}{c} \boxed{1} \\ \boxed{2} \ \boxed{2} \\ \boxed{3} \\ \boxed{3} \\ \boxed{4} \end{array}$  have the same biword. Let  $T$  be a tableau of skew-shape with  $n$  columns (possibly some are empty), where the word  $w$  has weight  $m = (m_1, \dots, m_t)$ . Let  $\begin{pmatrix} \pi_j \\ x_j \end{pmatrix}$  be the billetter where  $x_j$  is the letter in column  $\pi_j$  of  $T$ , for  $j = 1, \dots, \ell(w)$ . Any biword over the alphabet  $[n] \times [t]$  with these biletters is said to be a *biword of  $T$* . In particular,  $\Sigma = \begin{pmatrix} 1^{f_1} & \cdots & n^{f_n} \\ w_1 & \cdots & w_n \end{pmatrix}$  is the

biword of  $T$  where the factor  $\begin{pmatrix} i^{f_i} \\ w_i \end{pmatrix}$  indicates that  $w_i$  is a column word of length  $f_i$  filled in column  $i$ , counting from left to right, of the skew-shape  $T$ . For each  $i$  in  $[t]$ ,  $J_i \in [n]^*$  is the column word defined by the indices of the columns of  $T$  where the  $m_i$  letters  $i$  of  $w$  are placed. The column words  $J_1, \dots, J_t$  are called the *indexing sets* of  $T$  and  $J = J_t \cdots J_1$  the *indexing-set word*.

A biword class, over the alphabet  $[n] \times [t]$ , without repeated biletters, shall be represented in the lattice  $\mathbb{N}^2$  by identifying its set of biletters  $\binom{y}{i}$  with the set of points  $(y, i)$ , where  $y \in J_i$ , for  $i = 1, \dots, t$ , called the *lattice diagram* of the biword class. Those points, graphically represented as  $\bullet$ , are called *lattice points* or *vertices* of a biword in the class. For convenience, in drawing the lattice  $\mathbb{N}^2$ , we adopt the matrix convention, that is, the first coordinate, the row index, increases as one goes downwards, and the second coordinate, the column index, increases as one goes from left to right.

**Example 3.1.** The skew-tableau  $T$  in Example 2.1 has biword  $\Sigma = \begin{pmatrix} 11 & 2 & 3 & 44 \\ 43 & 3 & 2 & 21 \end{pmatrix}$  and the corresponding lattice diagram, contained in  $\{1, 2, 3, 4\} \times \{1, 2, 3, 4\}$ , is

$$(3.3) \quad T = \begin{array}{|c|c|c|} \hline & & 1 \\ \hline 3 & 3 & 2 \\ \hline 4 & & 2 \\ \hline \end{array} \longleftrightarrow \begin{array}{|c|c|} \hline 3 & 4 \\ \hline 3 & \\ \hline 2 & \\ \hline 1 & 2 \\ \hline \end{array} \quad \left( \begin{array}{cccc} 11 & 2 & 3 & 44 \\ 43 & 3 & 2 & 21 \end{array} \right) \longleftrightarrow \begin{array}{|c|c|c|c|} \hline 1 & 2 & 3 & 4 \\ \hline 1 & & \bullet & \bullet \\ \hline 2 & & \bullet & \\ \hline 3 & & \bullet & \\ \hline 4 & \bullet & \bullet & \\ \hline \end{array} .$$

Note that the lattice vertex marginal cardinalities, by rows respectively columns, are  $(2, 1, 1, 2)$ , which are the corresponding lengths of the columns of  $T$ , and  $(1, 2, 2, 1)$ , which are the corresponding multiplicities of each letter in  $T$ .

**3.2. Nonintersecting E-NE lattice paths.** In this section we consider biwords where the bottom word  $w$  of  $\Sigma$  is in the Knuth class of the key-tableau (2.2) and is a shuffle of its columns, that is,  $w$  is in the set  $\sqcup\sqcup(v_t^{\ell_t}, \dots, v_1^{\ell_1})$ . We are interested in the biwords  $\Pi$  equivalent to  $\Sigma$  with bottom word the key  $v_t^{\ell_t} \cdots v_1^{\ell_1}$ . They are in one to one correspondence with the shuffle decompositions of  $w$  in  $\sqcup\sqcup(v_t^{\ell_t}, \dots, v_1^{\ell_1})$ . Let  $\{F^{t, \ell_t}, \dots, F^{t, 1}, \dots, F^{1, \ell_1}, \dots, F^{1, 1}\}$  be a set partition of  $\{1, \dots, \sum_{k=1}^t k \ell_k\}$  such that  $\Sigma|F^{k, j} = \binom{I^{k, j}}{v_k}$ , for  $k \in [t]$  and  $j \in [\ell_k]$ , which generates the biword

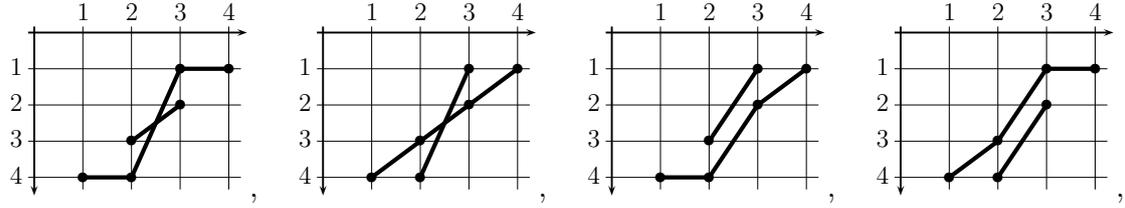
$$(3.4) \quad \Pi = \left( \begin{array}{cccccccccccc} I^{t, \ell_t} & \cdots & I^{t, 1} & \cdots & I^{2, \ell_2} & \cdots & I^{2, 1} & I^{1, \ell_1} & \cdots & I^{1, 1} \\ v_t & \cdots & v_t & \cdots & v_2 & \cdots & v_2 & v_1 & \cdots & v_1 \end{array} \right),$$

where each  $I^{j, i}$  is a weakly increasing word with  $|I^{j, i}| = j$ , for  $j = 1, \dots, t$ , and  $i = 1, \dots, \ell_j$ . In the lattice diagram of  $\Pi$  we link the vertices of any two consecutive biletters in the factor  $\binom{I^{j, i}}{v_j}$  of  $\Pi$ , for  $1 \leq i \leq \ell_j$ ,  $1 \leq j \leq t$ , with a straight line.

Considering the last and the first biletters of this factor, respectively, as the starting and the ending points, one gets a lattice path with  $j$  vertices made of eastward E and northeastward NE steps. These lattice paths do not have vertices in common since the biletters of  $\Pi$  are pairwise distinct. This means that the biword  $\Pi$  (3.4) is graphically represented by  $\ell_t + \cdots + \ell_1$  nonintersecting northeastward-eastward E-NE lattice paths, where  $\ell_j$  of them have length  $j$ , for  $j = 1, \dots, t$ . Conversely, every word in  $\sqcup\sqcup(v_t^{\ell_t}, \dots, v_1^{\ell_1})$  may be obtained in this way using the points of the lattice diagram of  $\Pi$  to draw all the  $\ell_t + \cdots + \ell_1$  nonintersecting northeastward-eastward lattice paths, where  $\ell_j$  of them have  $j$  vertices, for  $j = 1, \dots, t$ . For the rest of the

paper, we identify a biword  $\Pi$  with its lattice path representation in  $\mathbb{N}^2$  where each factor  $\begin{pmatrix} I^{j,i} \\ v_j \end{pmatrix}$ , for  $i = 1, \dots, \ell_j$ ,  $j = 1, \dots, t$ , is seen as the lattice path of length  $j$  whose vertices are the biletters of this factor.

**Example 3.2.** The word in the second row of the biword  $\Sigma$  in (3.3) is a shuffle of 4321 and 32. We may sort the biletters of  $\Sigma$  in several ways, in order to obtain all biwords  $\Pi$  as in (3.4), which is equivalent to draw in the lattice diagram (3.3) all the sets of two nonintersecting paths of lengths four and two using northeast-eastward steps. We have then four types of northeast-eastward nonintersecting paths:



each corresponding to the biwords

(3.5)

$$\Pi = \begin{pmatrix} 1144 & 23 \\ 4321 & 32 \end{pmatrix}, \Pi' = \begin{pmatrix} 1234 & 14 \\ 4321 & 32 \end{pmatrix}, \Pi^\dagger = \begin{pmatrix} 1244 & 13 \\ 4321 & 32 \end{pmatrix} \text{ and } \Pi^\ddagger = \begin{pmatrix} 1134 & 24 \\ 4321 & 32 \end{pmatrix},$$

respectively.

*Remark 3.1.* As the  $v_i$ 's are column words, the lattice paths do not have northward steps. Recalling that the column words of a key-tableau are pairwise comparable for the inclusion order, the sets of column indices of any two paths of  $\Pi$  are comparable for the inclusion order. In particular, any pair of two lattice paths of  $\Pi$  is such that, if one of the paths has a vertex in a column of  $\mathbb{N}^2$  and the other has not, then the set of the indices of the columns of the latter is contained in the set of the former. The lattice paths with the longest length have vertices in all columns of the lattice diagram of  $\Pi$ . We may therefore identify  $\Pi$  with a collection of nonintersecting E-NE paths in  $[n] \times [t]$  such that the sets of the indices of the columns are comparable for the inclusion order.

We still have to fix more terminology. Given two lattice paths  $u$  and  $v$  having each a vertex in column  $k$ , we say that  $u$  is *above* [respectively *below*]  $v$  in column  $k$  if the vertices  $(a, k) \in u$  and  $(x, k) \in v$  satisfy  $a < x$  [respectively  $a > x$ ]. Given two consecutive vertices  $(a, b)$  and  $(x, y)$  of a lattice path, we say that  $(a, b)$  is *positively-linked* to  $(x, y)$  if  $b < y$  (recall that we have always  $a \geq x$ ). Otherwise,  $(a, b)$  is said to be *negatively-linked* to  $(x, y)$ . If  $(a, b)$  is positively-linked to  $(x, y)$  and  $a > x$ , then  $(a, b)$  is called an *obstacle*. An obstacle is a vertex where a northeast step starts. A lattice path without obstacles is a lattice path with only east steps, that is, all the vertices of the path are in the same row.

For instance, in the lattice representation of  $\Pi'$  (3.5), the path  $\begin{pmatrix} 1234 \\ 4321 \end{pmatrix}$  is above the path  $\begin{pmatrix} 14 \\ 32 \end{pmatrix}$  in column 2, and below it in column 3. The vertex  $(4, 1)$  is an obstacle NE-linked to  $(3, 2)$ , and thus the vertex  $(3, 2)$  is negatively linked to  $(4, 1)$ .

We define now a map  $s : \Pi \longrightarrow \mathbb{N}$  which assigns to each vertex of  $\Pi$  the index of a specific row in the lattice diagram of  $\Pi$ . This map will be used in the matrix construction given by the Main Algorithm in Section 4. The algorithm below defining this map is based on the following two observations which are a consequence of Remark 3.1. Pick a vertex  $(\tilde{a}, \tilde{b})$  in  $\Pi$ . If  $(\tilde{a}, \tilde{b})$  is the initial vertex of a path, the first vertex of  $\Pi$  (if any) that one meets when one moves to the west is an obstacle. If  $(\tilde{a}, \tilde{b})$  is not the initial vertex of a path, and if the first vertex of  $\Pi$  one meets, when one moves to the west, does not belong to the path containing  $(\tilde{a}, \tilde{b})$ , then  $(\tilde{a}, \tilde{b})$  is negatively linked to an obstacle. Moreover, if the first vertex, that one has met in the west, is not an obstacle but it is positively linked, say to  $(\tilde{a}, \tilde{f})$ , then  $(\tilde{a}, \tilde{b})$ , or some vertex  $(\tilde{a}, \tilde{d})$ ,  $\tilde{d} \geq \tilde{b}$ , in the same path, is itself an obstacle in a path with a vertex in column  $\tilde{f} > \tilde{d} \geq \tilde{b}$ . See Example 3.5.

**Algorithm 3.1.** Consider the biword  $\Pi$  in (3.4) and its lattice path representation. Let  $(a, b)$  be a vertex of  $\Pi$ . Let  $y := 0$  if  $(a, b)$  is the initial vertex of a path, otherwise, let  $y$  be the index of the column containing the vertex positively linked to  $(a, b)$ . Write  $(\tilde{a}, \tilde{b}) := (a, b)$ , set  $x := \tilde{a}$  and implement the following algorithm:

Begin

1. Starting in  $(\tilde{a}, \tilde{b})$  move one step  $(1, 0)$  in  $\mathbb{N}^2$  to the west. Do  $\tilde{b} := \tilde{b} - 1$ .  
 If  $\tilde{b} = y$ , then  $x := \tilde{a}$  and stop.  
 Else  $\tilde{b} > y$  and go to 2.
  2. If  $(\tilde{a}, \tilde{b})$  is an obstacle of  $\Pi$ , then move northeast along the path linking positively the obstacle  $(\tilde{a}, \tilde{b})$  to the vertex  $(\tilde{c}, \tilde{d}) \in \Pi$ , with  $\tilde{d} \leq \tilde{b}$ . Set  $(\tilde{a}, \tilde{b}) := (\tilde{c}, \tilde{d})$  and go to 1.  
 Else go to 3.
  3. If  $(\tilde{a}, \tilde{b})$  is positively linked to some vertex  $(\tilde{a}, \tilde{d})$  with  $\tilde{b} + 1 < \tilde{d}$ , then set  $\tilde{b} := \tilde{d}$  and go to 1.  
 Else, go to 1.
- End.

*Remark 3.2.* When applying the above algorithm, we always move to the west until we reach column  $y$ , except when we meet an obstacle, as in step 2, in which case we move northeast, or we meet a vertex positively linked to some other vertex in the same row to the east of our current position, as in step 3, in which case we move east. After an east move, in step 3, we will be forced to make a northeast move along an obstacle. Since the number of obstacles is finite, the algorithm must terminate.

The last point in  $\mathbb{N} \times \mathbb{Z}_{\geq 0}$  that one reaches in the algorithm is  $(x, y)$  which may or may not be in  $\Pi$ . Nevertheless,  $x$  is the row index of the last vertex of  $\Pi \setminus \{(x, y)\}$ .

**Definition 3.1.** Let  $(a, b) \in \Pi$ . The sequence of vertices of  $\Pi \setminus \{(x, y)\}$  that one passes in the algorithm above from the vertex  $(a, b) \in \Pi$  to vertex  $(x, y) \in \mathbb{N} \times \mathbb{Z}_{\geq 0}$  is called the  $(a, b)$ -path.

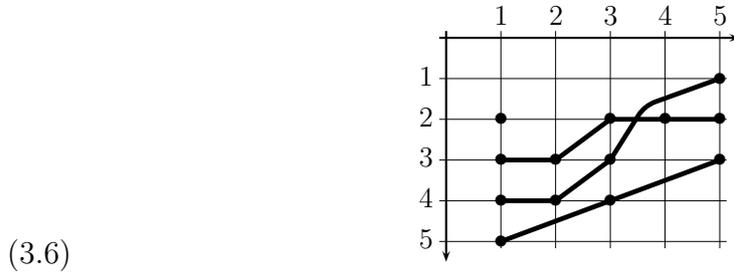
The  $(a, b)$ -path contains at least the vertex  $(a, b) \in \Pi$ , and has odd length. This procedure defines the map

$$s : \Pi \longrightarrow \mathbb{N},$$

where  $s(a, b) := x$  is the row index of the last vertex in the  $(a, b)$ -path.

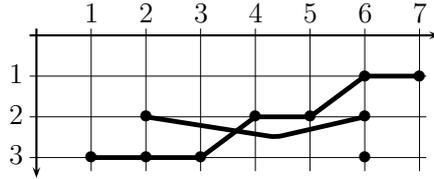
**Example 3.3.** Consider the biword  $\Pi$  in (3.5). The vertices in one row always belong to a unique path. Thus, we have always  $s(a, b) = a$  for every vertex  $(a, b)$  in  $\Pi$ . Consider now  $\Pi'$  in  $w(3.5)$ . There exists a row containing vertices belonging to different paths. To compute the  $(4, 2)$ -path, notice that  $(4, 1)$  is an obstacle. Thus, one must move along the line linking this obstacle to  $(3, 2)$ . Since there are no more obstacles in row 3 to the west of  $(3, 2)$ , the  $(4, 2)$ -path is  $(4, 2), (4, 1), (3, 2)$  and  $s(4, 2) = 3$ . Starting from any other vertex  $(a, b)$  of  $\Pi$ , we will face no obstacle, and thus  $s(a, b) = a$ .

**Example 3.4.** Consider the biword  $\Pi = \begin{pmatrix} 22233 & 1344 & 345 & 2 \\ 54321 & 5321 & 531 & 1 \end{pmatrix}$  whose lattice path representation is



We have  $s(4, 3) = 2$  since the obstacle  $(4, 2)$  is linked to a vertex in row 3, and  $(3, 2)$  is another obstacle that appears to its left, linked to a vertex in row 2. The  $(4, 3)$ -path is  $(4, 3), (4, 2), (3, 3), (3, 2), (2, 3)$  of length five. Similarly, we find that  $s(3, 3) = 2$ ,  $s(3, 5) = 1$ , and  $s(2, 1) = 1$  whose  $(2, 1)$ -path consists only of  $(2, 1)$ . For any other vertex  $(a, b)$  of  $\Pi$ , we get  $s(a, b) = a$ .

**Example 3.5.** Consider now the following lattice path representation:



We have  $s(3, 6) = s(2, 6) = 1$ . For any other vertex  $(a, b)$  of  $\Pi$ , we get  $s(a, b) = a$ .

**Lemma 3.1.** Let  $F$  be the set of all initial vertices of lattice paths of  $\Pi$ . Then the map  $s : F \rightarrow \mathbb{N}$  sending  $(a, b) \in F$  into  $s(a, b)$  is an injection.

*Proof.* Let  $A$  and  $B$  denote the  $(a, b)$  and  $(c, d)$ -paths, respectively, where  $(a, b)$  and  $(c, d)$  are distinct initial vertices of lattice paths in  $F$ . We will show that, after any common vertex, the two paths  $A$  and  $B$  follow distinct directions: one goes west and the other goes northeast, continuing their way in different rows. Thereby  $A$  and  $B$  can not have the last  $\Pi$  vertex in the same row. If  $A$  and  $B$  would have the last  $\Pi$  vertices in the same row, they should coincide. But then one goes west and stops in this row, and the other goes northeast and stops in a row above. This is absurd. Hence,  $s(a, b) \neq s(c, d)$ .

Let  $(e, f)$  in  $A \cap B$ . The vertex  $(e, f)$  appears in  $A$  either if  $e = a$ , or if it is negatively linked to an obstacle in  $A$ , or if it is negatively linked to a vertex  $(e, f - k)$ ,  $k \geq 2$ , in  $A$ . In the first case, the vertex  $(e, f)$  must be an obstacle in  $A$ , and therefore when constructing the path  $B$  we must move northeast by step 2, while in  $A$  we move west. In the second case, in the path  $A$  we move west by step 1, while in the path  $B$  we move northeast, since  $(e, f)$  must be an obstacle. If  $(e, f)$  is the first vertex where the two paths meet, then  $(e, f)$  has to be an obstacle. By induction assume that the paths  $A$  and  $B$  have never shared a common line connecting two consecutive vertices of a path in  $\Pi$ . Finally, in the third case, the path  $A$  arrives at vertex  $(e, f)$  after passing through a vertex, say  $(e, y)$  with  $f - k < y < f$ , negatively linked to an obstacle in  $A$ . After  $(e, f)$  we go to west by step 1 and then northeast. If  $(e, f) \in B$ , then the lattice path  $B$  arrives at this vertex by step 1. After this vertex, either we go northeast along an obstacle, or go east if  $(e, f)$  is positively linked to some vertex  $(e, f + k')$ . In any case, after these steps, we will move northeast using a lattice path distinct from the one used in  $A$ . Therefore, after  $(e, f)$  the two paths  $A$  and  $B$  follow distinct directions, and, in particular, the paths  $A$  and  $B$  never share a common line connecting two consecutive vertices of a path in  $\Pi$ . As the number of vertices is finite it follows that  $s(a, b) \neq s(c, d)$ .  $\square$

#### 4. MATRIX TABLEAU-PAIRS WITH KEY AND SHUFFLE CONDITIONS: STATEMENT OF RESULTS

Let  $\mathcal{R}_p$  be a local principal ideal domain with maximal ideal  $(p)$ . A unit in  $\mathcal{R}_p$  is an element that is not divisible by  $p$ . Every nonzero element  $x$  may be written in the form  $\mu p^k$  with  $k \geq 0$  and  $\mu$  a unit in  $\mathcal{R}_p$ , both uniquely determined by  $x$ . Two elements  $x$  and  $y$  are said to be *associated* if  $x = \mu y$  for some unit  $\mu$  in  $\mathcal{R}_p$ . In this paper, all matrices are  $n$  by  $n$  and nonsingular with entries over  $\mathcal{R}_p$ . Given a matrix  $A$ ,  $A^T$  denotes its transpose. A matrix is said to be unimodular if its determinant is a unit. Let  $\mathcal{U}_n$  be the group of  $n$  by  $n$  unimodular matrices over  $\mathcal{R}_p$ . Given  $n$  by  $n$  matrices  $A$  and  $B$ , we say that  $B$  is *left equivalent* to  $A$  (written  $B \sim_L A$ ) if  $B = UA$  for some  $U \in \mathcal{U}_n$ ;  $B$  is *right equivalent* to  $A$  (written  $B \sim_R A$ ) if  $B = AV$  for some  $V \in \mathcal{U}_n$ ; and  $B$  is *equivalent* to  $A$  (written  $B \sim A$ ) if  $B = UAV$  for some  $U, V \in \mathcal{U}_n$ . The relations  $\sim_L$ ,  $\sim_R$  and  $\sim$  are equivalence relations in the set of all  $n$  by  $n$  nonsingular matrices over  $\mathcal{R}_p$ . (For more details see [23].)

Let  $A$  be an  $n$  by  $n$  nonsingular matrix. By the Smith normal form theorem, Theorem [9, 23], there exist nonnegative integers  $a_1, \dots, a_n$  with  $a_1 \geq \dots \geq a_n \geq 0$  such that  $A$  is equivalent to the diagonal matrix  $\text{diag}_p(a_1, \dots, a_n)$  having  $p^{a_i}$  in the  $i$ -th diagonal position. The sequence  $a = (a_1, \dots, a_n)$  of the exponents of the  $p$ -powers in the Smith normal form of  $A$  is a partition of length  $\leq n$ , uniquely determined by the matrix  $A$ . We call  $a$  the *invariant partition* of  $A$ . Given a subset  $J \subseteq [n]$ , put  $D_J := \text{diag}_p(\chi^J)$ , where  $\chi^J$  is the nonnegative integer vector having 1 at position  $i$  if  $i \in J$ , and 0 otherwise. If  $\sigma \in \mathfrak{S}_n$ ,  $P_\sigma$  denotes the permutation matrix having  $\delta_{i\sigma(j)}$  in position  $(i, j)$ . As usual, if  $u$  and  $v$  are positive integers,  $(u \ v)$  denotes the transposition of  $u$  and  $v$ .

Let  $(i, j) \in [n] \times [n]$ . We denote by  $E_{ij}$  the  $n$  by  $n$  matrix having a 1 in position  $(i, j)$  and 0's elsewhere, and we define the elementary unimodular matrices  $T_{ij}(x)$  as

follows:

$$\begin{aligned} T_{ij}(x) &= I + xE_{ij}, \quad \text{where } i \neq j \text{ and } x \in \mathcal{R}_p; \\ T_{ii}(v) &= I + (v - 1)E_{ii}, \quad \text{where } v \text{ is a unit of } \mathcal{R}_p. \end{aligned}$$

It is obvious that  $E_{ij}E_{rs} = \delta_{jr}E_{is}$ . Therefore, if  $i \neq j$  and  $r \neq s$ , we find that  $T_{ij}(x)T_{rs}(y) = I + xE_{ij} + yE_{rs} + xy\delta_{jr}E_{is}$ . The lemma below states two kinds of rules to be used in the sequel: (I) for commuting either two elementary matrices  $T_{ij}(x)$ , or an elementary matrix  $T_{ij}(x)$  with an elementary diagonal matrix  $D_{[m]}$ ; (II) is concerned with the appearance of a permutation matrix when one elementary matrix passes past another. The proof is left to the reader.

**Lemma 4.1.** *Let  $i, j, r, s, m \in [n]$ , and  $x, y, v \in \mathcal{R}_p$ , such that  $v$  is a unit. Then,*

(I) (Commutation rules)

(i)  $T_{ij}(x)T_{rs}(y) = T_{rs}(y)T_{ij}(x)$ , whenever  $i \neq s$  and  $j \neq r$ .

(ii)  $T_{ij}(x)T_{js}(y) = T_{js}(y)T_{ij}(x)T_{is}(xy)$ , if  $i \neq s$ ,  $i \neq j$ .

(iii)  $T_{ii}(v)T_{rs}(x) = T_{rs}(ux)T_{ii}(v)$ , for some unit  $u$ .

(iv)  $T_{ji}(y)T_{ij}(xp) = T_{ij}(u_1xp)T_{ii}(u_2)T_{jj}(u_3)T_{ji}(u_4y)$   
 $= T_{ij}(u_1xp)T_{ji}(u_5y)T_{ii}(u_2)T_{jj}(u_3)$ , for some units  $u_i$ ,  $i = 1, \dots, 4, 5$ .

(v)  $T_{ij}(x)D_{[m]} = D_{[m]}T_{ij}(x)$ , if  $i, j > m$ .

(vi)  $T_{ij}(x)D_{[m]} = D_{[m]}T_{ij}(xp)$ , if  $i > m \geq j \geq 1$ .

(vii)  $T_{ij}(xp)D_{[m]} = D_{[m]}T_{ij}(x)$ , if  $j > m \geq i \geq 1$ .

(II) (Anti-commutation rule)

(viii)  $T_{ji}(-1)T_{ij}(1) = T_{jj}(-1)T_{ij}(1)P_{(ij)}$ ,  $i \neq j$ .

The lemma says that there exists always a unimodular elementary matrix  $E$  such that

$$T_{ji}(\tau)T_{ab}(\tau') = T_{ab}(\varrho')T_{ji}(\varrho)E, \quad \text{with } \varrho, \tau \text{ and } \varrho', \tau' \text{ pairs of associated elements,}$$

whenever  $\tau, \tau'$  are not both unities, or  $(a, b) \neq (i, j)$  if  $i \neq j$ . Otherwise, the passage of a matrix  $T_{ij}(\tau)$ , with  $i \neq j$ , to the left of a matrix  $T_{ji}(\tau')$ , with  $\tau$  and  $\tau'$  both unities, leads to the appearance of a nonidentity permutation matrix.

Following [3, 5, 7], we introduce the definition of a matrix realization of a pair of tableaux  $(T, F)$ , where  $T$  and  $F$  have the same weight  $m$  and  $F$  has partition shape  $b$ . Let  $B_r$  be a matrix with invariant partition  $(1^{m_r}, 0^{n-m_r})$ , for  $r = 1, \dots, t$ . If  $a^r$  is the conjugate of the invariant partition of  $A_0B_1 \cdots B_r$ ,  $0 \leq r \leq t$ , then it is clear that  $(a^0 = a, a^1, \dots, a^t = c)$  is a tableau with weight  $m = (m_1, \dots, m_t)$  and shape  $c/a$ . Similarly, if  $b^r$  is the conjugate of the invariant partition of  $B_1 \cdots B_r$ ,  $1 \leq r \leq t$ , then  $(b^1, \dots, b^t = b)$  is a tableau of shape  $b$  with weight  $(m_1, \dots, m_t)$  as well (see [7]).

**Definition 4.1.** Let  $T = (a^0, a^1, \dots, a^t)$  and  $F = (0, b^1, \dots, b^t)$  be tableaux of weight  $m = (m_1, \dots, m_t)$ . We say that a sequence of  $n$  by  $n$  nonsingular matrices  $A_0, B_1, \dots, B_t$  is a matrix realization of the pair  $(T, F)$  (or realizes  $(T, F)$ ) if:

I. For each  $r \in \{1, \dots, t\}$ , the invariant partition of the matrix  $B_r$  is  $(1^{m_r}, 0^{n-m_r})$ .

II. For each  $r \in \{0, 1, \dots, t\}$ , the invariant partition of the matrix  $A_r := A_0B_1 \cdots B_r$  is the conjugate of  $a^r$ .

III. For each  $r \in \{1, \dots, t\}$ , the invariant partition of the matrix  $B_1 \cdots B_r$  is the conjugate of  $b^r$ .

The tableau-pair  $(T, F)$  is said to be a *matrix-tableau pair*.

**Example 4.1.** Let  $a = (3, 2, 1)$  and  $F = 11122$ . The sequences

$$\text{diag}_p(a), UD_{[3]}, D_{\{4,5\}},$$

with  $U$  running over the unimodular matrices of order five, give rise to tableaux  $T$  of skew-shape whose words  $w$  are congruent with tableaux  $P$  running over  $\{11122; 21112; 21211\}$  the set of all tableaux of weight  $(3, 2)$  with partition shape. For instance, for  $U = I$ ,  $U = P_{4321}$ , and  $U = P_{12543}P_{4321}$ , respectively, we have:

$$\begin{aligned} (a) \quad T = \begin{array}{c} \boxed{2} \boxed{2} \\ \boxed{1} \\ \boxed{1} \\ \boxed{1} \end{array}, \quad w = 11122; \quad (b) \quad T = \begin{array}{c} \boxed{1} \boxed{2} \\ \boxed{1} \\ \boxed{1} \\ \boxed{2} \end{array}, \quad w = 21112; \text{ and} \\ (c) \quad T = \begin{array}{c} \boxed{1} \boxed{1} \\ \boxed{2} \\ \boxed{1} \\ \boxed{2} \end{array}, \quad w = 21211. \end{aligned}$$

Actually, I and II, in Definition 4.1, say that the invariant partitions of  $A_0, A_0B_1, \dots, A_0B_1 \cdots B_t$  define  $T$ , while III says that the invariant partitions of  $B_1, B_1B_2, \dots, B_1B_2 \cdots B_t$  define  $F$ . Recall that the key  $K(m)$  is the only tableau of weight  $m = (m_1, \dots, m_t)$  and conjugate shape  $\sum_{i=1}^t (1^{m_i})$ , and when written as a sequence of partitions it takes the form  $K(m) = (0, (1^{m_1}), \sum_{i=1}^2 (1^{m_i}), \dots, \sum_{i=1}^t (1^{m_i}))$ .

Thus, when  $F = K(m)$ , in order to verify property III, it is sufficient to show that  $B_1 \cdots B_t$  has invariant partition  $(1^{m_1}) + \cdots + (1^{m_t})$ . For the purpose of this paper, we shall consider only pairs  $(T, K(m))$  such that the word of  $T$  is a shuffle of the columns of  $K(m)$ . In [6], it has been shown that  $(T, K(m))$  is a matrix-tableau pair only if the word of  $T$  is an element of the Knuth class of  $K(m)$ . Therefore, the following problem arises:

*Given a tableau-pair  $(T, K(m))$  such that the word of  $T$  is a shuffle of the columns of the key  $K(m)$ , is  $(T, K(m))$  a matrix-tableau pair?*

The next algorithm and Theorem 4.4 give an answer to this problem. To this end, we need the following definition and lemma.

**Definition 4.2.** Given  $\sigma \in \mathfrak{S}_n$  and  $x \geq y$  in  $[n]$ , we define the  $n$  by  $n$  matrix  $S(x, y, \sigma) = E_{uv}$  if  $\sigma(u) = x > \sigma(v) = y$ , and the zero matrix otherwise.

Clearly,  $I + S(x, y, \sigma) = T_{uv}(1)$  with  $\sigma(u) = x$  and  $\sigma(v) = y$ . In particular,  $I + S(x, x, \sigma) = I$ .

**Lemma 4.2.** *Given  $\sigma \in \mathfrak{S}_n$  and  $x \geq y$  in  $[n]$ , we have, for any partition  $a$  of length  $\leq n$ ,*

$$(4.1) \quad \text{diag}_p(a)P_\sigma(I + S(x, y, \sigma))(I - S(x, y, \sigma)^T) \sim_L \text{diag}_p(a)P_{(xy)\sigma}.$$

*Proof.* Let  $u, v \in [n]$  be such that  $\sigma(u) = x$  and  $\sigma(v) = y$ . Then,  $(I + S(x, y, \sigma))(I - S(x, y, \sigma)^T) = T_{uv}(1)T_{vu}(-1)$ . By Lemma 4.1 (viii), the left hand side of (4.1) can

be written as

$$(4.2) \quad \text{diag}_p(a)P_\sigma T_{uv}(1)T_{vu}(-1) = \text{diag}_p(a)P_\sigma T_{vv}(-1)T_{vu}(-1)P_{(uv)}.$$

Since  $x \geq y$ , we have

$$(4.2) \sim_L \text{diag}_p(a_{\sigma(1)}, \dots, a_{\sigma(n)})P_{(uv)} \sim_L \text{diag}_p(a)P_\sigma P_{(uv)} = \text{diag}_p(a)P_{(xy)\sigma}.$$

□

**Main Algorithm.** Let  $\Pi$  be the biword (3.4). Our algorithm is presented as a three-step definition:

*Step 1.* For each  $k = t, \dots, 1$ , let  $X_k \subseteq \mathbb{N}^2$  be the set of initial vertices of the  $\ell_k$  lattice paths of length  $k$  of  $\Pi$ , and define

$$s(X_k) := \{s(x, j) : (x, j) \in X_k\} = \{s_{\ell_{t+1}+\dots+\ell_{k+1}+1}^0 < \dots < s_{\ell_{t+1}+\dots+\ell_{k+1}+\ell_k}^0\} \subseteq [n],$$

where we set  $\ell_{t+1} := 0$ . Let  $\sigma_1 \in \mathfrak{S}_n$  such that  $\sigma_1(i) = s_i^0$ , for  $i \in [\ell_1 + \dots + \ell_t]$ .

*Step 2.* For  $k = 1, \dots, t-1$ , let

$$J'_k := \{x \in J_k : (x, k) \text{ is positively-linked}\} = \{x_1^k < \dots < x_{q_k}^k\} \subseteq J_k$$

and  $\nu_0^k := id \in \mathfrak{S}_n$ . For each  $k = 1, \dots, t-1$  and  $j = 1, \dots, q_k$ , let  $(y_j^k, k_j)$  be the vertex negatively-linked to  $(x_j^k, k)$ , and define inductively

$$S_{x_j^k y_j^k}^{k+1} := S(x_j^k, s(y_j^k, k_j), \nu_{j-1}^k \sigma_k), \text{ and } \nu_j^k := (x_j^k s(y_j^k, k_j)) \nu_{j-1}^k.$$

Define  $\theta_{k+1} := \nu_{q_k}^k$ ,  $\sigma_{k+1} := \theta_{k+1} \sigma_k$ , and

$$S_{k+1} := \prod_{j=1}^{q_k} (I + S_{x_j^k y_j^k}^{k+1}) (I - S_{x_j^k y_j^k}^{k+1})^T.$$

*Step 3.* Let  $T$  be a tableau with biword  $\Pi$ , and let  $c/a$  be its skew-shape. Let  $A_0 := \text{diag}_p(a')$ , put  $B_k := S_k D_{[m_k]}$ , for  $k = 1, \dots, t$ , with  $S_1 := P_{\sigma_1}$ , and define inductively

$$A_k := A_{k-1} B_k.$$

*Remark 4.1.* (a) The Main Algorithm only depends on the biword  $\Pi$ .

(b) Lemma 3.1 asserts that the permutation  $\sigma_1 \in \mathfrak{S}_t$ , given in Step 1 of the Main Algorithm, is well defined.

(c) The matrix  $S_{x_j^k y_j^k}^{k+1}$ , defined in Step 2 of the Main Algorithm, is the zero matrix whenever  $s(y_j^k, k_j) = x_j^k$ , that is, if  $x_j^k = y_j^k$ . Therefore, the sequence  $A_0, B_1, \dots, B_t$ , obtained through the application of the Main Algorithm may be simplified if we use a biword  $\Pi$  whose graphical representation has a reduced number of links between vertices in distinct rows.

(d) Fix a  $k \in \{1, \dots, t-1\}$ , set  $u_j := \sigma_k^{-1}(\nu_{j-1}^k)^{-1}(x_j^k)$ , and  $v_j := \sigma_k^{-1}(\nu_{j-1}^k)^{-1}(s(y_j^k, k_j))$ . Then  $I + S_{x_j^k, y_j^k}^{k+1} = T_{u_j v_j}(1)$ ,  $(I - S_{x_j^k, y_j^k}^{k+1})^T = T_{v_j u_j}(-1)$ , and

$$S_{k+1} = \prod_{j=1}^{q_k} T_{u_j v_j}(1) T_{v_j u_j}(-1).$$

(e) For  $i = 1, \dots, t$ ,  $|J_i| = m_i = \sum_{k=\sigma^{-1}(i)}^t \ell_k \geq 0$ .

(f) Set  $I_k := [\ell_{t+1} + \dots + \ell_{k+1} + 1, \ell_{t+1} + \dots + \ell_{k+1} + \ell_k]$ , for all  $k = 1, \dots, t$ . Then,  $\sigma_1(I_k) = s(X_k)$ , for  $1 \leq k \leq t$ . In particular,  $s(X_k) = \emptyset$  if and only if  $\ell_k = 0$ .

(g)  $\sigma_{k+1} = \theta_{k+1} \cdots \theta_2 \sigma_1$ .

We say that a partition  $a$  *dominates* a partition  $b$  if  $|a| = |b|$  and the Young diagram of  $a$  is obtained by lifting some boxes in the Young diagram of  $b$ .

**Proposition 4.3.** *Let  $(T, K(m))$  be a tableau-pair such that  $T$  rectifies to  $K(m)$  and the word of  $T$  is a shuffle of the columns of  $K(m)$ . Let  $\Pi$  as in (3.4) be a biword of  $T$ . Then the sequence  $A_0, B_1, \dots, B_t$  generated by the Main Algorithm is a matrix realization for  $T$ . In particular,  $B_1, \dots, B_t$  is a matrix realization of a Young tableau  $F$  of weight  $m$  whose shape dominates  $\alpha(m)$ .*

Given a tableau  $T$  with rectification a key-tableau and whose word is a shuffle of its columns, there are, in general, several shuffle decompositions for the word of  $T$ , each corresponding to a biword  $\Pi$ . In the next section, we show that it is always possible to choose a biword  $\Pi$  where the shuffling decomposition satisfies additional properties so that the sequence  $B_1, \dots, B_t$  generated in the Main Algorithm is also a matrix realization of  $K(m)$ .

**Main Theorem 4.4.** *Let  $(T, K(m))$  be a tableau-pair such that the word of  $T$  is a shuffle of the columns of  $K(m)$ . Then there exists a biword  $\Pi$  of  $T$  such that the sequence  $A_0, B_1, \dots, B_t$  generated by the Main Algorithm is a matrix realization for the pair  $(T, K(m))$ .*

As already pointed out, it has been proved that  $(T, K)$  is a matrix tableau-pair only if the word of  $T$  is in the plactic class of  $K$  [6]. Thus, when the plactic class of  $K$  is the set of all shuffles of its columns, we obtain the following characterization of matrix tableau-pairs  $(T, K)$ .

**Corollary 4.5.** *Let  $T$  be a (skew) Young tableau with rectification  $P$ . Let  $K$  be a key whose plactic class is equal to the set of all shuffles of its columns. Then,  $(T, K)$  is a matrix tableau-pair if and only if  $P = K$ .*

Before giving the proofs of the statements above, we work out some examples.

**Example 4.2.** Consider the biword  $\Pi = \begin{pmatrix} 1144 & 23 \\ 4321 & 32 \end{pmatrix}$  as in (3.5) (see Example 3.2).

The word 433221 is in  $\sqcup \sqcup (v_4^1, v_3^0, v_2^1, v_1^0)$ . Let  $T$  be a tableau with biword  $\Pi$ . Then  $T$  has word 433221, indexing sets  $J_1 = \{4\}$ ,  $J_2 = \{3, 4\}$ ,  $J_3 = \{1, 2\}$ ,  $J_4 = \{1\}$ , and weight  $(1, 2, 2, 1)$ . The initial vertices of the lattice paths  $v_4$  and  $v_2$  are  $(4, 1)$  and  $(3, 2)$ , respectively. Recalling Example 3.3, we have  $s_1^0 = s(4, 1) = 4$  and  $s_2^0 =$

$s(3, 2) = 3$ , and thus  $s(X_4) = \{4\}$  and  $s(X_2) = \{3\}$ . Then, by Step 1 of the Main Algorithm, we must consider  $\sigma_1 \in \mathfrak{S}_4$  satisfying  $\sigma_1(1) = 4$  and  $\sigma_1(2) = 3$ . Take, for instance,  $\sigma_1 = (14)(23)$ . Next, define

$$\begin{aligned} I + S_{4,4}^2 &= I + S(4, 4, \sigma_1) = I, \quad \text{and} \quad \sigma_2 = (44)\sigma_1 = \sigma_1; \\ I + S_{3,2}^3 &= I + S(3, 2, \sigma_2) = T_{23}(1); \\ I + S_{4,1}^3 &= I + S(4, 1, (3\ 2)\sigma_2) = T_{14}(1), \quad \text{and} \quad \sigma_3 = (41)(32)\sigma_2; \\ I + S_{1,1}^4 &= I + S(1, 1, \sigma_3) = I, \quad \text{and} \quad \sigma_4 = (11)\sigma_3 = \sigma_3. \end{aligned}$$

Finally, let  $a' = (2, 2, 1, 0) + (2, 1, 1, 0)$  and define the matrices  $A_0 = \text{diag}_p(a')$ ,  $B_1 = P_{\sigma_1}D_{[1]}$ ,  $B_2 = D_{[2]}$ ,  $B_3 = T_{23}(1)T_{32}(-1)T_{14}(1)T_{41}(-1)D_{[2]}$ , and  $B_4 = D_{[1]}$ .

Clearly,  $A_1 = A_0B_1 \sim \text{diag}_p(a' + \chi^{J_1})$  and  $A_2 = A_0B_1B_2 \sim \text{diag}_p(a' + \chi^{J_1} + \chi^{J_2})$ . Since  $\sigma_3([2]) = J_3$ , by Lemma 4.2, we find that

$$\begin{aligned} A_3 &\sim_L \text{diag}_p(a' + \chi^{J_1} + \chi^{J_2})P_{\sigma_2}B_3 \sim_L \text{diag}_p(a' + \chi^{J_1} + \chi^{J_2})P_{\sigma_3}D_{[2]} \\ &\sim_R \text{diag}_p(a' + \chi^{J_1} + \chi^{J_2} + \chi^{J_3}). \end{aligned}$$

Since  $A_4 \sim \text{diag}_p(a' + \chi^{J_1} + \dots + \chi^{J_4})$ , the sequence  $A_0, B_1, \dots, B_4$  satisfies conditions I and II of Definition 4.1. It remains to show that this sequence also defines the key  $K = 4321\ 32$ . Bearing in mind Lemma 4.1, we may write

$$\begin{aligned} B_1B_2B_3B_4 &= P_{\sigma_1}D_{[1]}D_{[2]}T_{23}(1)T_{14}(1)T_{32}(-1)T_{41}(-1)D_{[2]}D_{[1]} \\ &= P_{\sigma_1}T_{23}(p)T_{14}(p^2)D_{[1]}D_{[2]}T_{32}(-1)T_{41}(-1)D_{[2]}D_{[1]} \\ &\sim_L D_{[1]}D_{[2]}T_{32}(-1)T_{41}(-1)D_{[2]}D_{[1]} \\ &\sim_R D_{[1]}D_{[2]}D_{[2]}D_{[1]}. \end{aligned}$$

Therefore, the sequence  $A_0, B_1, B_2, B_3, B_4$  is a matrix realization of the pair

$$(T = \begin{array}{cccc} & & & \boxed{1} \\ & & & \boxed{2} \\ & & \boxed{2} & \\ & \boxed{3} & & \\ \boxed{3} & & & \\ \boxed{4} & & & \end{array}, K = \begin{array}{cc} \boxed{1} & \boxed{2} \\ \boxed{2} & \boxed{3} \\ \boxed{3} & \\ \boxed{4} & \end{array} ).$$

In the next examples, we show that different shuffle decompositions may give different matrix realizations of the same tableau  $T$  described in the previous example. Moreover, there are shuffle decompositions that do not produce the tableau-pair  $(T, K)$  in Theorem 4.4. More precisely, we shall see that some shuffle decompositions may lead in our matrix construction of  $B_1B_2 \cdots B_t$  to condition (viii) in Lemma 4.1, and, therefore, we are driven away from the key-tableau  $K$ .

**Example 4.3.** Consider the biword  $\Pi^\dagger = \begin{pmatrix} 1244 & 13 \\ 4321 & 32 \end{pmatrix}$  as in (3.5), which is different from the one considered in Example 4.2. Note that  $s_1^0 = s(4, 1) = 4$  and  $s_2^0 = s(3, 2) =$

3. Following the Main Algorithm, we may define  $\sigma_1 = (14)(23) \in \mathfrak{S}_4$ , the matrices

$$\begin{aligned} I + S_{4,4}^2 &= I + S(4, 4, \sigma_1) = I, \\ I + S_{3,1}^3 &= I + S(3, 1, \sigma_1) = T_{24}(1), \\ I + S_{4,2}^3 &= I + S(4, 2, (31)\sigma_1) = T_{13}(1), \\ I + S_{2,1}^4 &= I + S(2, 1, (42)(31)\sigma_1) = T_{12}(1), \end{aligned}$$

and the permutations  $\sigma_2 = \sigma_1$ ,  $\sigma_3 = (42)(31)\sigma_1$  and  $\sigma_4 = (21)\sigma_3$ . Finally, define  $A_0 = \text{diag}_p(a')$ ,  $B_1 = P_{\sigma_1}D_{[1]}$ ,  $B_2 = D_{[2]}$ ,  $B_3 = T_{24}(1)T_{42}(-1)T_{13}(1)T_{31}(-1)D_{[2]}$ , and  $B_4 = T_{12}(1)T_{21}(-1)D_{[1]}$ . Applying the same arguments as the ones used in Example 4.2, we may show that the sequence  $A_0, B_1, B_2, B_3, B_4$  is a matrix realization for  $(T, K)$ , and it is clearly different from the one obtained in Example 4.2.

**Example 4.4.** Consider now the biword  $\Pi'$  as in (3.5) (see Example 3.2). We shall see that, in this case, the sequence of matrices obtained by the Main Algorithm is not a matrix realization of  $(T, K)$ . Following Step 1 of the Main Algorithm, we may define  $\sigma_1 = (14)(23)$ , since  $s_1^0 = s(4, 1) = 4$  and  $s_2^0 = s(4, 2) = 3$ . By Step 2, we define the matrices

$$\begin{aligned} I + S_{4,3}^2 &= I + S(4, 3, \sigma_1) = T_{12}(1), \\ I + S_{3,2}^3 &= I + S(3, 2, (43)\sigma_1) = T_{13}(1), \\ I + S_{4,1}^3 &= I + S(4, 1, (32)(43)\sigma_1) = T_{24}(1), \\ I + S_{2,1}^4 &= I + S(2, 1, (41)(32)(43)\sigma_1) = T_{12}(1), \end{aligned}$$

and the permutations  $\sigma_2 = (43)\sigma_1$ ,  $\sigma_3 = (41)(32)\sigma_2$  and  $\sigma_4 = (21)\sigma_3$ . Finally, we define the matrices  $A_0 = \text{diag}_p(a')$ ,  $B_1 = P_{\sigma_1}D_{[1]}$ ,  $B_2 = T_{12}(1)T_{21}(-1)D_{[2]}$ ,  $B_3 = T_{13}(1)T_{31}(-1)T_{24}(1)T_{42}(-1)D_{[2]}$ , and  $B_4 = T_{12}(1)T_{21}(-1)D_{[1]}$ .

As in Example 4.2, using Lemma 4.2, we find that  $A_0, B_1, B_2, B_3, B_4$  satisfy conditions I and II of Definition 4.1. Let us now compute the invariant partition of  $B_1B_2B_3B_4$ . Using Lemma 4.1, (ii), (v), and (vii), we may write

$$\begin{aligned} (4.3) \quad B_1B_2B_3B_4 &\sim_L D_{\{1\}}T_{21}(-1)D_{\{1,2\}}T_{31}(-1)T_{42}(-1)D_{\{1,2\}}T_{12}(1)T_{21}(-1)D_{\{1\}} \\ &= D_{\{1\}}T_{21}(-1)T_{12}(1)D_{\{1,2\}}T_{31}(-1)T_{32}(-1)T_{42}(-1)D_{\{1,2\}}T_{21}(-1)D_{\{1\}}, \end{aligned}$$

and by Lemma 4.1, (vi), we find that

$$(4.3) \sim_R D_{\{1\}}T_{21}(-1)T_{12}(1)D_{\{1,2\}}D_{\{1,2\}}D_{\{1\}}.$$

Since by Lemma 4.1, (viii),  $T_{21}(-1)T_{12}(1) = T_{22}(-1)T_{12}(1)P_{(12)}$ , we obtain

$$\begin{aligned} B_1B_2B_3B_4 &\sim D_{\{1\}}T_{22}(-1)T_{12}(1)P_{(12)}D_{\{1,2\}}D_{\{1,2\}}D_{\{1\}} \\ &\sim_L D_{\{1\}}P_{(12)}D_{\{1,2\}}D_{\{1,2\}}D_{\{1\}} \\ &\sim_R D_{\{1\}}D_{\{1,2\}}D_{\{1,2\}}D_{\{2\}}. \end{aligned}$$

$$F = \begin{array}{|c|c|} \hline 1 & 2 \\ \hline 2 & 3 \\ \hline 3 & 4 \\ \hline \end{array} \neq K = \begin{array}{|c|c|} \hline 1 & 2 \\ \hline 2 & 3 \\ \hline 3 & \\ \hline 4 & \\ \hline \end{array}$$

The sequence  $B_1, B_2, B_3, B_4$  gives the Young tableau  $F \neq K$  and, therefore,  $A_0, B_1, B_2, B_3, B_4$  is a matrix realization for the pair  $(T, F)$  with  $F \neq K$ .

## 5. PROOF OF THE MAIN THEOREM

For any biword  $\Pi$  as in (3.4), the sequence of matrices  $A_0, B_1, \dots, B_t$  generated by the Main Algorithm always defines a tableau-pair  $(T, F)$ , where  $T$  is a Young tableau with biword  $\Pi$  which rectifies to a key-tableau  $K(m)$ , and  $F$  is some Young tableau with weight  $m$  whose shape dominates  $\alpha(m)$ . To prove that  $F$  is the rectification of  $T$ , we need restrictions on the biword  $\Pi$ , as we have seen in Example 4.4. The proof of the main theorem, Theorem 4.4, has therefore two parts. Firstly, we show that the matrix sequence  $(A_0, B_1, \dots, B_t)$  generated in the Main Algorithm, by a biword  $\Pi$  as in (3.4), defines any Young tableau  $T$  with that biword. Secondly, we show that we may always select an appropriate biword of  $T$  such that the matrix sequence  $B_1, \dots, B_t$  defines its rectification.

**5.1. Biwords generate matrix-tableaux.** In this section, we only assume that we have applied the Main Algorithm to a biword  $\Pi$ , without any additional conditions.

**Lemma 5.1.** *Let  $t \geq 1$  and  $k \in \{1, \dots, t\}$ . Let  $(x, k)$  and  $(y, k + \varepsilon)$ ,  $\varepsilon \geq 1$ , be two consecutive vertices of a lattice path of  $\Pi$ . Then  $\theta_{k+\varepsilon} \cdots \theta_{k+1}(x) = y$ .*

*Proof.* Case 1. Assume  $s(y, k + \varepsilon) = y$ . This means that there are no vertices in positions  $(y, j)$ ,  $k < j < k + \varepsilon$ . Thus, following Step 2 of the Main Algorithm, we may write

$$\theta_{k+\varepsilon} \cdots \theta_{k+1} = \rho_2(x y) \rho_1,$$

for some permutations  $\rho_1, \rho_2 \in \mathfrak{S}_n$ , such that  $\rho_1(x) = x$ .

We claim that  $\rho_2(y) = y$ . Note that, for this equality to be false, we should have a vertex  $(y, j)$ , with  $j > k + \varepsilon$ , negatively-linked to a vertex  $(a, i)$ , for some  $i \in [k, k + \varepsilon - 1]$ , with  $a > x$ , and  $(y, k + \varepsilon)$  should be an obstacle. This is not possible since the sets of column vertices of the lattice paths containing  $(y, k + \varepsilon)$  and  $(a, i)$ , respectively, must be comparable for the inclusion order. Therefore,  $\rho_2(y) = y$  and the result follows.

Case 2. Assume now that  $s(y, k + \varepsilon) \neq y$ , and let  $(y, k + \varepsilon), (a_1, b_1), \dots, (a_r, b_r)$  be the  $(x, k)$ -path (see Definition 3.1). Clearly,  $s(y, k + \varepsilon) = a_r$ , the integer  $r$  is even, and, for each odd integer  $i \in [r]$ , the vertices  $(a_i, b_i)$  and  $(a_{i+1}, b_{i+1})$  are consecutive vertices of a lattice path, with  $a_{i-1} = a_i$  and  $a_0 = y$ . Without loss of generality, assume that  $b_1 \leq b_3 \leq \dots \leq b_{r-1}$ . Then, following Step 2 of the Main Algorithm, we have

$$\theta_{k+\varepsilon} \cdots \theta_{k+1} = \rho_1(a_1 a_2) \rho_2(a_2 a_3) \cdots \rho_{r-2}(a_{r-2} a_{r-1}) \rho_{r-1}(a_{r-1} a_r) \rho_r(x a_r) \rho_0,$$

for some  $\rho_i \in \mathfrak{S}_n$ , such that  $\rho_0(x) = x$ . By an argument similar to the one used in Case 1, we find that  $\rho_i(a_i) = a_i$ , for  $i = 1, \dots, r$ , and the result follows.  $\square$

**Lemma 5.2.** *Let  $t \geq 1$  and  $1 \leq j \leq t$ . Let  $(x, j)$  be the initial vertex of a lattice path. Then,*

- (a)  $\sigma_j(\sigma_1^{-1}(s(x, j))) = x$ .
- (b)  $\sigma_j[m_j] = J_j$ .

*Proof.* (a) Note that  $\sigma_j(\sigma_1^{-1}(s(x, j))) = \theta_j \cdots \theta_2(s(x, j))$ . If there are no vertices in row  $x$  to the left of column  $j$ , then  $s(x, j) = x$ ,  $\theta_j \cdots \theta_2(x) = x$ , and thus  $\sigma_j(\sigma_1^{-1}(s(x, j))) = x$ .

Suppose there is a vertex in row  $x$  to the left of column  $j$ . Extend the biword  $\Pi$  to a biword  $\Pi_0$  by adding extra biletters  $\binom{n+u}{0}$  for all  $u \in [\ell_1 + \dots + \ell_t]$ , such that  $\binom{x}{j}$  and  $\binom{n+1}{0}$  are consecutive vertices in  $\Pi_0$ , and the initial vertex of each lattice path is linked to a distinct vertex  $(n+u, 0)$ . In particular, the vertices  $(x, j)$  and  $(n+1, 0)$  are linked. Put  $\theta_1 := \nu(n+1 \ s(x, j))$ , where  $\nu$  is defined as in Step 2 of the Main Algorithm. By Lemma 5.1,  $\theta_j \cdots \theta_2 \theta_1(n+1) = x$ , that is,  $\theta_j \cdots \theta_2(s(x, j)) = x$ , and the result follows. (As the lattice diagram of  $\Pi_0$  is contained in  $([n] \cup [\ell_1 + \dots + \ell_t]) \times (\{0\} \cup [t])$ , in Algorithm 3.1, we have to make a shift and put  $y := -1$  in the case of an initial vertex.)

(b) Let  $k \in \{1, \dots, t\}$ . Recall that  $I_k := [\ell_{t+1} + \dots + \ell_{k+1} + 1, \ell_{t+1} + \dots + \ell_{k+1} + \ell_k]$  and  $|J_j| = m_j = \sum_{k=\sigma^{-1}(j)}^t \ell_k$ . Thus,  $I_k \subseteq [m_j]$  for  $k \in \{t, \dots, \sigma^{-1}(j)\}$ . We use induction on  $j$  to prove that, if  $x \in J_j$  is such that  $(x, j)$  is a vertex of a lattice path of length  $k$  with  $k \in \{t, \dots, \sigma^{-1}(j)\}$ , there exists  $i_x \in I_k$  such that  $\sigma_j(i_x) = x$ .

When  $j = 1$ , the letters in  $J_1$  correspond to initial vertices of the lattice paths of length  $k$ , for  $k = t, \dots, \varepsilon$ , where  $\varepsilon := \sigma^{-1}(1)$ . Therefore, we may write

$$\{(x, 1) : x \in J_1\} = \bigcup_{k=t}^{\varepsilon} X_k.$$

By the definition of  $\sigma_1$ , we have  $\sigma_1(I_k) = s(X_k)$ . Since  $s(X_k) = \{x : (x, 1) \in X_k\}$ , for  $k = t, \dots, \varepsilon$ , the result follows.

Fix now  $j$  in  $\{2, \dots, t\}$ , and let  $y \in J_j$  be the row index of a letter  $j$  belonging to a lattice path  $\mathbf{v}$  of length  $k$ . Clearly, we must have  $k \in \{t, \dots, \sigma^{-1}(j)\}$ . If  $(y, j)$  is the initial vertex of  $\mathbf{v}$  then, by (a), we have  $\sigma_j(\sigma_1^{-1}(s(y, j))) = y$ , with  $\sigma_1^{-1}(s(y, j)) \in I_k$ .

Assume now that  $(y, j)$  is negatively-linked to a vertex  $(x, j - \varepsilon)$  with  $x \geq y$ , and  $\varepsilon \geq 1$ . By induction, there exists  $i_x \in I_k$  such that  $\sigma_{j-\varepsilon}(i_x) = x$ , and by Lemma 5.1 we find that  $\theta_j \cdots \theta_{j-\varepsilon+1}(x) = y$ . Thus,

$$\sigma_j(i_x) = \theta_j \cdots \theta_{j-\varepsilon+1} \sigma_{j-\varepsilon}(i_x) = y.$$

By induction, our claim is proved. Thus, for each  $x \in J_j$  there is  $i_x \in [m_j]$  such that  $\sigma_j(i_x) = x$ . Since  $|J_j| = m_j$ , we must have  $\sigma_j([m_j]) = J_j$ .  $\square$

**Corollary 5.3.** *If  $(x, j)$  is a vertex of a lattice path of length  $k$  with initial vertex  $(y, j - \varepsilon)$ ,  $\varepsilon \geq 0$ , then  $\sigma_j^{-1}(x) = \sigma_{j-\varepsilon}^{-1}(y) = \sigma_1^{-1}(s(y, j - \varepsilon)) \in I_k$ .*

*Proof.* This follows from the proof of the previous lemma.  $\square$

*Proof of Proposition 4.3.* Let  $t \geq 1$ . We have to show that, for each  $k \in \{0, 1, \dots, t\}$ , the matrix  $A_k$ , given by Step 3 of the Main Algorithm, is left equivalent to  $\text{diag}_p(a' + \chi^{J_0} + \chi^{J_1} \cdots + \chi^{J_k})P_{\sigma_k}$ , where  $\sigma_0 := \text{id} \in \mathfrak{S}_n$  and  $J_0 := \emptyset$ . The proof is by induction on  $k$ . For  $k = 0$ ,  $P_{\text{id}} = I$  and  $A_0 := \text{diag}_p(a')$ , and in this case there is nothing to prove. So let  $k$  be in  $\{1, \dots, t\}$ . By induction,  $A_{k-1}$  is left equivalent to  $\text{diag}_p(a' + \chi^{J_1} \cdots + \chi^{J_{k-1}})P_{\sigma_{k-1}}$ . Therefore, by definition of  $B_k$ ,  $A_k$  is left equivalent to

$$(5.1) \quad \text{diag}_p(a' + \chi^{J_1} \cdots + \chi^{J_{k-1}})P_{\sigma_{k-1}}S_kD_{[m_k]}.$$

Recall that  $S_1 = P_{\sigma_1}$  and  $S_k = \prod_{i=1}^{q_k-1} (I + S_{x_i^k y_i^k}^k) (I - S_{x_i^k y_i^k}^k)^T$ , for  $k > 1$ . Therefore, by Lemma 4.2 and (5.1),  $A_k$  is left equivalent to

$$(5.2) \quad \text{diag}_p(a' + \chi^{J_1} \cdots + \chi^{J_{k-1}}) P_{\sigma_k} D_{[m_k]} = \text{diag}_p(a' + \chi^{J_1} \cdots + \chi^{J_{k-1}}) D_{\sigma_k[m_k]} P_{\sigma_k}.$$

By Lemma 5.2 (b), we have  $\sigma_k[m_k] = J_k$ , and the result follows.  $\square$

**5.2. Biwords again: a canonical biword.** In the proof of Proposition 4.3 there are no restrictions on the shuffle decomposition that we have considered in  $\Pi$ . The sequence of matrices  $A_0, B_1, \dots, B_t$  given by the Main Algorithm satisfies always conditions I and II of Definition 4.1 for any biword  $\Pi$  as in (3.4) of  $T$ . To prove that  $B_1, \dots, B_t$  defines the rectification of  $T$ , we need restrictions on the biword  $\Pi$ , as we have seen in Example 4.4. We shall show that we may always choose a biword equivalent to  $\Pi$  obeying conditions that are convenient for our purpose.

In a biword  $\Pi$  some vertices play a special role in the proof of our Theorem 4.4. They are identified in the following definitions.

**Definition 5.1.** Given distinct lattice paths  $\mathbf{u}, \mathbf{v}$  of  $\Pi$ , a vertex  $(a, b) \in \mathbf{u}$  is said to be a *right critical vertex* of the pair  $(\mathbf{u}, \mathbf{v})$  if it is positively-linked to a vertex  $(a', b')$  such that  $(s(a', b'), b)$  belongs to  $\mathbf{v}$ . If  $b' = b + 1$  the right critical vertex is of type I, otherwise it is of type II.

Given pairwise distinct lattice paths  $\mathbf{v}_1, \dots, \mathbf{v}_k$ ,  $k \geq 2$ , of  $\Pi$ , a vertex  $(a_1, b_1) \in \mathbf{v}_1$  is said to be a right critical vertex of type I (respectively of type II) of  $(\mathbf{v}_1, \dots, \mathbf{v}_k)$  with *critical components*  $\mathbf{v}_2, \dots, \mathbf{v}_{k-1}$ , if

- $(a_1, b_1) \in \mathbf{v}_1$  is a right critical vertex of type I (respectively of type II) of  $(\mathbf{v}_1, \mathbf{v}_2)$ , positively-linked to  $(a_2, b_2)$ ;
- for  $i = 2, \dots, k-2$ ,  $(s(a_i, b_i), b_1)$  is a right critical vertex of  $(\mathbf{v}_i, \mathbf{v}_{i+1})$  positively-linked to  $(a_{i+1}, b_{i+1})$ , and
- $(s(a_{k-1}, b_{k-1}), b_1)$  is a right critical vertex of  $(\mathbf{v}_{k-1}, \mathbf{v}_k)$ .

Since the lattice paths are distinct, a right critical vertex of type I is always an obstacle. In the case of a right critical vertex of type II,  $b' - b > 1$  and we may have  $a' = a$  with an obstacle of another lattice path between  $(a, b)$  and  $(a, b')$ . The following diagrams are schematic representations of right critical vertices of types I and II, respectively, both with one critical component,



**Definition 5.2.** Given distinct lattice paths  $\mathbf{u}, \mathbf{v}$  of  $\Pi$ , a vertex  $(a, b) \in \mathbf{u}$  is said to be a *left critical vertex* of  $(\mathbf{u}, \mathbf{v})$  if it is not negatively-linked to a vertex in column  $b - 1$  and there is an obstacle on its left in position  $(a, b - 1)$  belonging to  $\mathbf{v}$ .

Given pairwise distinct lattice paths  $\mathbf{v}_1, \dots, \mathbf{v}_k$ ,  $k \geq 2$ , of  $\Pi$ , a vertex  $(a_1, b) \in \mathbf{v}_1$  is a left critical vertex of  $(\mathbf{v}_1, \dots, \mathbf{v}_k)$  with *critical components*  $\mathbf{v}_2, \dots, \mathbf{v}_{k-1}$  if

- $(a_1, b) \in \mathbf{v}_1$  is a left critical vertex of  $(\mathbf{v}_1, \mathbf{v}_2)$ ;
- for  $i = 1, \dots, k - 2$ , the obstacle  $(a_i, b - 1) \in \mathbf{v}_{i+1}$  is a right critical vertex of type I of  $(\mathbf{v}_{i+1}, \mathbf{v}_{i+2})$  positively-linked to a vertex  $(a_{i+1}, b)$ , and
- $(a_{k-1}, b - 1)$  is an obstacle of  $(\mathbf{v}_{k-1}, \mathbf{v}_k)$ .

Schematically, with  $k = 2$ , either  $(a_1, b)$  is the initial vertex of  $\mathbf{v}_1$  or it is negatively linked to a vertex not in column  $b - 1$ .



For instance, in the biword  $\Pi'$  as in (3.5) the vertices  $(4, 2)$  and  $(2, 3)$  are a left critical vertex and a right critical vertex of type I of the two lattice paths, respectively. They are the only critical vertices in this biword. Consider now the biword represented in Example 3.4. The vertex  $(4, 3)$  is a left critical vertex of  $\left(\begin{smallmatrix} 345 \\ 531 \end{smallmatrix}, \begin{smallmatrix} 22233 \\ 54321 \end{smallmatrix}\right)$  with critical component  $\begin{smallmatrix} 1344 \\ 5321 \end{smallmatrix}$ , and clearly it is also a left critical vertex of  $\left(\begin{smallmatrix} 345 \\ 531 \end{smallmatrix}, \begin{smallmatrix} 1344 \\ 5321 \end{smallmatrix}\right)$ . In this biword, there is no other left critical vertex. Notice that  $(4, 3)$  is also a right critical vertex of type II of  $\left(\begin{smallmatrix} 345 \\ 531 \end{smallmatrix}, \begin{smallmatrix} 1344 \\ 5321 \end{smallmatrix}\right)$ . This example shows that a vertex may be simultaneously a left and a right critical vertex. The vertex  $(5, 1)$  is a right critical vertex of type II of  $\left(\begin{smallmatrix} 345 \\ 531 \end{smallmatrix}, \begin{smallmatrix} 2 \\ 1 \end{smallmatrix}\right)$ , and the vertex  $(4, 2)$  is a right critical vertex of type I of  $\left(\begin{smallmatrix} 1344 \\ 5321 \end{smallmatrix}, \begin{smallmatrix} 22233 \\ 54321 \end{smallmatrix}\right)$ .

In what follows, given a biword  $\Pi$  as in (3.4), we provide a procedure to adjust the links between their vertices, to form a new biword  $\Pi'$  equivalent to  $\Pi$ , where the new lattice path decomposition satisfies some properties needed in the sequel.

**Algorithm 5.1.** Let  $\Pi$  be the biword (3.4) identified with its lattice path representation in  $[n] \times [t]$ . For each vertex  $(a, b)$  of  $\Pi$ , we move along each column, from top to bottom, and left to right, to check whether  $(a, b)$  is a left or a right critical vertex of type I of a pair of lattice paths  $(\mathbf{u}, \mathbf{v})$  such that both have vertices in columns  $b, b + 1, \dots, b + r$ , with  $r \geq 1$ ,  $\mathbf{u}$  is below  $\mathbf{v}$  in columns  $b, \dots, b + r - 1$ , and  $\mathbf{u}$  is above  $\mathbf{v}$  in column  $b + r$ . If so, we have two situations:

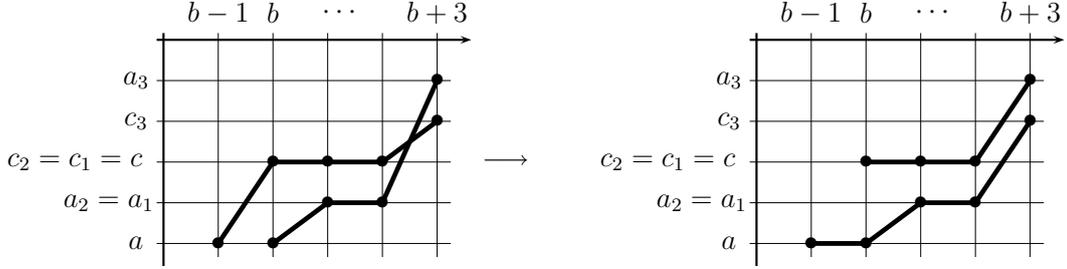
1-  $(a, b)$  is a *left critical vertex* of  $(\mathbf{u}, \mathbf{v})$ , and then  $(a, b - 1)$  is an obstacle of  $\mathbf{v}$ . The lattice paths  $\mathbf{u}$  and  $\mathbf{v}$ , restricted to the vertices in columns  $[b, b + r]$  and  $[b - 1, b + r]$ , respectively, are the factors of the sub-biword

$$\begin{aligned} \Pi_L &= \mathbf{u}\mathbf{v} \\ &= \begin{pmatrix} a_r & a_{r-1} & \cdots & a_1 & a \\ b+r & b+r-1 & \cdots & b+1 & b \end{pmatrix} \begin{pmatrix} c_r & c_{r-1} & \cdots & c_1 & c & a \\ b+r & b+r-1 & \cdots & b+1 & b & b-1 \end{pmatrix} \end{aligned}$$

of  $\Pi$ , respectively, where  $a > c$ ,  $a_i > c_i$ , for  $i = 1, \dots, r - 1$ , and  $a_r < c_r$ . Then, relink the vertices of  $\mathbf{u}$  and  $\mathbf{v}$ , between columns  $b$  and  $b + r$ , so that the sub-biword  $\Pi_L$  is

replaced by

$$\Pi'_L = \begin{pmatrix} a_r & c_{r-1} & \cdots & c_1 & c \\ b+r & b+r-1 & \cdots & b+1 & b \end{pmatrix} \begin{pmatrix} c_r & a_{r-1} & \cdots & a_1 & a & a \\ b+r & b+r-1 & \cdots & b+1 & b & b-1 \end{pmatrix},$$

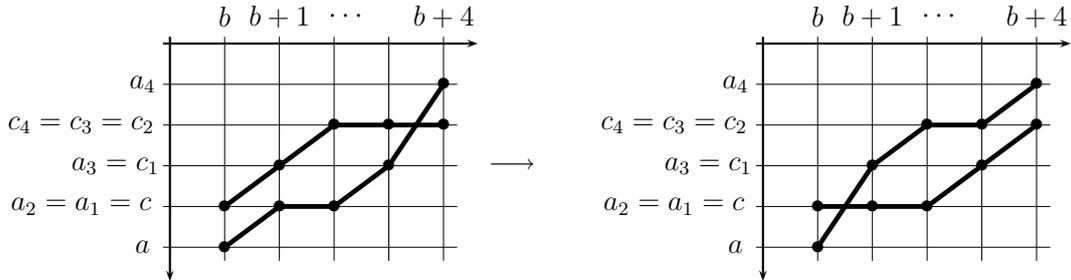


2-  $(a, b)$  is a *right critical vertex of type I* of  $(\mathbf{u}, \mathbf{v})$ . The lattice paths  $\mathbf{u}$  and  $\mathbf{v}$  restricted to the vertices in columns  $[b, b+r]$  are factors of the sub-biword

$$\Pi_R = \mathbf{u}\mathbf{v} = \begin{pmatrix} a_r & a_{r-1} & \cdots & a_1 & a \\ b+r & b+r-1 & \cdots & b+1 & b \end{pmatrix} \begin{pmatrix} c_r & c_{r-1} & \cdots & c_1 & c \\ b+r & b+r-1 & \cdots & b+1 & b \end{pmatrix}$$

of  $\Pi$ , where  $a > c = a_1$ ,  $a_i > c_i$ , for  $i = 2, \dots, r-1$ , and  $a_r < c_r$ . Then, relink the vertices of  $\mathbf{u}$  and  $\mathbf{v}$ , between columns  $b$  and  $b+r$ , so that the sub-biword  $\Pi_R$  is replaced by

$$\Pi'_R = \begin{pmatrix} a_r & c_{r-1} & \cdots & c_1 & a \\ b+r & b+r-1 & \cdots & b+1 & b \end{pmatrix} \begin{pmatrix} c_r & a_{r-1} & \cdots & a_1 & c \\ b+r & b+r-1 & \cdots & b+1 & b \end{pmatrix},$$



Let  $\mathbf{u}'$  and  $\mathbf{v}'$  be the lattice paths obtained from  $\mathbf{u}$  and  $\mathbf{v}$  replacing, in the first case, the factor  $\Pi_L$  by  $\Pi'_L$ , and, in the second case, the factor  $\Pi_R$  by  $\Pi'_R$ . In this way, we obtain a new biword  $\Pi'$  equivalent to  $\Pi$ , where the lattice paths  $\mathbf{u}$  and  $\mathbf{v}$  were replaced by the new lattice paths  $\mathbf{u}'$  and  $\mathbf{v}'$ , respectively. The vertex  $(a, b)$  belonging to the new lattice path  $\mathbf{u}'$  is no longer a critical vertex of  $(\mathbf{u}', \mathbf{v}')$ . Moreover, during this process no new critical points will occur in the vertices already checked. Before moving to the top of the next column, we continue this analysis, restarting with the same vertex  $(a, b)$  and moving down along column  $b$ .

As there is only a finite number of vertices, this algorithm produces a new biword  $\Pi'$ , equivalent to  $\Pi$ , under the conditions of (3.4), satisfying the following proposition.

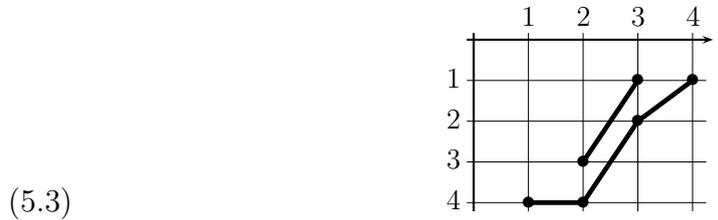
**Proposition 5.4.** *The biword  $\Pi'$  obtained by the application of Algorithm 5.1 to  $\Pi$  is such that, if  $\mathbf{u}$  and  $\mathbf{v}$  are two lattice paths each having vertices in all columns  $b, \dots, b+r$ , for some  $b, b+r \in [t]$ , with  $r \geq 1$ , and, if  $(a, b)$  is either a left critical vertex, or a right critical vertex of type I of  $(\mathbf{u}, \mathbf{v})$ , then the vertices of  $\mathbf{u}$  are below the vertices of  $\mathbf{v}$  in columns  $b, b+1, \dots, b+r$ .*

*Remark 5.1.* A biword  $\Pi$  satisfies trivially the conditions of the proposition above if either (1) two vertices in the same row belong to the same lattice path; or, (2) if  $(a, b')$  and  $(a, b)$ ,  $b' < b$ , are two vertices of  $\Pi$  in distinct lattice paths, then  $(a, b')$  is single or only negatively linked, that is,  $s(a, b) = a$ .

When  $T$  rectifies to  $K(\sigma\alpha(m))$ , with  $\sigma$  the identity, or the reverse permutation in  $\mathfrak{S}_t$ , or  $\sigma \in \mathfrak{S}_3$ , the conditions in the proposition above are satisfied [4, 5, 7].

**Example 5.1.** (1) The biword  $\Pi$  as in (3.5) used in the matrix construction of Example 4.2 satisfies Proposition 5.4.

(2) In Example 4.4 the biword  $\Pi'$  as in (3.5) was used in the matrix construction, but the sequence  $B_1, \dots, B_4$  generated in the Main Algorithm does not define  $K$ . Notice that both lattice paths  $\mathbf{u} = \begin{pmatrix} 14 \\ 32 \end{pmatrix}$  and  $\mathbf{v} = \begin{pmatrix} 1234 \\ 4321 \end{pmatrix}$  have vertices in columns 3 and 2; the vertex in position  $(4, 2)$  is a left critical vertex of  $(\mathbf{u}, \mathbf{v})$ , and  $\mathbf{u}$  is above  $\mathbf{v}$  in column 3. Therefore  $\Pi'$  does not satisfy Proposition 5.4. Applying the algorithm above, we get  $\Pi^\dagger$  as in (3.5), which already satisfies the conditions of Proposition 5.4,

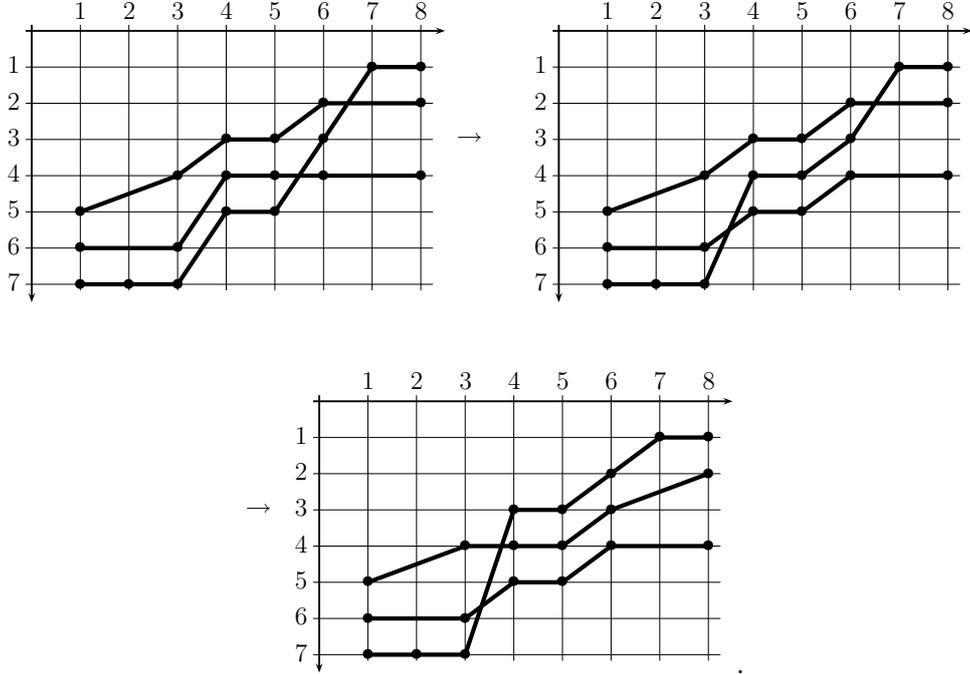


$\Pi^\dagger$  was used in the matrix construction of Example 4.3 where the sequence  $B_1, \dots, B_t$  produces  $K$ .

(3) Consider now the lattice path representation of the biword

$$\begin{pmatrix} 11355777 & 444466 & 223345 \\ 87654321 & 865431 & 865431 \end{pmatrix}.$$

This biword does not satisfy the conditions of Proposition 5.4, since, for example, the vertex  $(7, 3)$  is a right critical vertex of the lattice paths  $\mathbf{u} = \begin{pmatrix} 11355777 \\ 87654321 \end{pmatrix}$  and  $\mathbf{v} = \begin{pmatrix} 444466 \\ 865431 \end{pmatrix}$ . Both have vertices in columns  $3, 4, \dots, 7$ , and  $\mathbf{u}$  is above  $\mathbf{v}$  in column 7. Below are the successive steps of the application of Algorithm 5.1 to this biword. Notice that the lattice path representation of the final biword satisfies the conditions of Proposition 5.4,



In the following, we assume that the biword  $\Pi$  satisfies the conditions of Proposition 5.4. Recall that we have encoded the tableau  $T$  by a biword which we identify with its lattice path representation. This lattice path representation is used in the Main Algorithm to construct a sequence of matrices  $A_0, B_1, \dots, B_t$ , for which we have already proved that they define the tableau  $T$ . In order to show that  $A_0, B_1, \dots, B_t$  is a matrix realization for the pair  $(T, K)$ , with  $K$  the rectification of  $T$ , we must show that  $B_1, \dots, B_t$  defines the key  $K$ . The main pieces of our matrix construction, as observed in Remark 4.1 (d), are the elementary matrices  $T_{ij}(1)$  whose indices  $i, j$  depend on the lattice path representation of  $\Pi$ . The next lemmas characterize these elementary matrices by relating their indices  $i, j$  with the particular lattice path representation that we have chosen in Proposition 5.4. This allows us to use the commutation rules of Lemma 4.1, and it allows us to avoid any occurrence of the situation (viii) in Lemma 4.1, and, thereby, to show that the invariant partition of  $B_1, \dots, B_t$  is exactly the shape of the key-tableau  $K$ .

**Lemma 5.5.** *Let  $(x, k)$  and  $(y, k + \varepsilon)$ ,  $\varepsilon \geq 1$ , be two consecutive vertices of a lattice path of length  $q$  in  $\Pi$  with initial vertex  $(u, k')$ . Then the permutation  $\nu$  and the matrix*

$$S_{x,y}^{k+1} = S(x, s(y, k + \varepsilon), \nu\sigma_k)$$

*defined in Step 2 of the Main Algorithm satisfy:*

- (a)  $\nu(x) = x$ .
- (b)  $\nu^{-1}(s(y, k + \varepsilon)) \notin J'_k = \{a \in J_k : (a, k) \text{ is a positively-linked vertex}\}$ .
- (c)  $I + S_{x,y}^{k+1} = T_{ij}(1)$ , with  $i = \sigma_k^{-1}(x) = \sigma_1^{-1}(s(u, k')) \in I_q \subseteq [m_k]$ , and  $j \notin [J'_k]$ .

*Proof.* (a) and (b) are obvious. To prove (c), recall that by the definition of the matrix  $I + S_{x,y}^{k+1} = (s_{ij})$ , we have  $s_{ij} \neq 0$  only if  $\nu\sigma_k(i) = x$  and  $\nu\sigma_k(j) = s(y, k + \varepsilon)$ .

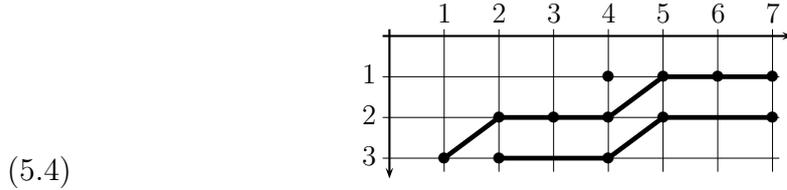
Using (a) and Corollary 5.3, we obtain

$$i = \sigma_k^{-1}(x) = \sigma_1^{-1}(s(u, k')) \in I_q \subseteq [m_k].$$

Now, by (b), we find that  $\sigma_k(j) = \nu^{-1}(s(y, k + \varepsilon)) \notin J'_k$ . If  $\sigma_k(j) \notin J_k$ , by Lemma 5.2 (b), we must have  $j \notin [m_k]$ . Assume now that  $\sigma_k(j) \in J_k \setminus J'_k$ . Then, we must have  $\sigma_k(j) = a$ , where  $(a, k)$  is a not positively-linked vertex of a lattice path of length  $q'$ , with  $q > q'$ . Thus  $j = \sigma_k^{-1}(a) \in I_q \subseteq [m_k] \setminus [|J'_k|]$ .  $\square$

**Corollary 5.6.** *Each matrix  $I + S_{x,y}^{k+1} = T_{ij}(1)$ , defined in Step 2 of the Main Algorithm, is upper triangular with  $i \in [m_k]$  and  $j > i$ . Moreover,  $j \in [m_k]$  if and only if there is a vertex  $(a, k)$ , not positively-linked, belonging to a lattice path  $\mathbf{v}$ , and such that  $(x, k)$  is a right critical vertex of type I or II of  $(\mathbf{u}, \mathbf{v})$ , where  $\mathbf{u}$  is the lattice path containing  $(x, k)$ .*

**Example 5.2.** Let  $\Pi = \begin{pmatrix} 1112223 & 2233 & 1 \\ 7654321 & 7542 & 4 \end{pmatrix}$  and its lattice path representation



Applying the Main Algorithm, we must set  $\sigma_1 = (13)$ , since  $s_1^0 = s(3, 1) = 3$ ,  $s_2^0 = s(3, 2) = 2$  and  $s_3^0 = s(1, 4) = 1$ , and thus  $\sigma_i = (32)\sigma_1$ , for  $i = 2, 3, 4$  and  $\sigma_5 = (32)(21)\sigma_4$ . Consider the matrices  $S_{3,2}^2 = S(3, s(2, 2), \sigma_2)$ ,  $S_{2,1}^5 = S(2, s(1, 5), \sigma_4)$ , and  $S_{3,2}^5 = S(2, s(2, 5), (21)\sigma_4)$ , produced in Step 2 of the Main Algorithm. Since  $s(2, 2) = 2$ ,  $s(1, 5) = 1$  and  $s(2, 5) = 2$ , we get

$$(5.5) \quad I + S_{3,2}^2 = T_{12}(1), \quad I + S_{2,1}^5 = T_{13}(1), \quad \text{and} \quad I + S_{3,2}^5 = T_{23}(1).$$

Each matrix in (5.5) is upper triangular and satisfies Lemma 5.5. Notice that in the last two matrices the column index  $3 \in [3] = [m_4]$ . This is consistent with the previous corollary, since the vertex  $(1, 4)$  is not positively-linked, and  $(2, 4)$  and  $(3, 4)$  are right critical vertices of type I of  $\left( \begin{pmatrix} 1112223 \\ 7654321 \end{pmatrix}, \begin{pmatrix} 1 \\ 4 \end{pmatrix} \right)$  and of  $\left( \begin{pmatrix} 2233 \\ 7542 \end{pmatrix}, \begin{pmatrix} 1 \\ 4 \end{pmatrix} \right)$ , respectively.

In the next lemmas, we analyse the relationship between the critical vertices of our chosen biword  $\Pi$  and the matrices defined in Step 2 of the Main Algorithm. This analysis is important in order to prove that the invariant partition of  $B_1, \dots, B_t$  is the shape of  $K$ .

**Lemma 5.7.** *For  $q = 1, r$ , let  $(x_q, k)$  and  $(y_q, k + \varepsilon_q)$  be a pair of linked vertices belonging to the lattice path  $\mathbf{u}_q$ , with  $x_1 < x_r$ , and consider the matrices  $I + S(x_q, s(y_q, k + \varepsilon_q), \nu_q \sigma_k) = T_{i_q j_q}(1)$ . Then  $j_1 = j_r$  if and only if  $(x_r, k)$  is a right critical vertex of type I or II of  $(\mathbf{u}_r, \dots, \mathbf{u}_2, \mathbf{u}_1)$ ,  $k \geq 2$ , with critical components  $\mathbf{u}_{r-1}, \dots, \mathbf{u}_2$ .*

*Proof.* The *if* part. Without loss of generality assume that  $(x_2, k)$  is a right critical vertex of type I or II of  $(\mathbf{u}_2, \mathbf{u}_1)$ , without critical components. Then  $x_1 = s(y_2, k + \varepsilon_2)$  and we may write  $\nu_2\sigma_k = \theta(x_1 s(y_1, k + \varepsilon_1))\nu_1\sigma_k$ , where  $\theta \in \mathfrak{S}_n$  satisfy  $\theta(x_1) = x_1$ . Since  $\nu_2\sigma_k(j_1) = x_1$  we get  $\nu_2\sigma_k(j_1) = \theta(x_1) = x_1$ , and the result follows.

The *only if* part. Without loss of generality, assume that in the matrix  $B_k$ , defined in steps 2 and 3 of the Main Algorithm, there are no matrices  $T_{i_j 1}(1)$  between  $T_{i_1 j_1}(1)$  and  $T_{i_2 j_1}(1)$ . Then, we may write  $\nu_2\sigma_k = \theta(x_1 s(y_1, k + \varepsilon_1))\nu_1\sigma_k$ , for some permutation  $\theta \in \mathfrak{S}_n$  satisfying  $\theta(x_1) = x_1$ . By definition, we must have  $\nu_2\sigma_k(j_1) = s(y_2, k + \varepsilon_2)$ . On the other hand,  $\nu_2\sigma_k(j_1) = \theta(x_1 s(y_1, k + \varepsilon_1))\nu_1\sigma_k(j_1) = x_1$ . Therefore, we find that  $x_1 = s(y_2, k + \varepsilon_2)$ , and this equality means that  $(x_2, k)$  is a right critical vertex of type I or II of  $(\mathbf{u}_2, \mathbf{u}_1)$ .  $\square$

**Lemma 5.8.** *Suppose  $m_k < m_{k+1}$ , and let  $(x, k)$  and  $(y, k+1)$  be consecutive vertices of the lattice path  $\mathbf{v}_l$ . Then  $I + S(x, y, \nu\sigma_k) = T_{ij}(1)$  with  $m_k < j \leq m_{k+1}$  if and only if there is a vertex  $(a, k+1)$  belonging to a lattice path  $\mathbf{v}_1$  which is a left critical vertex of  $(\mathbf{v}_1, \dots, \mathbf{v}_l)$ , with  $\sigma_{k+1}^{-1}(a) = j$ .*

*Proof.* The *if* part. Without loss of generality, assume that  $(x, k+1)$  is a left critical vertex of  $(\mathbf{u}, \mathbf{v})$ , with no critical components. Since  $\sigma_{k+1}([m_{k+1}]) = J_{k+1}$ , there is  $j \in [m_{k+1}]$  such that  $\sigma_{k+1}(j) = x$ . Consider the matrix  $S(x, y, \nu\sigma_k)$ . Clearly, we may write  $\sigma_{k+1} = \theta(x y)\nu\sigma_k$ , where  $\theta \in \mathfrak{S}_n$  satisfy  $\theta(x) = x$ . Therefore,  $\nu\sigma_k(j) = (x y)\theta^{-1}\sigma_{k+1}(j) = y$ , and thus  $I + S(x, y, \nu\sigma_k) = T_{ij}(1)$ , for some integer  $i \in [m_k]$ . Finally, note that since the lattice path  $\mathbf{v}$  has no vertex in column  $k$ , all vertices in column  $k$  must be positively-linked. Thus, by Lemma 5.5, we find  $m_k < j$ .

The *only if* part. Assume now that  $I + S(x, y, \nu\sigma_k) = T_{ij}(1)$ , for some integer  $m_k < j \leq m_{k+1}$ . Again, it is clear that we must have

$$\sigma_{k+1} = \rho(x_r x_{r-1}) \cdots (x_2 x_1)(x_1 x)(x y)\nu\sigma_k,$$

for some integers  $x_r > \cdots > x_1 > x > y$  and a permutation  $\rho \in \mathfrak{S}_n$  such that  $\rho(x_r) = x_r$ . Thus,  $\sigma_{k+1}(j) = x_r \in J_{k+1}$ , since  $j \in [m_{k+1}]$ . Denote by  $\mathbf{v}$  the lattice path containing the vertex  $(x_r, k+1)$ , and note that since  $\rho(x_r) = x_r$ ,  $\mathbf{v}$  cannot have a vertex in column  $k$ . Thus,  $(x_r, k+1)$  is a left critical vertex of  $(\mathbf{u}, \mathbf{v})$ .  $\square$

Compare this result with Example 5.2, where the subindex 2 of  $T_{12}(1) = I + S_{3,2}^2$  satisfies  $1 = m_1 < 2 = m_2$ , and  $(3, 2)$  is a left critical vertex of  $\left(\begin{pmatrix} 2233 \\ 7542 \end{pmatrix}, \begin{pmatrix} 1112223 \\ 7654321 \end{pmatrix}\right)$ .

We are now ready to start the proof of our Main Theorem.

*Proof of Theorem 4.4.* It was shown in Proposition 4.3 that  $A_0, B_1, \dots, B_t$  satisfy conditions I and II of Definition 4.1. So it remains to show that  $B_1 \cdots B_t$  is equivalent to the diagonal matrix  $D_{[m_1]} \cdots D_{[m_t]}$ , that is,  $B_1, \dots, B_t$  is a matrix construction of the key with weight  $(m_1, \dots, m_t)$ . Using Corollary 5.3, Lemma 5.7, and the commutation rules in Lemma 4.1, we may write each matrix

$$B_k = \prod_{l=1}^{q_{k-1}} (T_{i_l j_l}(1) T_{j_l i_l}(-1)) D_{[m_k]},$$

$k = 2, \dots, t$ , as

$$B_k = \left( \prod_{l=1}^{q_{k-1}} T_{i_l j_l}(1) \right) C_k D_k \left( \prod_{l=1}^{q_{k-1}} T_{j_l i_l}(-1) \right) D_{[m_k]},$$

where each  $C_{k+1}$  [respectively  $D_{k+1}$ ] is a product of upper [respectively lower] triangular elementary matrices  $T_{ij}(1)$ , with  $i, j \in \{i_1, \dots, i_{q_k}\}$ ,  $i \neq j$ . Moreover, note that  $T_{ij}(1)$  is a factor of  $C_{k+1}D_{k+1}$  only if there are lattice paths  $\mathbf{u}_l$  with vertices  $(a_l, k)$  such that  $\sigma_k^{-1}(a_l) = l$ , for  $l = i, j$ , and where  $(a_i, k)$  is a right critical vertex of type I or II of  $(\mathbf{u}_i, \mathbf{u}_j)$ . We will refer to  $T_{i_l j_l}(1)$  and  $T_{j_l i_l}(-1)$ ,  $l = 1, \dots, q_{k-1}$ , as original matrices of  $B_k$ , and we call the triangular matrices in  $C_k D_k$  secondary matrices.

Denote by  $\mathbf{u}_l$  the lattice path for which  $\sigma_0^{-1}(u_l^0) = l$ , with  $u_l^0$  the column index of its initial vertex, for  $l \in [\ell_1 + \dots + \ell_t]$ . Next, for  $k = 2, \dots, t$ , using Lemma 5.8 write

$$(5.6) \quad B_k = \left( \prod_{l=1}^{q_{k-1}} T_{i_l j_l}(1) \right) C_k D_k D_{[m_k]} \left( \prod_{l=1}^{q_{k-1}} T_{j_l i_l}(\phi_l) \right),$$

where  $\phi_l = -1$  if  $T_{j_l i_l}(-1) = S(x, y, \nu\sigma_k)$ , if  $(x, k)$  and  $(y, k+1)$  are linked vertices of  $\mathbf{u}_{i_l}$  and there is a vertex  $(a, k+1) \in \mathbf{u}_{j_l}$ , which is a left critical vertex of  $(\mathbf{u}_{j_l}, \mathbf{u}_{i_l})$ ; otherwise  $\phi_l = -p$ .

The next step consists in using the commutation rules in Lemma 4.1 and Corollary 5.6 to eliminate, by left equivalence, the left most upper triangular matrix  $T_{ij}(x)$  in  $D_{[m_1]}B_2 \cdots B_t$ . Notice that this operation may generate new matrices  $T_{uv}(y)$ . Repeat this operation, eliminating all upper triangular matrices.

By Lemma 4.1, there exists always a unimodular elementary matrix  $E$  such that

$$T_{ji}(\tau)T_{ab}(\tau') = T_{ab}(\varrho')T_{ji}(\varrho)E, \text{ with } \varrho, \tau \text{ and } \varrho', \tau' \text{ pairs of associated elements,}$$

whenever  $\tau, \tau'$  are not both unities, or  $(a, b) \neq (i, j)$  if  $i \neq j$ . Thus, *the matrix operations described in the previous paragraph are feasible if they do not involve the passage of a matrix  $T_{ij}(\tau)$ , with  $i < j$ , to the left of a matrix  $T_{ji}(\tau')$ , with  $\tau = -\tau' = \pm 1$ .*

We begin by noticing that  $D_{[m_l]}T_{ij}(1) = T_{ij}(p)D_{[m_l]}$ , whenever  $m_l < j$ . Thus, the situation described above only happens if the diagonal matrices, say  $D_{[m_{k+1}]}, \dots, D_{[m_{k+\varepsilon}]}$ , between  $T_{ji}(\tau')$  and  $T_{ij}(\tau)$  satisfy  $j \leq m_{k+1}, \dots, m_{k+\varepsilon}$ . If  $T_{ji}(\tau') = I - S(x, y, \nu\sigma_{k-1})^T$  is a original matrix of  $B_k$ , (5.6) implies the existence of a left critical vertex  $(a, k+1) \in \mathbf{u}_j$  of  $(\mathbf{u}_j, \mathbf{u}_i)$ . If  $T_{ji}(\tau')$  is a secondary matrix of  $B_{k+1}$ , by Lemma 5.7, there is a right critical vertex of type I of  $\mathbf{u}_j$ , in column  $k$ , of  $(\mathbf{u}_j, \mathbf{u}_i)$ . By (5.6) and Lemma 4.1, any other case implies the existence of a sequence of matrices

$$(5.7) \quad T_{i_0 i_1}(\tau_0), T_{i_1 i_2}(\tau_1), \dots, T_{i_s i_{s+1}}(\tau_s),$$

with  $i_l \leq m_{k+1}, \dots, m_{k+\varepsilon}$ , for  $l = 0, 1, \dots, s+1$ ,  $s \geq 1$ , which generates the matrix  $T_{ji}(\tau')$  by the process of elimination described above. By Corollary 5.6, and Lemmas 5.7 and 5.8, the matrices in (5.7) are the result of left or right critical vertices of type I. In particular, we find that  $\mathbf{u}_{i_0}$  is below  $\mathbf{u}_{i_1}$  in columns  $k+1, \dots, k+\varepsilon$ ,  $\mathbf{u}_i$  is above  $\mathbf{u}_{i_1}$  in columns  $k+\zeta, \dots, k+\varepsilon$  and  $\mathbf{u}_j$  is below  $\mathbf{u}_{i_0}$  in columns  $k+\xi, \dots, k+\varepsilon$ , for some  $\varepsilon \geq \zeta, \xi \geq 1$ . In any case, the existence of  $T_{ji}(\tau')$  means that  $\mathbf{u}_j$  is below  $\mathbf{u}_i$  in columns  $k+\varepsilon', \dots, k+\varepsilon$ , for some  $\varepsilon \geq \varepsilon' \geq 1$ .

The same reasoning shows that the existence of  $T_{ij}(\tau)$  means that  $\mathbf{u}_i$  is below  $\mathbf{u}_j$  in column  $k + \varepsilon$ , contradicting Proposition 5.4. Therefore, whenever  $T_{ji}(\tau)$ ,  $i < j$ , is on the left of  $D_{[m_{k+1}]}$  and satisfies  $j \leq m_{k+1}$ , any elementary matrix  $T_{ij}(x)$  is transformed, during the process of its elimination by left equivalence, into  $T_{ij}(xp)$  when reaching the diagonal matrix  $D_{[m_{k+1}]}$ . We may thus eliminate all upper triangular matrices, and obtain

$$(5.8) \quad B_1 \cdots B_t \sim_E D_{[m_1]} T_2 D_{[m_2]} \cdots T_t D_{[m_t]},$$

where each  $T_k$  is a product of lower elementary triangular matrices,  $k = 2, \dots, t$ . Finally, using Lemma 4.1, together with Corollary 5.6, we may eliminate by right equivalence all these lower elementary triangular matrices, beginning with the right most triangular matrix in (5.8), and ending in the left most triangular matrix in  $T_2$ . Therefore, (5.8) is equivalent to  $D_{[m_1]} \cdots D_{[m_t]}$ .  $\square$

## 6. MATRIX TABLEAU-PAIRS WITH KEY CONDITION OVER A FOUR-LETTER ALPHABET

**6.1. An injection of the Knuth class of a key over the four-letter alphabet into a key over the six-letter alphabet.** According to Theorem 2.1, in the case of a four-letter alphabet, the problem of a matrix construction of a tableau-pair  $(T, K)$ , where  $T$  rectifies to the key-tableau  $K$ , is reduced to the key-tableau  $K(\sigma, (\ell_4, \dots, \ell_1))$  with  $\sigma \in S$ ,  $\ell_2 > 0$ , and  $\ell_4 > 0$ , since its Knuth class is the union of the sets  $\sqcup (\widehat{v}_5^{n_5}, v_4^{\ell_4 - n_5}, v_3^{\ell_3}, v_2^{\ell_2 - n_5}, v_1^{\ell_1})$  for  $0 \leq n_5 \leq \min\{\ell_2, \ell_4\}$ . Recall, from Section 2, that  $\widehat{v}_5 = 431421$  and  $S = \{1423, 1432, 4123, 4132\}$ .

Consider  $w \in \sqcup (\widehat{v}_5^{n_5}, v_4^{n_4}, \dots, v_1^{n_1})$ , for some  $0 \leq n_5 \leq \min\{\ell_2, \ell_4\}$ ,  $n_i = \ell_i$ ,  $i = 1, 3$ , and  $n_i = \ell_i - n_5$ ,  $i = 2, 4$ . Let  $\{X_5^{n_5}, \dots, X_5^1, \dots, X_4^{n_4}, \dots, X_4^1, \dots, X_3^{n_3}, \dots, X_3^1, \dots, X_2^{n_2}, \dots, X_2^1, \dots, X_1^{n_1}, \dots, X_1^1\}$  be a set partition of  $[\sum_{i=1}^4 n_i + 6n_5]$ , with  $w|X_j^i = v_j$ ,  $i \in [n_j]$ ,  $j \in \{1, 2, 3, 4\}$ , and  $w|X_5^i = \widehat{v}_5$ ,  $i \in [n_5]$ . Let  $I^{j,i}$  be a row of length  $j$ , for  $i \in [n_j]$ ,  $j = 1, \dots, 4$ , and  $I^{5,i} = a^i b^i c^i d^i e^i f^i$ , with  $a^i \leq b^i \leq c^i < d^i \leq e^i \leq f^i$ , for  $i \in [n_5]$ . We consider the biword

$$(6.1) \quad \Pi^+ = \left( \begin{array}{cccc|cccc} I^{5,n_5} & \cdots & I^{5,1} & I^{4,n_4} & \cdots & I^{4,1} & \cdots & I^{1,n_1} & \cdots & I^{1,1} \\ \widehat{v}_5 & \cdots & \widehat{v}_5 & v_4 & \cdots & v_4 & \cdots & v_1 & \cdots & v_1 \end{array} \right)$$

with bottom row  $\widehat{v}_5^{n_5} v_4^{n_4} \cdots v_1^{n_1}$ . The lattice path representation  $\Pi^+$  as in (6.1) consists of  $n_5 + \cdots + n_1$  nonintersecting lattice paths: the  $n_j$  lattice paths  $\begin{pmatrix} I^{j,i} \\ v_j \end{pmatrix}$  of length  $j$  have nonnegative slope, for  $i \in [n_j]$ ,  $j \in \{1, 2, 3, 4\}$ , and the  $n_5$  lattice paths  $\begin{pmatrix} I^{5,i} \\ \widehat{v}_5 \end{pmatrix}$ ,  $i \in [n_5]$ , of length 6, are such that the line linking the vertices  $\begin{pmatrix} c^i \\ 1 \end{pmatrix}$  and  $\begin{pmatrix} d^i \\ 4 \end{pmatrix}$  has negative slope.

Given a biword  $\Pi^+$  (6.1) it is a simple task to adjust, if necessary, the links between the vertices of this biword in order to form a new biword  $\widehat{\Pi}^+$  satisfying the conditions of the following lemma. (To avoid cumbersome notation we shall drop the super indices in this statement.)

**Lemma 6.1.** *There is a biword  $\widehat{\Pi}^+$  equivalent to  $\Pi^+$  as in (6.1) satisfying the following conditions:*

(a) *Any two lattice paths have no critical vertices.*

(b) *If  $\begin{pmatrix} abcdef \\ 431421 \end{pmatrix}$  is a lattice path of  $\widehat{\Pi}^+$  then:*

(i) *in row  $c$  there is at most one more vertex, placed in column 4, which must be negatively-linked to a vertex in column 3;*

(ii) *if  $d \neq e$ , in row  $d$  there is at most one more vertex, placed in column 1, which must be positively-linked;*

(iii) *if  $d \neq e$ , in row  $e$ , to the right of  $(e, 2)$ , there are no vertices.*

*Proof.* We prove only condition (b)(iii). All other conditions are proven in a similar way. Recall that we must have  $a \leq b \leq c < d \leq e \leq f$ . If  $(e, 3)$  is also a vertex of  $\Pi^+$ , then it must belong to one of the words 431421, 4321 or 431. In these cases,  $\Pi^+$  must have a sub-biword either of the form

$$\Gamma = \begin{pmatrix} abcdef & gehjkl \\ 431421 & 431421 \end{pmatrix}, \text{ or } \Delta = \begin{pmatrix} abcdef & gehi \\ 431421 & 4321 \end{pmatrix}, \text{ or } \Lambda = \begin{pmatrix} abcdef & geh \\ 431421 & 431 \end{pmatrix},$$

with  $g \leq e \leq h \leq i < j \leq k \leq l$ . We may re-link the vertices of the lattice paths of  $\Gamma$ ,  $\Delta$ , or  $\Lambda$  such that they are replaced in  $\Pi^+$  by

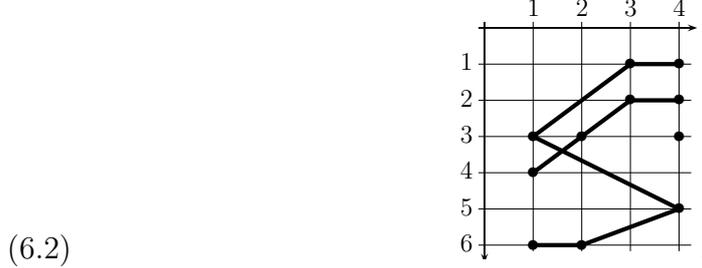
$$\begin{pmatrix} abcjkl & gee f & dh \\ 431421 & 4321 & 41 \end{pmatrix}, \begin{pmatrix} abcdhi & gee f \\ 431421 & 4321 \end{pmatrix}, \text{ or } \begin{pmatrix} gee f & abc & dh \\ 4321 & 431 & 41 \end{pmatrix},$$

respectively. In either case, the vertex  $(e, 2)$  belongs now to the new lattice path  $\begin{pmatrix} gee f \\ 4321 \end{pmatrix}$ , and it is linked to  $(e, 3)$ . We have only changed the links in the two lattice paths. Therefore, we may assume, without loss of generality, that, if  $(e, 2)$  is a vertex of a lattice path  $\begin{pmatrix} I^{5,i} \\ \widehat{v}_5 \end{pmatrix}$ , there is no vertex in position  $(e, 3)$ .

Assume now that  $(e, 4)$  is a vertex, but  $(e, 3)$  is not. An analysis similar to the one done above shows that if  $(e, 4)$  is a vertex of  $\begin{pmatrix} I^{5,i} \\ \widehat{v}_5 \end{pmatrix}$  or  $\begin{pmatrix} I^{j,i} \\ v_j \end{pmatrix}$ , for some  $j = 1, \dots, 4$ , then we may re-link the vertices of the corresponding lattice paths in such a way that the vertex  $(e, 2)$  is linked to  $(e, 4)$ . Therefore, we may assume that  $(e, 4)$  is not a vertex of  $\Pi^+$ .  $\square$

**Example 6.1.** The biword  $\Pi^+ = \begin{pmatrix} 113566 & 2234 & 3 \\ 431421 & 4321 & 4 \end{pmatrix}$ , whose lattice path representation is given below, fails to satisfy condition (b)(i) of the lemma above, since the vertex  $(3, 1)$  belongs to the lattice path of  $\begin{pmatrix} 113566 \\ 431421 \end{pmatrix}$ , and there is a vertex in

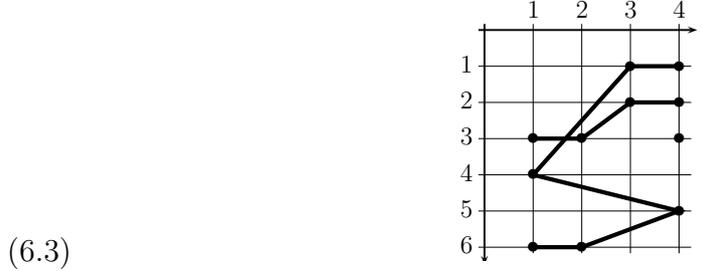
row 3, column 2



Rearranging the links between the vertices, we obtain the biword

$$\widehat{\Pi}^+ = \begin{pmatrix} 114566 & 2233 & 3 \\ 431421 & 4321 & 4 \end{pmatrix},$$

which satisfy all the required conditions of Lemma 6.1



In what follows, we assume that our biword  $\Pi^+$  satisfies the conditions of Lemma 6.1.

**Corollary 6.2.** (a) If  $(a, i)$  is the initial vertex of a lattice path of length  $k$ , then  $s(a, i) \neq a$  only if  $\sigma \in \{4123, 4132\}$ ,  $i = 4$ ,  $k = 1$ , and the nearest vertex in row  $a$  is either in column 1, which is the rightmost vertex of the lattice path  $\begin{pmatrix} I^{5,i} \\ \widehat{v}_5 \end{pmatrix}$  or of  $\begin{pmatrix} I^{j,i} \\ v_j \end{pmatrix}$ , for some  $j = 1, \dots, 4$ , or in column 2, which is a vertex of  $\begin{pmatrix} I^{4,i} \\ v_4 \end{pmatrix}$ .

(b) If  $(b, j)$ ,  $(a, i)$  are consecutive vertices of a lattice path of length  $k$ , then  $s(a, i) \neq a$  only if  $j = 1$ ,  $i = 4$ , and the nearest vertex in row  $a$  is in column 2, and it is a vertex of  $\begin{pmatrix} I^{4,i} \\ v_4 \end{pmatrix}$ .

*Proof.* This follows from Lemma 6.1 (a). □

For each  $\sigma = \sigma(1)\sigma(2)\sigma(3)\sigma(4)$  in  $\mathfrak{S}_4$ , consider the permutation

$$\bar{\sigma}(1)\bar{\sigma}(2)\bar{\sigma}(3)\bar{\sigma}(4)34 \in \mathfrak{S}_6,$$

where

$$\bar{\sigma}(k) = \begin{cases} \sigma(k), & \text{if } \sigma(k) = 1, 2, \\ \sigma(k) + 2, & \text{if } \sigma(k) = 3, 4. \end{cases}$$

This correspondence is a bijection between  $\mathfrak{S}_4$  and the set  $\{\bar{\sigma} = \alpha 34 \in \mathfrak{S}_6 : \alpha \in \mathfrak{S}_{\{1256\}}\} \subseteq \mathfrak{S}_6$ . In particular,  $S = \{1423, 1432, 4123, 4132\}$  is transformed into

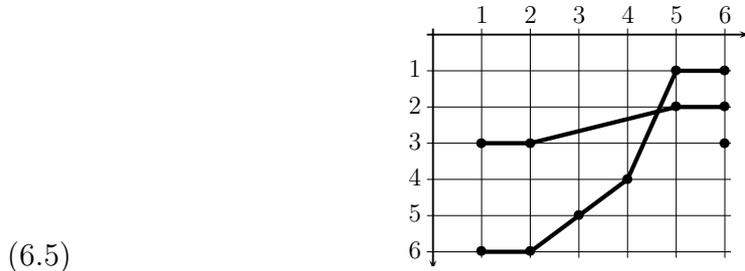
$\{162534, 165234, 612534, 615234\} \subseteq \mathfrak{S}_6$ . Let  $\rho : \{\widehat{v}_5, v_j : j \in [4]\} \rightarrow \{\bar{v}_j : j \in [6]\}$ , such that  $\rho(\widehat{v}_5) = \bar{v}_6$  and  $\rho(v_j) = \bar{v}_j$ ,  $j = 1, 2, 3, 4$ , where  $\bar{v}_6, \dots, \bar{v}_1$  are the columns of the key encoded by the permutation  $\bar{\sigma}$  and the vector  $(1^6)$ .

For each  $n_5 = 0, \dots, \min\{\ell_2, \ell_4\}$ , the map  $\rho$  can be extended, by shuffling, to a bijection between the sets  $\sqcup\sqcup(\widehat{v}_5^{n_5}, v_4^{n_4}, \dots, v_1^{n_1})$  and  $\sqcup\sqcup(\bar{v}_6^{n_5}, \bar{v}_4^{n_4}, \dots, \bar{v}_1^{n_1})$ . This last set is a subset of the plactic class of the key encoded by the permutation  $\bar{\sigma}$  and the vector  $(n_5, 0, n_4, n_3, n_2, n_1)$ . Thus every word in the plactic class of  $K(\sigma\alpha(m))$ ,  $\sigma \in \mathfrak{S}_4$ , has a copy in the shuffle of the columns of some key over the alphabet  $[6]$ . Let  $w \in \sqcup\sqcup(\widehat{v}_5^{n_5}, v_4^{n_4}, \dots, v_1^{n_1})$ . Fix an equivalent biword  $\Pi^+$  as in (6.1) satisfying the conditions of Lemma 6.1, and let

$$(6.4) \quad \bar{\Pi}^+ = \begin{pmatrix} I^{5,n_5} & \dots & I^{5,1} & I^{4,n_4} & \dots & I^{4,1} & \dots & I^{1,n_1} & \dots & I^{1,1} \\ \bar{v}_6 & \dots & \bar{v}_6 & \bar{v}_4 & \dots & \bar{v}_4 & \dots & \bar{v}_1 & \dots & \bar{v}_1 \end{pmatrix}$$

be the biword that we obtain when we transform each lattice path  $\begin{pmatrix} I^{j,i} \\ v_j \end{pmatrix}$  and  $\begin{pmatrix} a^i b^i c^i d^i e^i f^i \\ 431421 \end{pmatrix}$  of  $\Pi^+$  into  $\begin{pmatrix} I^{j,i} \\ \bar{v}_j \end{pmatrix}$  and  $\begin{pmatrix} a^i b^i c^i d^i e^i f^i \\ 654321 \end{pmatrix}$ , respectively, where  $a^i \leq b^i \leq c^i < d^i \leq e^i \leq f^i$ . Let  $\bar{\Sigma}$  be the biword obtained by sorting the biletters of  $\bar{\Pi}^+$  by weakly increasing rearrangement with respect to the anti-lexicographic order with priority in the first row. The second row of  $\bar{\Sigma}$ , denoted  $\rho(w)$ , is a shuffle of columns  $\bar{v}_j$ ,  $j = 6, 4, 3, 2, 1$ . Under this injection the lattice path  $\begin{pmatrix} a^i b^i c^i d^i e^i f^i \\ 431421 \end{pmatrix}$ , which has some negative slope steps, is stretched northeastward to  $\begin{pmatrix} a^i b^i c^i d^i e^i f^i \\ 654321 \end{pmatrix}$ , where now every step has nonnegative slope.

**Example 6.2.** Let  $\sigma = 4123$  and consider the biword  $\widehat{\Pi}^+ = \begin{pmatrix} 114566 & 2233 & 3 \\ 431421 & 4321 & 4 \end{pmatrix}$  in Example 6.1. We have  $\bar{\sigma} = 612534$ . Applying the map  $\rho$ , we obtain the biword  $\widehat{\bar{\Pi}}^+ = \begin{pmatrix} 114566 & 2233 & 3 \\ 654321 & 6521 & 6 \end{pmatrix}$ , whose lattice path representation is given below



The bottom word of the biword  $\bar{\Sigma}$ , defined above, and equivalent to  $\widehat{\bar{\Pi}}^+$ , is  $\rho(w) = 65656214321$ , a shuffle of the columns 654321, 6521 and 6, which are exactly the words read along the lattice paths, now all having nonnegative slope.

In the next proposition, we state some properties of the biword  $\bar{\Pi}^+$ .

**Proposition 6.3.** *Let  $\bar{\Pi}^+$  be the biword defined above, and let  $\begin{pmatrix} abcdef \\ 654321 \end{pmatrix}$  be one of its lattice paths. Then*

(a)  $\bar{\Pi}^+$  has no critical vertices.

(b) There are no vertices in row  $c$  to the left of  $(c, 4)$ , nor are there in row  $d$  to the right of  $(d, 3)$ , nor in row  $e$  to the right of  $(e, 2)$ .

*Proof.* This is a consequence of the definition of  $\bar{\Pi}^+$  and Lemma 6.1.  $\square$

We are now in position to prove the following result.

**Theorem 6.4.** *Let  $T$  be a Young tableau over a four-letter alphabet with rectification  $P$ , and let  $K$  be the key with the same weight as  $T$ . Then,  $(T, K)$  is a matrix tableau-pair if and only if  $P = K$ .*

*Proof.* The “only if” part was proved in [6]. Conversely, assume  $K = K(\sigma\alpha(m))$ , for some  $\sigma \in \{1423, 1432, 4123, 4132\}$  and  $\ell_4, \ell_2 > 0$ . Recall that in any other case the plactic class of  $K(\sigma\alpha(m))$  is the set of all shuffles of its columns and, thus, by Theorem 4.4,  $(T, K(\sigma\alpha(m)))$  is a matrix-tableau pair.

Without loss of generality let  $w$  be the word of  $T$  in the set  $\sqcup\sqcup(\widehat{v}_5^{n_5}, v_4^{n_4}, \dots, v_1^{n_1})$ , for some  $0 < n_5 \leq \min\{\ell_2, \ell_4\}$ , where  $n_i = \ell_i$ ,  $i = 1, 3$ , and  $n_i = \ell_i - n_5$ ,  $i = 2, 4$ . Notice that either  $m_1 = 2n_5 + n_4 + n_3 + n_2 + n_1$  and  $m_4 = 2n_5 + n_4 + n_3 + n_2$ , if  $\sigma \in \{1423, 1432\}$ , or  $m_4 = 2n_5 + n_4 + n_3 + n_2 + n_1$  and  $m_1 = 2n_5 + n_4 + n_3 + n_2$  otherwise. Let  $\Pi^+$  be the biword (6.1) of  $T$  satisfying Lemma 6.1. Consider the injection  $\rho$  and the corresponding biword  $\bar{\Pi}^+$  as in (6.4), whose bottom word is the key  $K(\bar{\sigma}, (n_5, 0, n_4, n_3, n_2, n_1))$ . Let  $\bar{T}$  be the Young tableau with the same skew-shape as  $T$  and biword  $\bar{\Pi}^+$ . For each  $i = 1, \dots, n_5$ , let  $\begin{pmatrix} a^i b^i c^i d^i e^i f^i \\ 6\ 5\ 4\ 3\ 2\ 1 \end{pmatrix}$  be a lattice

path of  $\bar{\Pi}^+$ . Then, if  $J_4, J_3, J_2, J_1$  are the indexing sets of  $T$ , and if  $\bar{J}_6, \dots, \bar{J}_1$  are the indexing sets of  $\bar{T}$ , we have  $\bar{J}_1 = J_1 \setminus \{c^i : i \in [n_5]\}$ ,  $\bar{J}_2 = J_2$ ,  $\bar{J}_5 = J_3$ ,  $\bar{J}_3 = \{d^i : i \in [n_5]\}$ ,  $\bar{J}_4 = \{c^i : i \in [n_5]\}$ , and  $\bar{J}_6 = J_4 \setminus \{d^i : i \in [n_5]\}$ . Note also that, by Proposition 6.3, there are no vertices in row  $c^i$  to the left of  $(c^i, 4)$ , nor are there in row  $d^i$ , to the right of  $(d^i, 3)$ . Then, we may apply the Main Algorithm to  $\bar{\Pi}^+$ , choosing a permutation  $\sigma_1 \in \mathfrak{S}_n$  satisfying the conditions of Step 1, and satisfying in addition that  $\sigma_1(n_1 + \dots + n_5 + i) = c^i$ , for  $i = 1, \dots, n_5$ . Denote by  $A_0, B_1, \dots, B_6$  the sequence of matrices obtained by this procedure. Then  $A_0 = \text{diag}_p(a')$ ,  $B_1 = P_{\sigma_1} D_{[m_1 - n_5]}$ ,  $B_2 = S_2 D_{[m_2]}$ ,  $B_k = S_k D_{[n_5]}$ , for  $k = 3, 4$ ,  $B_5 = S_5 D_{[m_3]}$ , and  $B_6 = S_6 D_{[m_4 - n_5]}$ . By Theorem 4.4, this sequence is a matrix realization for the pair  $(\bar{T}, K(\bar{\sigma}, (n_5, 0, n_4, n_3, n_2, n_1)))$ . In particular, this means that

$$(6.6) \quad \sigma_1([m_1 - n_5]) = J_1 \setminus \{c^i : i \in [n_5]\}, \quad \sigma_2([m_2]) = J_2,$$

$$(6.7) \quad \sigma_3([n_5]) = \{d^i : i \in [n_5]\}, \quad \sigma_4([n_5]) = \{c^i : i \in [n_5]\},$$

$$(6.8) \quad \sigma_5([m_3]) = J_3, \quad \text{and} \quad \sigma_6([m_4 - n_5]) = J_4 \setminus \{d^i : i \in [n_5]\}.$$

Consider the sequence of matrices  $A'_0, B'_1, B'_2, B'_3, B'_4$ , defined by  $A'_0 = \text{diag}_p(a')$ ,  $B'_1 := P_{\sigma_1} D_{\Gamma_1}$ ,  $B'_2 := B_2$ ,  $B'_3 := S_3 S_4 S_5 D_{[m_3]}$ , and  $B'_4 := S_6 D_{\Gamma_4}$ , where  $\Gamma_1 =$

$[m_1], \Gamma_4 = [m_4 - n_5] \cup \{m_4 + n_1 + 1, \dots, m_4 + n_1 + n_5\}$  if  $\sigma \in \{1423, 1432\}$ , and  $\Gamma_1 = [m_1 - n_5] \cup \{m_1 + n_1 + 1, \dots, m_1 + n_1 + n_5\}, \Gamma_4 = [m_4]$  otherwise.

By (6.6) and the definition of  $\sigma_1$ , it is clear that

$$A'_0 B'_1 = \text{diag}_p(a') P_{\sigma_1} D_{\Gamma_1} = \text{diag}_p(a' + \chi^{J_1}) P_{\sigma_1},$$

and

$$A'_0 B'_1 B'_2 \sim_L \text{diag}_p(a' + \chi^{J_1}) P_{\sigma_2} D_{[m_2]} = \text{diag}_p(a' + \chi^{J_1} + \chi^{J_2}) P_{\sigma_2}.$$

Next, by (6.7) and (6.8), we find that

$$A'_0 B'_1 B'_2 B'_3 \sim_L \text{diag}_p(a' + \chi^{J_1} + \chi^{J_2}) P_{\sigma_5} D_{[m_3]} = \text{diag}_p(a' + \chi^{J_1} + \chi^{J_2} + \chi^{J_3}) P_{\sigma_5}.$$

Finally, consider the product

$$A'_0 B'_1 B'_2 B'_3 B'_4 \sim \text{diag}_p(a' + \chi^{J_1} + \chi^{J_2} + \chi^{J_3}) P_{\sigma_6} D_{\Gamma_4},$$

and notice that, for each  $i = 1, \dots, n_5$ , we may write  $\sigma_6 = \alpha_2^i (c^i d^i) \alpha_1^i$ , with  $(c^i d^i)$  the transposition of  $c^i$  with  $d^i$ , for some  $\alpha_2^i, \alpha_1^i \in \mathfrak{S}_n$  satisfying  $\alpha_1^i(c^i) = c^i$  and  $\alpha_2^i(d^i) = d^i$ , since by Lemma 6.3 there are no vertices of  $\overline{\Pi}^+$  nor are there in row  $c^i$  to the left of  $(c^i, 4)$ , nor in row  $d^i$  to the right of  $(d^i, 3)$ . Therefore,

$$\sigma_6(\{m_4 + n_1 + 1, \dots, m_4 + n_1 + n_5\}) = \{d^i : i \in [n_5]\},$$

and, by (6.8), we must have

$$A'_0 B'_1 B'_2 B'_3 B'_4 \sim_L \text{diag}_p(a' + \chi^{J_1} + \chi^{J_2} + \chi^{J_3} + \chi^{J_4}).$$

Thus,  $A'_0, B'_1, B'_2, B'_3, B'_4$  satisfy conditions I and II of Definition 4.1. It remains to show that  $B'_1 \cdots B'_4$  is equivalent to the diagonal matrix  $D_{[m_1]} \cdots D_{[m_4]}$ . We start by using Lemma 4.1 to write, for  $k = 2, 3, 4$ ,

$$B'_k = \prod_{l=1}^{q_{k-1}} T_{i_l j_l}(1) C_k D_k \prod_{l=1}^{q_{k-1}} T_{j_l i_l}(-1) D_{\Gamma_k},$$

where  $\Gamma_l = [m_l]$  for  $l = 2, 3$ ,  $C_k D_k$  is the identity matrix for  $k = 2, 4$ , and  $C_3$  [respectively  $D_3$ ] is a product of upper [respectively lower] elementary matrices  $T_{jj'}(\tau)$ , for some integers  $j, j' \in [m_2]$ , and  $\tau = \pm 1$ . Next, use again Lemma 4.1 to write  $B'_1 \cdots B'_4$  as

$$(6.9) \quad \prod_{l=1}^{q_1} T_{i_l j_l}(p^{\nu_l}) D_{[\Gamma_1]} \prod_{l=1}^{q_1} T_{j_l i_l}(-1) \prod_{l=1}^{q_2} T_{i_l j_l}(p) D_{[m_2]} C_3 D_3 \prod_{l=1}^{q_2} T_{j_l i_l}(-1) \\ \cdot \prod_{l=1}^{q_3} T_{i_l j_l}(p) D_{[m_3]} \prod_{l=1}^{q_3} T_{j_l i_l}(-1) D_{[\Gamma_4]},$$

for some  $\nu_l \geq 0$ . Notice that we may eliminate any upper triangular matrix  $T_{ij}(p^\nu)$ ,  $\nu > 0$ , by left equivalence, using Lemma 4.1. This operation may create new elementary matrices  $T_{uv}(p^{\nu'})$ , but, as we have mentioned, using Lemma 4.1, we may assume without loss of generality that this is not the case. Thus, we may write

$$(6.10) \quad (6.9) \sim_L D_{[\Gamma_1]} \prod_{l=1}^{q_1} T_{j_l i_l}(-1) D_{[m_2]} C_3 D_3 \prod_{l=1}^{q_2} T_{j_l i_l}(-1) D_{[m_3]} \prod_{l=1}^{q_3} T_{j_l i_l}(-1) D_{[\Gamma_4]}.$$

Use again Lemma 4.1 to eliminate all upper triangular matrices in  $C_3$  by left equivalence. New elementary matrices may be created. Among these, eliminate all those by left equivalence which are upper triangular. Finally, starting from right and moving to the left, eliminate all lower triangular matrices left in the product by right equivalence. It is now clear that  $B_1 \cdots B_4$  is equivalent to the diagonal matrix  $D_{[\Gamma_1]} D_{[m_2]} D_{[m_3]} D_{[\Gamma_4]}$ . Since  $m_k \leq n_5 + n_4 + n_3 + n_2$ ,  $k = 2, 3$ , and  $\Gamma_1 \subseteq \Gamma_4$  or  $\Gamma_4 \subseteq \Gamma_1$ , we find that this last diagonal matrix is equivalent to  $D_{[m_1]} \cdots D_{[m_4]}$ .  $\square$

## 7. CONCLUDING REMARKS AND OPEN QUESTIONS

The Main Algorithm in Section 4 generates a matrix realization for the pair  $(T, K)$  whenever the word of  $T$  is a shuffle of the columns of the key-tableau  $K$ . In particular, when the Knuth class of  $K$  is the set of shuffles of its columns, we may construct a matrix realization for the pair  $(T, K)$  whenever  $T$  rectifies to  $K$ . Over a four-letter alphabet, there are key-tableaux whose plactic class is bigger than the set of the shuffles of its columns. However, in those cases, we just add the single word 431421 to the set of the columns of the key-tableau  $K$  in order to describe its Knuth class as the union of the shuffles of subsets of words comprising the columns of  $K$  and the word 431421.

Working out some examples gives us an indication that this procedure might be generalized to a  $t$ -letter alphabet,  $t \geq 5$ , once shuffling generators for the Knuth class are known. Unfortunately one rapidly sees that the number of words needed to describe the Knuth class of a key-tableau, for  $t \geq 5$ , as a set of shuffles containing the shuffles of its columns, increases very fast. For example, over the alphabet [5], we need to add the words 5415321 and 5431521 to the columns of the key-tableau  $K = 54321\ 51$  to describe its Knuth class. If we instead consider the key-tableau  $K = 54321\ 54321\ 51$ , we now need to add 17 new words to the columns of  $K$  in order to describe its Knuth class. Nevertheless, a simple adaptation of the procedure used over a four-letter alphabet works also fine.

Although it seems to be a difficult problem, it would be interesting to describe the Knuth class of a key-tableau  $K$  as the union of shuffles of subsets of words where the set of columns of  $K$  is included. With this description, one believes that the general problem of finding a matrix construction for the pair  $(T, K)$ , when  $T$  is Knuth equivalent to  $K$ , can be solved using a technique similar to that which we used over a four-letter alphabet. We point out that recently the enumeration of the distinct shuffles of two permutations of any given lengths has been provided in [8]. In particular, this result gives the number of distinct shuffles of the columns  $v_1$  and  $v_2$  in the Knuth class of the key  $K = v_2 v_1$  with only two columns. This number is exactly the number of elements of the Knuth class when the support of  $v_1$  is either an interval or is obtained from an interval by removing a single letter (see Theorem 4.1 in [6]).

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CMUC, DEPARTMENT OF MATHEMATICS, UNIVERSITY OF COIMBRA, 3001-454 COIMBRA,  
PORTUGAL